

Essays on Stochastic Inventory Systems

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Dedication

To my parents for the great support and continuous care.

To my wife Shan Gao and my son Matthew Chen for the unyielding love, support,
and encouragement.

Abstract

This thesis consists of three essays in stochastic inventory systems. The first essay is on the impact of input price variability and correlation on stochastic inventory systems. For a general class of such systems, we show that the expected cost function is concave in the input price. From this, it follows that higher input price variability in the sense of the convex order always leads to lower expected cost. We show that this is true under a wide range of assumptions for price evolution, including cases with i.i.d. prices and cases where prices are correlated and evolve according to an AR(1) process, a geometric Brownian motion, or a Markovian martingale. In addition, the result holds in cases where there is just a single period. We also examine the impact of price correlation over time and across inputs, and we find that expected cost is increasing in price correlation over time and decreasing in price correlation across components. We present results of a numerical study that provide insights on how various parameters influence the effects of price variability and correlation.

The second essay is on the optimal control of inventory systems with stochastic and independent leadtimes. We show that a fixed base-stock policy is sub-optimal and can perform poorly. For the case of exponentially distributed leadtimes, we show that the optimal policy is state-dependent and specified in terms of an inventory-dependent threshold function. Moreover, we show that this threshold function is non-increasing in the inventory level and characterized by at most m parameters. That is, once the threshold function starts to decrease it continues to

decrease with a rate that is at least one. Taking advantage of this structure, we develop an efficient algorithm for computing these parameters. In characterizing the structure of the optimal policy, we rely on an application of the Banach fixed point theorem. We compare the performance of the optimal policy to that of simpler heuristics. We also extend our analysis to systems with lost sales and systems with order cancellations.

The third essay is on the optimal policies for inventory systems with concave ordering costs. By extending the Scarf (1959) model to systems with piecewise linear concave ordering costs, we characterize the structure of optimal policies for periodic review inventory systems with concave ordering costs and general demand distributions. We show that, except for a bounded region, the generalized (s, S) policy is optimal. We do so by (a) introducing a conditional monotonicity property for the optimal order-up-to levels and (b) applying the notion of c -convexity. We also provide conditions under which the generalized (s, S) policy is optimal for all regions of the state space.

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Chapter 1

Introduction

This thesis consists of five chapters. In Chapter 2, 3 and 4, we present three completed research projects. In Chapter 5, we describes other ongoing research projects and future research directions. Chapters 2, 3 and 4 are self-contained, independent, and deal with separate topics. The following paragraphs are a brief summary of Chapters 2, 3, 4 and 5.

In Chapter 2, we explore the impact of input price variability in the context of an inventory system with stochastic demand and stochastic input prices. For a general class of such systems, we show that the expected cost function is concave in the input price. This implies that higher input price variability always leads to lower expected cost. We show that this is true under a wide range of assumptions for price evolution, including cases with i.i.d. prices and cases where prices are correlated and evolve according to an AR(1) process or a geometric Brownian motion. More significantly, we show that the result is true when prices evolve according to a Markovian martingale so that the expected price in the next period is equal to the realized price in the current period. This is perhaps surprising

Chapter 2

On the Impact of Input Price Variability and Correlation in Stochastic Inventory Systems

2.1 Introduction

Stochastic input prices are common in practice. The prices of raw materials, precious metals, grain commodities, and electronic components, among many others, can fluctuate considerably over short periods. Such fluctuations may result from variations in supply and demand, changes in market conditions, or the introduction of new technology. Firms in some industries face input price variability because of their reliance on spot markets for procurement and, in the case of firms with global supply chains, because of exchange rate fluctuation.

The presence of stochastic input prices raises several important questions.

First, how does the presence of variability in input prices affect input ordering decisions and the nature of the optimal ordering policy? Second, how does price variability affect performance, and particularly cost? Does higher price variability increase or decrease overall costs? How does price correlation, over time or across inputs, interact with price variability and what is the net effect on cost? Is the effect of price variability more pronounced with higher correlation?

There is literature dealing with inventory systems with stochastic input prices; see Zhang (2012) for a comprehensive review. In a periodic review inventory system, Kalymon (1971) considers a single-item model with setup costs in which future input prices are determined by a Markovian stochastic process, and establishes that the optimal policy is a price-dependent (s, S) policy. Golabi (1985) considers a problem with an independent price process, negligible setup cost, and deterministic demand. He shows that the optimal policy is to always purchase a quantity that covers demands for the next several periods, and that this number of periods is decreasing in the current price. Gavirneni (2004) develops an efficient recursive procedure to calculate the base stock level when there are no setup costs and shows that myopic solutions are very effective under a non-speculative assumption. For continuous review inventory systems, Song and Zipkin (1993) characterize the optimal policy and develop algorithms for settings with Markov modulated purchasing price and Markov modulated demand. Yang and Xia (2009) consider a problem in which the input price follows a discrete-state Markov process and demand is a compound Poisson process. They show that the optimal policy is of the order-up-to type and identify conditions under which the order-up-to levels are

of cost parameters evolves as a stochastic process. They show that if the single period cost is concave with respect to this vector, then the optimal cost is bounded above by the optimal cost for the dynamic program in which these stochastic cost parameters are replaced by their expectations in each period. However, the approach they employ cannot be used to compare two dynamic programs each with stochastic cost parameters. Boyabatlı et al. (2011) study optimal procurement, processing, and production policies for a meat-processing company which sources input through long-term contracts and from a spot market. They assume that the spot price follows a normal distribution and show that the optimal expected profit of the firm increases in the spot price variability under certain conditions.

In the economics literature, there is a stream of research that examines a firm's behavior when price or cost fluctuates. Sandmo (1971) and Batra and Ullah (1974) study the optimal output and input decisions for a competitive firm under price uncertainty and risk aversion. Anderson and Danthine (1981, 1983), Meyer (1987) and Kamara (1993) study how firms can use futures to hedge or speculate against price uncertainty. This literature relies on aggregate models of demand and supply and does not model operational decisions.

In this paper, we show that for a wide range of inventory problems and assumptions, higher input price variability (as measured by convex ordering of prices) leads to lower expected inventory costs over the planning horizon, where inventory costs include ordering, inventory holding, and shortage costs. One may initially attribute this phenomenon to the fact that higher variability affords more frequent opportunities to place large (small) orders in periods in which prices are anticipated to be higher (lower) in subsequent periods. Although we do

We let $y_t^*(s, x)$ denote a minimizer of (2.2). Then an optimal policy uses order-up-to level $y_t^*(s, x)$ if the state is (s, x) in period t , and the optimal order quantity is $y_t^*(s, x) - s$. The optimal expected total cost for the entire planning horizon (computed before learning the first ordering price) with starting inventory s is given by $V_1(s) = Ev_1(s, X_1)$.

In preparation for our analysis of the impact of input price variability, we next describe the form of the optimal policy for this inventory system. We begin with the following lemma.

Lemma 1. *The function $w_t(y, x)$ is convex in y for all x and $t = 1, \dots, T$.*

The proof of Lemma 1 (and all other proofs not provided in the paper) can be found in the appendix. Let $y_t^\circ(x)$ denote a minimizer of $w_t(y, x)$ over $y \in (-\infty, \infty)$. An optimal policy is described in the following proposition, which follows immediately from Lemma 1.

Proposition 1. *There exists an optimal ordering policy for the multi-period inventory system with stochastic input prices that is a state-dependent base stock policy with base stock levels $y_t^\circ(x)$. That is, $y_t^*(s, x) = \max\{s, y_t^\circ(x)\}$ and the optimal order quantity in state (s, x) at time t is $\max\{0, y_t^\circ(x) - s\}$.*

The optimal base stock level $y_t^\circ(x)$ need not be decreasing in the realized price x . For example, consider a case where $T = 2$, $b = 0.5$, $h = 0.5$, $D_1 = D_2 = 10$ and $X_2 = 2X_1 - 5$, and suppose that the marginal distribution for the ordering price in period 1 is $P(X_1 = 4) = P(X_1 = 6) = 0.5$ and thus the marginal distribution of the ordering price in period 2 is $P(X_1 = 3) = P(X_1 = 7) = 0.5$. In this case, it is easy to check that it is optimal to order nothing if the realized ordering price

in period 1 is 4 ($y_1^\circ(4) = 0$) and to order up to 20 if the realized ordering price in period 1 is 6 ($y_1^\circ(6) = 20$). Therefore, the optimal base stock is increasing with respect to the realized price. This is due to the strong positive correlation in the ordering price across periods. In the following proposition, we provide a sufficient condition under which this phenomenon does not occur and the base stock level is decreasing in the realized price.

Proposition 2. *If $E|f(\epsilon_t)| \leq 1$ for $t = 1, \dots, T$, then $y_t^\circ(x)$ is decreasing in x for $t = 1, \dots, T$.*

Examples that satisfy the condition $E|f(\epsilon_t)| \leq 1$ for $t = 1, \dots, T$ include the case of i.i.d. input prices and the case where the input prices evolve according to an AR(1) process. In the first case $f(\epsilon) = 0$, and in the second case $f(\epsilon) = \rho \in [-1, 1]$. If the condition in the proposition is not satisfied, for example, if $Ef(\epsilon_t) > 1$, then it is possible that a high (low) price in one period would lead to a even higher (lower) expected price in the next period. In this case, it may be optimal to order more (less) when the price is high (low). Or, if $Ef(\epsilon_t) < -1$, then a low price in one period (say, now) would lead to a high expected price in the next period and an even lower expected price after two periods. In that case, one may wish to order more now in anticipation of a high price in the next period but to order less now in anticipation of an even lower price after two periods. It is possible that the second of these two effects is stronger. Therefore, it is possible that it is optimal to order more as price increases.

2.3 Impact of Price Variability

In this section, we discuss the impact of input price variability on the expected total cost and show that higher variability yields lower expected total cost. In our analysis, we use the tool of convex ordering to compare different levels of price variability. A random variable X is said to be smaller than \widehat{X} in the *convex order* (written $X \leq_{cx} \widehat{X}$) if $Eu(X) \leq Eu(\widehat{X})$ for all convex functions $u(\cdot)$ such that the expectations exist. The concept of convex order is reviewed in, for example, Müller and Stoyan (2002) and Shaked and Shanthikumar (2007). If $X \leq_{cx} \widehat{X}$, then it is well known that $EX = E\widehat{X}$ and $\text{Var}(X) \leq \text{Var}(\widehat{X})$. For random variables drawn from various common distributions, convex ordering is equivalent to having ordered variances and identical means. For example, if we compare two normal random variables with the same mean, then the one with the smaller variance is smaller in the convex order. The same holds true for uniform, gamma and lognormal random variables as well. Below, we will frequently make use of the fact that if $u(\cdot)$ is concave and $X \leq_{cx} \widehat{X}$ then $Eu(X) \geq Eu(\widehat{X})$.

The next lemma establishes the concavity of the cost function $v_t(s, x)$ with respect to the ordering price x .

Lemma 2. $v_t(s, x)$ is concave in x for all s and $t = 1, \dots, T + 1$.

To study the impact of ordering price variability, we compare two different inventory systems with ordering price sequences $\{X_t\}$ and $\{\widehat{X}_t\}$ and noise sequences $\{\epsilon_t\}$ and $\{\widehat{\epsilon}_t\}$ satisfying $X_{t+1} = f(\epsilon_t)X_t + g(\epsilon_t)$ and $\widehat{X}_{t+1} = f(\widehat{\epsilon}_t)\widehat{X}_t + g(\widehat{\epsilon}_t)$ respectively. All other parameters of the two systems are the same. We assume that each of the two systems individually satisfies the assumptions after (2.1) in Section 2.2. Let $\widehat{v}_t(s, x)$ be the optimal expected total cost-to-go in period

as follows:

$$\pi'(s+n+i, k_l-j) = A(k_l-j)\pi'(s+n+i-1, k_l-j+1) + B(k_l-j)\pi'(s+n+i+1, k_l-j), \quad (3.11)$$

for $l = n, \dots, m-1$, $j = 1, \dots, k_l - k_{l+1}$ and $i = j, j-1, \dots, 1$, and

$$\pi'(s+n+a, k_{n+a}) = \begin{cases} \frac{k_n\mu}{\lambda}\pi'(s+n, k_n) - \sum_{j=k_{n+1}+1}^{k_n-1} \pi'(s+n+1, j) & \text{if } a = 1, k_{n+a} \neq 0, \\ \frac{\pi'(s+n+a-1, k_{n+a-1})}{B(k_{n+a-1})} - \sum_{j=k_{n+a}+1}^{k_{n+a}-1} \pi'(s+n+a, j) & \text{if } a > 1, k_{n+a} \neq 0. \end{cases} \quad (3.12)$$

where $A(t) = (t+1)\mu/(\lambda+t\mu)$ and $B(t) = \lambda/(\lambda+t\mu)$.

Note that for states (x, y) such that $s+n < x \leq s+m$ and $k_{x-s} \leq y \leq m - (x-s)$, the computations of the $\pi'(x, y)$'s are carried out in a specific order. For each row, where a row corresponds to a value of y (see Figure 3.2), the $\pi'(x, y)$'s are computed in decreasing values of x . Once all the values of a row have been computed, computations for row $y-1$ begin and so on until all $\pi'(x, y)$'s have been computed. It is worth mentioning that $\pi'(s+n+a, k_{n+a})$, for $a = 1, \dots, m-n-2$, is computed from the balance equation of state $(s+n+a-1, k_{n+a-1})$ (given by (3.12)) since at this stage all $\pi'(x, y)$'s involved in the balance equation of state $(s+n+a-1, k_{n+a-1})$ have been computed except for state $(s+n+a, k_{n+a})$.

Given the steady state probabilities, $\pi(x, y)$, we obtain the marginal distribution of the inventory level as follows

$$p_{m-i} = \begin{cases} \sum_{j=k_i}^{m-i} \pi(s+i, j) & \text{if } i = 0, \dots, m, \\ \pi(s+i, m) & \text{if } i = -1, -2, \dots, -\infty, \end{cases} \quad (3.13)$$

where p_i is the probability that the net inventory level is $s+m-i$. The average

total cost $J(s, k)$ can then be written as

$$J(s, \mathbf{k}) = h \sum_{i=0}^{s+m} (s+m-i)p_i + b \sum_{i=s+m+1}^{\infty} (i-s-m)p_i + \lambda c, \quad (3.14)$$

where the three terms represent the expected inventory holding cost, the expected backorder cost, and the expected production cost, respectively. It is not difficult to show that $J(s, \mathbf{k})$ is convex in s . Therefore, the optimal stock level s^* is given by the smallest integer s for which

$$J(s, \mathbf{k}) - J(s+1, \mathbf{k}) \leq 0. \quad (3.15)$$

Here, it is important to note that the distribution specified by $\pi(x, y)$ is independent of the choice of the value s since none of the expressions (3.9)-(3.12) involves s in the computation of the value of the probabilities. This is valuable because it allows for the algorithm to be run only once using a large enough value of s (any $s \geq -m$ will do). The parameter s^* can then be obtained as follows:

$$s^* = \max \left\{ s \geq -m \mid \sum_{i=0}^{s+m-1} p_i \leq \frac{b}{h+b} \right\} - m.$$

To determine the optimal vector \mathbf{k} , we carry out an exhaustive search of all feasible vectors \mathbf{k} . This search is significantly expedited by taking advantage of Property 3 of Theorem 6 and the fact that the decrease of $r^*(x)$ to zero takes place over no more than m steps. Table 3.1 shows the number of feasible \mathbf{k} vectors, using full enumeration (a total of m^{m-1} possible \mathbf{k} vectors) and using Property 3 of Theorem 6. As we can see, using Property 3 dramatically reduces the number of feasible \mathbf{k} vectors. In Table 3.2, we compare the computational performance of our search algorithm to the performance of a standard value iteration algorithm (see for example Puterman (2014, Sec. 8.5.1)) for solving the dynamic program in

(3.3). We do so for an illustrative range of values of m and $\lambda/m\mu$. As we can see, our algorithm can significantly outperform the standard value iteration algorithm. This is particularly the case when the recurring region of the state space is large. This is true when $\lambda/m\mu$ is large, leading to high backorder levels (a high $\lambda/m\mu$ corresponds to either high demand or long lead time). In fact, the computational effort for the value iteration algorithms grows exponentially with $\lambda/m\mu$ while it grows more modestly for our algorithm.

Table 3.1: Number of feasible \mathbf{k} vectors

m	Full enumeration	Enumeration using Property 3 of Theorem 6
5	625	16
10	109	512
15	2.9193×10^{16}	16384
20	5.2429×10^{24}	524288

The above approach can also be adapted to the lost sales case (details are omitted for brevity). Note that in the case of lost sales, the optimal policy requires specifying an additional parameter $\bar{m} \leq m$ which corresponds to the optimal maximum number of units on order. In other words, the optimal policy may never need to place an order of size m . The average total cost, given parameters \bar{m} , s , and \mathbf{k} , can then be expressed as follows

$$J^L(\bar{m}, s, k) = \lambda L p_0 + h \sum_{i=0}^{s+\bar{m}} i p_i + \lambda c(1 - p_0), \quad (3.16)$$

where the three terms in the above expression correspond, respectively, to the expected lost sales cost, the expected inventory holding cost and the expected production cost.

We conclude this section by noting that we observed numerically, and in all the cases tested, that $r^*(x)$ is concave in x (this is the case in both the backorder and

Table 3.2: Computational performance comparisons

m	λ/μ	CPU time (seconds)	
		Value iteration algorithm	Proposed algorithm
5	0.6	0.415	0.002
	0.7	0.71	0.002
	0.75	0.994	0.002
	0.8	1.487	0.002
	0.85	7.29	0.002
	0.9	42.678	0.003
	0.95	633.844	0.003
10	0.6	0.819	0.055
	0.7	1.406	0.065
	0.75	1.96	0.074
	0.8	2.928	0.086
	0.85	14.384	0.088
	0.9	79.017	0.094
	0.95	1194.145	0.111
15	0.6	1.257	0.208
	0.7	2.148	0.245
	0.75	2.993	0.275
	0.8	4.461	0.312
	0.85	21.114	0.321
	0.9	123.098	0.338
	0.95	1953.74	0.389
20	0.6	1.888	7.61
	0.7	2.983	8.791
	0.75	4.16	9.776
	0.8	6.198	11.043
	0.85	30.511	11.332
	0.9	180.525	11.85
	0.95	2886.58	13.477

lost sale cases). That is, $r^*(x+2) - r^*(x+1) \leq r^*(x+1) - r^*(x)$. In terms of the k_l parameters, the concavity of $r^*(x)$ translates into the following constraints:

$$0 \leq k_l - k_{l+1} \leq k_{l+1} - k_{l+2}, \text{ for } l = 0, \dots, m-3, k_l \leq m, \text{ and } k_0 = m. \quad (3.17)$$

These constraints, if they were to be included in our search algorithm, would further reduce the number of feasible \mathbf{k} vectors (see Table 3.3). This would also lead to further improvements in the computational performance of our algorithm (see Table 3.4).

Table 3.3: Number of feasible \mathbf{k} vectors with concave $r^*(x)$

m	Number of feasible \mathbf{k} vectors
5	12
10	97
15	508
20	2087

3.5 Heuristics

In this section, we evaluate the performance of two heuristic policies that are simpler to implement and communicate than the optimal policy. Both policies belong to the same class of policies as the optimal one. Namely, both can be specified in terms of a threshold s on the inventory level and a vector of thresholds $\mathbf{k} = (k_0, \dots, k_{m-1})$ on the inventory on order.

Heuristic H1: Under this heuristic, we set $k_0 = m$ and $k_j = 0$ for $j \neq 0$. In other words, if $x < s$, we order the maximum number of units to bring the number of units on order to its maximum value m ; otherwise, we do not order. Hence,

Table 3.4: Computational performance of the proposed algorithm with concave $r^*(x)$

m	λ/μ	CPU time (seconds)
5	0.6	0.001
	0.7	0.001
	0.75	0.001
	0.8	0.002
	0.85	0.002
	0.9	0.002
	0.95	0.002
10	0.6	0.01
	0.7	0.012
	0.75	0.014
	0.8	0.016
	0.85	0.017
	0.9	0.018
	0.95	0.021
15	0.6	0.065
	0.7	0.076
	0.75	0.085
	0.8	0.097
	0.85	0.1
	0.9	0.105
	0.95	0.121
20	0.6	0.303
	0.7	0.35
	0.75	0.389
	0.8	0.44
	0.85	0.451
	0.9	0.472
	0.95	0.536

starts decreasing from its maximum value m , it continues to decrease by one unit for each unit increase in inventory. Hence, under this heuristic, the inventory position stays constant and equals to $s + m$ if $s \leq x \leq s + m$. This implies that the policy is a modified base-stock policy with base-stock level $s + m$ (ordering in the way such that the inventory position is as close to $s + m$ as possible). As in heuristic H1, the policy here is specified by the single parameter s .

For a given s , we can again follow the approach described in Section 3.4 to obtain the average cost. However, in this case, the analysis simplifies. In particular, the state of the system can be described by the net inventory level. This allows us to characterize, in closed form, the probabilities p_i , where p_i is the probability that the net inventory level is $s + m - i$:

$$p_i = \begin{cases} (\lambda/\mu)^i \frac{p_0}{i!} & \text{for } i = 1, \dots, m, \\ (\lambda/\mu)^m \frac{\rho_m^{i-m}}{m!} p_0 & \text{for } i = m + 1, \dots, \end{cases} \quad (3.19)$$

and

$$p_0 = \left(1 + \sum_{i=1}^m \frac{(\lambda/\mu)^i}{i!} + \frac{(\lambda/\mu)^m \rho_m}{m! (1 - \rho_m)} \right)^{-1}. \quad (3.20)$$

where $\rho_m = \lambda/m\mu$. The expected total cost of the system under this policy is given as follows

$$J^{\text{H2}}(s) = h \sum_{i=0}^{s+m} (s + m - i)p_i + b \sum_{i=s+m+1}^{\infty} (i - s - m)p_i + \lambda c. \quad (3.21)$$

Noting again that the average cost function is convex in s , the optimal value of the threshold s^{H2} can be determined as follows. Let

$$s^+ = \left\lceil \log \left(\left(\frac{1 - \rho_m}{p_0 (\lambda/\mu)^m / m!} \right) \left(\frac{h}{h + b} \right) \right) / \log(\rho_m) \right\rceil,$$

where the notation $\lceil w \rceil$ indicates the smallest integer that is greater than or equal to w . If $s^+ \geq 0$ then $s^{\text{H2}} = s^+$. Otherwise, $s^{\text{H2}} = \max\{s \geq$

$-m|\sum_{i=0}^{s+m-1} p_i \leq b/(h+b)\} - m$. Note that when $m = 1$, $p_0 = 1 - \lambda/\mu$, the optimal base-stock level $s^{\text{H2}} + m$ reduces to $\lceil \log(h/(h+b))/\log(\lambda/\mu) \rceil$. For an integrated production-inventory system, this corresponds to the optimal base-stock level for a system with a single facility (see Buzacott and Shanthikumar for a similar result).

To test the performance of heuristics H1 and H2, we choose a base system with parameters $m = 20$, $\mu = 1.0$, $\lambda = 18$, $h = 2$, $b = 15$, and $c = 0$. We vary parameter values one at a time and obtain the percentage difference between the average cost of the heuristic and that of the optimal policy:

$$\text{Percentage diff.} = \frac{\text{Heuristic policy average cost} - \text{Optimal policy average cost}}{\text{Optimal policy average cost}} \times 100.$$

Representative results are shown in Table 3.5. Note that we set the procurement cost c to zero since it is always incurred in the case of backorders and can be incorporated into the lost sales cost in the case of lost sales. From Table 3.5, we first note that $s^{\text{H1}} \geq s^* \geq s^{\text{H2}}$. This means that Heuristic H1 is associated with the highest maximum attainable inventory level $s^{\text{H1}} + m$, and Heuristic H2 with the lowest level $s^{\text{H2}} + m$. This is because heuristic H1 lacks the ability to adjust the number of units on order in the way heuristic H2 and the optimal policy do. Heuristic H2 must adjust its threshold for the number of units on order one unit at a time and lacks the flexibility of multiple unit increase or decrease that the optimal policy has. These results can also be explained by the shape of the threshold function $r^*(x)$ (illustrated by the vector \mathbf{k}^* as shown in Table 3.5. Note that when the demand rate is low, the optimal vector \mathbf{k}^* tends to have more nonzero values and approaches the corresponding vector for Heuristic H2. On the other hand, for systems with high utilization, \mathbf{k}^* tends to have fewer

Table 3.5: Performance of Heuristics H1 and H2 in the case of backorders

		s^*, \mathbf{k}^*	s^{H1}	s^{H2}	Percentage difference	
					H1	H2
h	2	16, (20, 17, 12, 5, 0, ..., 0)	17	14	0.045	0.991
	5	9, (20, 16, 11, 4, 0, ..., 0)	10	7	0.071	1.512
	7	7, (20, 15, 9, 1, 0, ..., 0)	8	5	0.083	1.837
	10	4, (20,19,14,8,0, ..., 0)	6	3	0.105	2.368
	12	3, (20, 19, 14, 8, 0, ..., 0)	5	2	0.118	2.669
	13	3, (20, 17, 12, 5, 0, ..., 0)	4	1	0.122	2.715
	15	2, (20, 18, 13, 7, 0, ..., 0)	4	0	0.127	3.088
	17	2, (20, 16, 10, 2, 0, ..., 0)	3	0	0.153	3.267
	20	1, (20, 17, 12, 5, 0, ..., 0)	2	-1	0.165	3.697
50	-2, (20, 18, 13, 6, 0, ..., 0)	-1	-4	0.327	6.873	
b	2	2, (20, 18, 13, 7, 0, ..., 0)	4	0	0.127	3.088
	5	8, (20, 15, 9, 1, 0, ..., 0)	9	6	0.077	1.68
	7	10, (20, 17, 12, 5, 0, ..., 0)	11	8	0.063	1.4
	10	13, (20, 16, 10, 2, 0, ..., 0)	14	11	0.056	1.175
	15	16, (20, 17, 12, 5, 0, ..., 0)	17	14	0.045	0.991
	20	18, (20, 19, 14, 8, 0, ..., 0)	20	17	0.04	0.902
	25	20, (20, 19, 14, 8, 0, ..., 0)	22	19	0.038	0.844
	30	22, (20, 17, 12, 5, 0, ..., 0)	23	20	0.035	0.766
	35	23, (20, 19, 14, 8, 0, ..., 0)	25	22	0.034	0.755
	40	25, (20, 15, 9, 1, 0, ..., 0)	26	23	0.032	0.698
	50	27, (20, 15, 9, 1, 0, ..., 0)	28	25	0.03	0.652
	75	30, (20, 19, 14, 8, 0, ..., 0)	32	28	0.027	0.609
100	33, (20, 17, 12, 5, 0, ..., 0)	34	31	0.025	0.542	
λ	4	-7, (20, 19, 17, 15, 13, 11, 9, 6, 3, 0, ..., 0)	-2	-14	96.896	29.265
	6	-5, (20, 19, 17, 15, 12, 9, 6, 2, 0, ..., 0)	-1	-11	43.244	32.535
	8	-3, (20, 18, 15, 12, 8, 4, 0, ..., 0)	-1	-9	17.134	36.441
	10	-2, (20, 18, 15, 11, 7, 2, 0, ..., 0)	0	-6	7.426	33.331
	12	-1, (20, 19, 16, 12, 7, 1, 0, ..., 0)	1	-4	2.844	24.242
	14	1, (20, 19, 15, 11, 5, 0, ..., 0)	3	-1	1.058	12.716
	16	5, (20, 18, 13, 8, 0, ..., 0)	7	3	0.285	4.807
	18	16, (20, 17, 12, 5, 0, ..., 0)	17	14	0.045	0.991
	19	37, (20, 19, 14, 8, 0, ..., 0)	39	36	0.011	0.223

nonzero values and approaches the corresponding vector for Heuristic H1.

As shown in Table 3.5 for the case of backorders, the heuristics perform well except when the holding cost is high, the backorder cost is low, or the demand rate is high. The heuristics lack the flexibility of the optimal policy to adjust the ordering threshold levels. This can lead to higher inventory levels when the heuristics are used, with the associated costs increasing with higher holding costs, lower backorder costs, or lower demand rates. The effect of the holding and backorder costs is more pronounced for heuristic H2 because the heuristic cannot adjust down the ordering thresholds sufficiently quickly. The effect of the low demand rate is more pronounced for heuristic H1 because, under H1, the order up to level cannot be smaller than m . Once delivered, ordered units that are not used to fulfill demand immediately tend to be held in inventory longer leading to higher holding costs.

Heuristics H1 and H2 can easily be adapted to the case of lost sales using the results of Section 3.3. Numerical results comparing the performance of the heuristics to that of the optimal policy are shown in Table 3.6 (in this case, the base system has parameters $m = 20$, $\mu = 1.0$, $\lambda = 19$, $h = 5$, $L = 150$, and $c = 0$). The results can be explained similarly to the case with backorders. Some of the differences in the relative performance of the two heuristics appear to be due the fact that parameters s^{H1} and s^{H2} can no longer be negative as in the backorder case.

Table 3.6: Performance of Heuristics H1 and H2 in the case of lost sales

		s^*, \mathbf{k}^*	s^{H1}	s^{H2}	Percentage difference	
					H1	H2
h	5	17, (20, 15, 10, 2, 0, ..., 0)	17	15	0.0487	1.062
	10	12, (20, 18, 13, 6, 0, ..., 0)	13	10	0.061	1.49
	15	10, (20, 16, 11, 3, 0, ..., 0)	10	8	0.082	1.79
	20	8, (20, 19, 14, 8, 0, ..., 0)	9	7	0.087	2.194
	25	7, (20, 18, 13, 7, 0, ..., 0)	8	5	0.102	2.5
	30	6, (20, 19, 14, 8, 0, ..., 0)	7	5	0.121	2.728
	35	6, (20, 16, 10, 3, 0, ..., 0)	6	4	0.136	2.927
	40	5, (20, 18, 13, 6, 0, ..., 0)	6	3	0.137	3.335
	45	5, (20, 18, 13, 6, 0, ..., 0)	5	3	0.159	3.435
	50	4, (20, 19, 14, 7, 0, ..., 0)	5	3	0.153	4.851
	70	3, (20, 18, 13, 6, 0, ..., 0)	4	1	0.196	4.851
	80	3, (20, 16, 10, 2, 0, ..., 0)	3	1	0.234	5.067
90	2, (20, 19, 14, 8, 0, ..., 0)	3	1	0.243	5.614	
100	2, (20, 17, 12, 5, 0, ..., 0)	3	0	0.26	6.175	
L	25	2, (20, 15, 9, 1, 0, ..., 0)	2	0	0.296	6.68
	50	4, (20, 15, 10, 2, 0, ..., 0)	4	2	0.191	4.149
	75	5, (20, 18, 13, 6, 0, ..., 0)	6	3	0.137	3.335
	100	6, (20, 19, 14, 8, 0, ..., 0)	7	5	0.121	2.728
	150	8, (20, 19, 14, 7, 0, ..., 0)	9	7	0.087	2.194
	175	9, (20, 18, 12, 6, 0, ..., 0)	10	7	0.081	1.966
	200	10, (20, 16, 11, 3, 0, ..., 0)	10	8	0.082	1.79
	250	11, (20, 17, 12, 6, 0, ..., 0)	12	9	0.067	1.618
	300	12, (20, 18, 13, 6, 0, ..., 0)	13	10	0.061	1.49
	350	14, (20, 18, 12, 6, 0, ..., 0)	14	11	0.056	1.317
	400	14, (20, 17, 11, 4, 0, ..., 0)	14	12	0.056	1.274
	500	15, (20, 19, 14, 8, 0, ..., 0)	16	14	0.052	1.181
l	4	0, (11, 8, 5, 2, 0, ..., 0)	1	0	5.538	61.392
	6	0, (15, 12, 9, 5, 1, 0, ..., 0)	1	0	6.223	74.245
	8	0, (18, 15, 12, 9, 5, 0, ..., 0)	1	0	7.909	88.66
	10	1, (20, 16, 12, 8, 4, 0, ..., 0)	2	0	4.358	56.123
	12	2, (20, 17, 13, 8, 3, 0, ..., 0)	3	0	1.95	22.314
	14	3, (20, 19, 15, 10, 5, 0, ..., 0)	4	1	0.858	10.331
	16	5, (20, 19, 15, 10, 0, ..., 0)	6	3	0.314	5.256
	18	8, (20, 19, 14, 7, 0, ..., 0)	9	7	0.087	2.194
	19	10, (20, 18, 13, 6, 0, ..., 0)	11	9	0.043	1.205
	21	16, (20, 16, 10, 0, ..., 0)	16	15	0.008	0.233
	23	25, (20, 18, 11, 1, 0, ..., 0)	25	24	0	0.01
	25	38, (20, 15, 7, 0, ..., 0)	38	37	0.022	0.798
30	64, (20, 0, ..., 0)	64	71	0	0	

Proof. It is easy to show that the optimal cost function is convex in x . That is, the difference $\Delta v^*(x) = v^*(x+1) - v^*(x)$ is non-decreasing in x . In turn, this implies that when $\Delta v^*(x) + c \geq 0$, it is optimal to bring the number of units on order to zero (by cancelling pending orders if necessary). On the other hand, when $\Delta v^*(x) + c < 0$, it is optimal to bring the number of units on order up to m . Thus, in each decision epoch, the optimal policy is a so-called bang-bang policy, where the optimal number of orders to place is either 0 or m . The convexity of $v^*(x)$ also implies that the optimal policy is a base-stock policy with base-stock level s^* , where $s^* = \min\{x \mid \Delta v^*(x) + c \geq 0\}$, such that it is optimal to bring the number of units on order to m if $x < s^*$ and to bring it to 0 otherwise. \square

Noting that once one unit is produced, it is possible to cancel the production of all remaining ones. It is not difficult to see that the dynamics of the system are the same as those where leadtime is exponentially distributed with rate $m\mu$ and only one unit can be on order at any time. In the case of an integrated production-inventory system, this means that the system is equivalent to one with a single facility with a production rate $m\mu$. The dynamics of such a system can be described by a simple Markov chain and various performance measures can be obtained in this case in closed form. In particular, given a base-stock level s , the average cost is given by (we again omit the details for the sake of brevity):

$$J(s) = \lambda c + h \left(s - \rho_m \frac{(1 - \rho_m^s)}{(1 - \rho_m)} \right) + b \frac{\rho_m^{s+1}}{(1 - \rho_m)} s, \quad (3.24)$$

where $\rho_m = \lambda/m\mu$. Noting that the average cost is convex in s , the optimal base-stock level is given by the smallest integer for which $J(s+1) - J(s) \geq 0$. It

is not difficult to show that the base-stock level, s^* , is given by:

$$s^* = \left\lceil \frac{\log\left(\frac{h}{h+b}\right)}{\log(\lambda/m\mu)} \right\rceil. \quad (3.25)$$

Similarly, for systems with lost sales, we can show that the optimal cost function v^* satisfies the following optimality equation:

$$v^*(x) = \begin{cases} g(x) + \lambda v^*(x-1) + m\mu v^*(x) \\ \quad + \min_{0 \leq k \leq m} \{k\mu(v^*(x+1) - v^*(x) + c)\} & \text{if } x > 0, \\ g(x) + \lambda(v^*(x) + L) + m\mu v^*(x) \\ \quad + \min_{0 \leq k \leq m} \{k\mu(v^*(x+1) - v^*(x) + c)\} & \text{otherwise.} \end{cases} \quad (3.26)$$

As in the corresponding backlog case, the optimal policy is bang-bang. It is optimal to bring the number of units on order to m if $x < s^*$ and to zero otherwise. Here too, the dynamics of the system can be modeled using a Markov Chain. Various performance measures can be obtained in closed form. In particular, for a given base-stock level s , the average cost is given by

$$J(s) = \lambda \frac{(1 - \rho_m)\rho_m^s}{(1 - \rho_m^{s+1})} L + \frac{(1 - \rho_m)(s + \rho_m^{s+1}) - \rho_m(1 - \rho_m^{s+1})}{(1 - \rho_m)(1 - \rho_m^{s+1})} h + \lambda c \frac{1 - \rho_m^s}{(1 - \rho_m^{s+1})}, \quad (3.27)$$

where $\rho_m = \lambda/m\mu$. Noting again that the average cost is convex in s , the optimal base-stock level can be easily computed.

We conclude this section by providing numerical results that examine the benefit from order cancellation. To do so, we compare the performance of the original model (a system with backlog and no order cancellation) and its lost sale counterpart to a system that allows for order cancellation. Using a base system with parameters $m = 20$, $\mu = 1.0$, $\lambda = 18$, $h = 2$, $b = 15$, and $c = 0$, Tables 3.7 and 3.8 show the percentage difference in average cost between systems without

Table 3.7: Percentage difference in average cost between systems without and with order cancellation (the backorder case)

h	2	0.881
	5	1.328
	7	1.629
	10	2.093
	12	2.365
	13	2.45
	15	2.783
	17	2.816
	20	3.375
50	7.346	
b	2	2.783
	5	1.478
	7	1.248
	10	0.994
	15	0.881
	20	0.795
	25	0.736
	30	0.68
	35	0.657
	40	0.607
	50	0.562
	75	0.527
	100	0.48
λ	4	141.049
	6	98.15
	8	65.426
	10	37.318
	12	21.346
	14	9.901
	16	4.095
	18	0.881
	19	0.21

Table 3.8: Percentage difference in average cost between systems without and with order cancellation (the lost sale case)

h	5	0.914
	10	1.329
	15	1.552
	20	1.955
	25	2.22
	30	2.498
	35	2.577
	40	2.98
	45	3.078
	50	3.51
	70	4.373
	80	4.705
	90	5.4
100	5.6	
L	25	6.74
	50	3.85
	75	2.98
	100	0.025
	150	0.02
	175	0.018
	200	0.016
	250	0.0144
	300	0.0133
	350	0.0122
	400	0.011
500	0.01	
λ	4	59.662
	6	49.969
	8	35.418
	10	23.105
	12	14.43
	14	8.067
	16	4.356
	18	1.955
	19	1.136
	21	0.221
	23	0.012
25	0	
30	0	

and with order cancellation for the backlog and lost sales cases respectively, where

$$\text{Percentage diff.} = \frac{\text{Average cost without cancellation} - \text{Average cost with cancellation}}{\text{Average cost with cancellation}} \times 100.$$

Results from a more extensive set of experiments reveal similar observations.

As expected, a system in which order cancellations are possible results in lower costs since it has the ability to quickly adjust the number of orders (or, in the case of a production-inventory system, the production capacity). As shown in Table 3.7, for the case of backlogs, the benefit of order cancellation increases as the holding to the backorder cost ratio increases. Without order cancellation, all placed orders eventually show up in inventory. The cost implication of the resulting inventory is higher with higher inventory holding cost or with lower backorder cost. The benefit of order cancellation increases with decreases in the demand rate. This is because, when the demand rate is low, any inventory held tends to be held for longer periods of time. Systems with order cancellation can mitigate the need for holding inventory by placing multiple orders when demand arises, thereby expediting deliveries, but then cancelling pending orders once demand is satisfied. This ability to expedite deliveries without repercussion on inventory holding cost is not available to the system without order cancellations. Such systems end up carrying more inventory on average than systems without order cancellation. Table 3.8 tells a similar story for systems with lost sales.

Chapter 4

Optimal Policies for Inventory Systems with Concave Ordering Costs

4.1 Introduction

Most of the literature on inventory systems usually assumes a linear ordering cost or a linear ordering cost with a setup cost. As Scarf (1963) argues, “This type of cost functions has appeared in inventory theory not necessarily because of its realism, but because it provides one of the few examples of cost functions with a decreasing average cost for which the analysis of inventory policies is relatively easy.” In this paper, we consider inventory systems with general concave ordering cost functions, where the type of ordering costs described in Scarf (1959) is a special case under our setting. The class of concave ordering cost functions is a

special type of a decreasing average cost and there are many examples of concave ordering costs in practice. Consider the following examples.

Quantity Discounts: Quantity discounts provide a practical foundation for coordinating inventory decisions in supply chains. Sellers usually employ quantity discount schemes or contracts to give buyers the incentive to buy more. That is, the larger the order is, the lower the marginal price will be. Ordering costs with quantity discounts can usually be expressed by piecewise linear concave functions: first, there is a setup cost; then the first few items have the same per-unit cost; the next few items have a lower per-unit cost, and so on.

The Effect of Economies of Scale: In economics, one of the common assumptions on production functions is that they have the economies of scale feature. Basically, the more a firm produces the same item, the more efficient the production technology will be. The transportation costs in supply chains also exhibit economies of scale: the more volume of goods to be shipped, the cheaper the marginal cost will be. Concave functions are an important class of functions exhibiting the feature of economies of scale.

Procurement with Multiple Suppliers: In many cases in practice, there are multiple suppliers available for a buyer. Usually local suppliers offer relatively lower setup costs but with higher per-unit costs and overseas or distant suppliers offer higher setup costs but with lower per-unit costs. Hence, if the buyer chooses the suppliers optimally, the resulting ordering cost is a piecewise linear concave function. A related case of a buyer that purchases from both long-term suppliers and spot markets is treated in Yi and Scheller-Wolf (2003) and the references therein. There are many other examples of concave costs due to the availability of

multiple choices of labor and production, see Fox et al. (2006) for examples and references therein.

Scarf (1959) proves that the (s, S) policy is optimal for an inventory system with a fixed ordering cost and a unit ordering cost and does it by introducing the notion of K -convexity. This type of ordering cost is a special case of concave ordering costs. Karlin (1958) analyzes the optimal ordering policy for a one-period inventory problem with concave ordering costs. Scarf (1963) points out that it is difficult to generalize the result to the dynamic multiperiod setting. There has been only limited research on stochastic inventory systems with concave ordering costs. Porteus (1971) analyzes inventory systems with piecewise linear concave ordering costs. He shows that a generalized (s, S) policy is optimal for a multi-period periodic review inventory system under some mild assumption on cost functions and that demand has a one-sided Polya density. He does it by introducing a generalized notion of K -convexity called quasi- K -convexity. However, the class of one-sided Polya densities does not include many densities encountered in practice, for example, the normal distribution, beta distribution and most gamma distributions, although it does include the exponential distribution and all its finite convolutions. Porteus (1972) also shows that the generalized (s, S) policy is optimal for uniform demand distributions.

Fox et al. (2006) consider the optimal policy for an inventory system with two suppliers: the buyer incurs a high variable cost but negligible fixed cost for the first supplier (HVC) and a lower variable cost but a substantial fixed cost for the second supplier (LVC). The resulting ordering cost is a two-piece linear concave function.

They show that the optimal policy is a (s, S_{HVC}, S_{LVC}) policy, which is a special case of the generalized (s, S) policy, under the condition that the demand density is log-concave. Their proof relies on K -convexity and quasi-convex properties since they consider a two-piece linear concave function. Although the class of log-concave densities is less restrictive than the class of one-sided Polya densities, it still only covers a limited range of distributions. Furthermore, their results do not cover general piecewise linear concave ordering costs. Hence, whether or not the generalized (s, S) policy is optimal for general demand distributions remained an open question.

Recently, Chen et al. (2010) consider joint pricing and inventory control for inventory systems with concave ordering costs. They utilize quasi- K -convexity to show that the optimal policy is a generalized (s, S, p) policy when demand distributions are Polya or uniform. Another related paper is Yi and Scheller-Wolf (2003), where they also consider a two-supplier inventory problem: the buyer has a long-term contract from a regular supplier with a minimum and maximum purchasing quantity, and the buyer can also purchase from a spot market that has no quantity limitation but with a fixed entry fee. They partially characterize the structure of the optimal policy and their proof relies on a closure property of K -convexity. Note that the ordering cost in their case is no longer concave since they assume a limited capacity for the regular supplier and the corresponding optimal policy is not a generalized (s, S) policy.

Chen and Simchi-Levi (2004) consider the joint inventory-pricing control problem with fixed ordering costs. They introduce the concept of sym- K -convexity, which is a generalization of K -convexity, and show that the

optimal policy can be fully characterized except for a bounded interval for the multiplicative demand model. Chao and Zipkin (2008) study a model with a fixed cost function that is neither convex nor concave: the fixed cost is incurred only if the order quantity exceeds a threshold, and hence the cost function can be written as $c(x) = K\delta(x - C)$ for some constant C . They apply the property of K -convexity and partially characterize the optimal policy with three critical points which divide the state space into five regions.

In contrast to the existing literature, we characterize the structure of optimal policies for inventory systems with concave ordering costs with general demand distributions. In order to analyze the structure of the optimal policy, we first introduce a monotone condition that ensures the optimality of a generalized (s, S) policy. We then introduce the concept of c -convexity, a generalization of K -convexity, and use it to show that the value function for this problem is c -convex with respect to a modified ordering cost function. Based on the c -convexity of the value function, we show that, except for a bounded region of the state space, the generalized (s, S) policy is optimal. We also provide conditions under which the generalized (s, S) policy is optimal for all regions of the state space. Our results can be readily extended to systems with time-varying cost parameters, systems with fixed leadtimes and to systems with lost sales. The notion of c -convexity we introduce in this paper may also have usefulness to other inventory control problems.

4.2 Inventory Systems with Concave Ordering Costs

We consider a single product single stage inventory problem with multiple periods, stochastic demands, and zero leadtime. The assumption on zero leadtime is not critical and is made for ease of exposition (see Section 4.6 for extensions). Demand ξ_t in each period t is a continuous random variable with $E[\xi_t] < \infty$ and distribution function $F_t(x), x \geq 0$, where $t = 1, \dots, T$ and T corresponds to the length of the planning horizon. Demands in different periods are independent but not necessarily identically distributed (i.e., demand can be time-varying). Inventory is replenished from an outside supplier immediately (i.e., with zero leadtime) with ample stock. Demand is satisfied from on-hand inventory, if any is available; otherwise it is backordered. In each period, the inventory manager must decide on the quantity to order to minimize the expected discounted cost over the entire planning horizon. There are three types of costs in each period t : (1) an ordering cost $c(z)$ if the order quantity is $z, z \geq 0$, (2) a holding cost $h_t(x^+)$ and (3) a backordering cost $b_t(x^-)$ given the inventory level x in period t , where $x^+ = \max\{0, x\}$ and $x^- = \max\{0, -x\}$. Finally, we allow a discount factor $\alpha \in (0, 1]$.

In order to simplify our presentation, we first consider a piecewise linear concave ordering cost $c(\cdot)$ with n linear pieces. Specifically, we can express

$$c(x) = \min_{i=1, \dots, n} \{K_i \delta(x) + c_i x\},$$

with $0 \leq K_1 < K_2 < \dots < K_n$ and $c_1 > c_2 > \dots > c_n \geq 0$, where δ is defined as

follows

$$\delta(x) = \begin{cases} 1 & \text{if } x > 0, \\ 0 & \text{if } x \leq 0. \end{cases} \quad (4.1)$$

We will discuss how we can deal with general concave ordering costs at the end of Section 4.3 and time-varying costs in the section on extensions (see Section 4.6).

Let

$$g_t(x - \xi_t) = \begin{cases} h_t(x - \xi_t) & \text{if } x \geq \xi_t, \\ b_t(\xi_t - x) & \text{otherwise.} \end{cases}$$

Let x_t be the starting inventory level and y_t be the post-ordering inventory level for period t , with $x_{t+1} = y_t - \xi_t$. Given $x_1, \dots, x_T, y_1, \dots, y_T$, i.e., the ordering quantities being $q_t = y_t - x_t, t = 1, \dots, T$, the expected discounted total cost is given by

$$E \left\{ \sum_{t=1}^T \alpha^t [c(y_t - x_t) + g_t(y_t - \xi_t)] \right\}. \quad (4.2)$$

Let $v_t^*(x)$ be the value function (the optimal expected discounted cost) in period t when the inventory level in period t is x . Then the corresponding dynamic programming formulation is given by

$$v_t^*(x) = \min_{y \geq x} \{c(y - x) + E g_t(y - \xi_t) + \alpha E v_{t+1}^*(y - \xi_t)\}. \quad (4.3)$$

Finally we let $v_{T+1}^*(x) = 0$ for all x .

Assumption 1. *We assume that $L_t(y) = E g_t(y - \xi_t)$ is convex in y and finite for any y .*

For example, this assumption is satisfied if h_t and b_t are linear and $E[\xi_t] < \infty$. The finiteness of L_t ensures that L_t is continuous on $(-\infty, \infty)$ (by the dominated convergence theorem).

Let

$$H_t(y) = Eg_t(y - \xi_t) + \alpha Ev_{t+1}^*(y - \xi_t).$$

Then the optimality equation is given by

$$v_t^*(x) = \min_{y \geq x} [c(y - x) + H_t(y)]. \quad (4.4)$$

Given x , let $y_t(x)$ be the smallest minimizer of $c(y - x) + H_t(y)$, i.e.,

$$y_t(x) = \min \arg \min_{y \geq x} \{c(y - x) + H_t(y)\}. \quad (4.5)$$

Hence, given the current inventory level is x , it is optimal to order $y_t(x) - x$ quantity in period t .

4.3 The Structure of the Optimal Policy

In this section, we show that, except for a bounded region, the optimal policy can be described by a generalized (s, S) policy.

First, we show a conditional monotone property for $y_t(x)$ for any concave function c .

Theorem 9. *Suppose that $y_t(x) > x$, then $y_t(z) \leq y_t(x)$ for $z \in (x, y_t(x))$.*

Proof. Suppose that for some x , we have $y_t(x) - x > 0$. Let $y_t(x) > z > x$. Since we know that $c(x)$ is concave in x , for $\omega > 0$ we have

$$c(y_t(x) + \omega - z) - c(y_t(x) - z) \geq c(y_t(x) + \omega - x) - c(y_t(x) - x),$$

which implies that

$$\begin{aligned} & c(y_t(x) + \omega - z) + H_t(y_t(x) + \omega) - [c(y_t(x) - z) + H_t(y_t(x))] \\ & \geq c(y_t(x) + \omega - x) + H_t(y_t(x) + \omega) - [c(y_t(x) - x) + H_t(y_t(x))] \geq 0. \end{aligned} \quad (4.6)$$

Since inequality (4.6) is true for all $\omega > 0$ and c is continuous in $(0, \infty)$, it follows that $y_t(z) \leq y_t(x)$. \square

As far as we know, this is a new result in the literature. The interesting aspect of this result is that the conditional monotone property in period t holds for general concave ordering costs. However we do not know what will happen for $z \notin (x, y_t(x))$. Next, we describe a monotone condition that is the key to characterizing the structure of optimal policies.

Condition 1. $y_t(x_2) > x_2$ implies that $y_t(x_1) > x_1$ for any $x_1 < x_2$. In words, if it is optimal to order a positive amount when the starting inventory level is x_2 , then it must be optimal to order a positive amount when the starting inventory level is less than x_2 in period t .

It turns out that if we know that $y_t(x_2) > x_2$ implies that $y_t(x_1) > x_1$ for any $x_1 < x_2$, then coupled with Theorem 9, we can show that $y_t(x_1) \geq y_t(x_2)$ for all $x_1 < x_2$ such that $y_t(x_2) > x_2$.

Lemma 10. *Under Condition 1, we have*

- (1) $y_t(x_1) \geq y_t(x_2)$ for all $y_t(x_2) > x_2$ and $x_1 < x_2$.
- (2) There exists some x_0 such that $y_t(x) = x$ for all $x \geq x_0$ and $y_t(x)$ is non-increasing in x for $x \in (-\infty, x_0]$.

Proof. We prove the first part by contradiction. Suppose we have $y_t(x_1) < y_t(x_2)$ given that $y_t(x_2) > x_2$, $y_t(x_1) > x_1$ and $x_1 < x_2$. We differentiate two cases. (1) $y_t(x_1) \leq x_2$. This case is impossible, since under Condition 1, we must have $y_t(y_t(x_1)) > y_t(x_1)$, i.e., $y_t(x_1)$ is not an optimal order-up-to level for x_1 , which violates the optimality of $y_t(\cdot)$. (2) $y_t(x_1) \in (x_2, y_t(x_2))$. This case is impossible since it violates Theorem 9. We know that by Theorem 9, we must have $y_t(x_2) < y_t(x_1)$ since $x_2 \in (x_1, y_t(x_1))$.

Since we have $\lim_{x \rightarrow \infty} H_t(x) = \infty$, it follows that for sufficiently large x , we must have $y_t(x) = x$. Let x_0 the smallest value such that $y_t(x) = x$. Then $y_t(x) > x$ for all $x < x_0$ by Condition 1. It follows that $y_t(x)$ is non-increasing in x in that domain by part (1) of this lemma. It can also be shown that $y_t(x) = x$ for all $x > x_0$, otherwise if $y_t(x) > x > x_0$ then we must have $y_t(x_0) > x_0$. This contradicts the definition of x_0 . \square

Theorem 10. *If Condition 1 is satisfied, then the optimal inventory policy in period t is a generalized (s, S) policy, i.e., there exists $(s_{m,t}, \dots, s_{1,t}, S_{1,t}, \dots, S_{m,t})$ with $s_{m,t} < s_{m-1,t} < \dots < s_{1,t} \leq S_{1,t} < S_{2,t} < \dots < S_{m,t}$ for some $m \leq n$ such that if $x < s_{m,t}$ then we order up to $S_{m,t}$ and if $x \in [s_{i,t}, s_{i-1,t})$ then we order up to $S_{i-1,t}$ for $i = 2, \dots, m$, and finally we order nothing for $x \geq s_{1,t}$. Hence, we have at most n distinctive such order-up-to levels $S_{i,t}$.*

Proof. Let $s_{1,t} = \min\{x : H_t(x) \leq c(y-x) + H_t(y), y > x\}$, i.e., $s_{1,t}$ is the minimum starting inventory level such that it is optimal to order nothing (the existence of $s_{1,t}$ is due to $\lim_{x \rightarrow \infty} H_t(x) = \infty$ and H_t is continuous). It follows that if $x > s_{1,t}$, then it is also optimal to order nothing, since otherwise it would violate Condition 1. Also if $x < s_{1,t}$, then it must be optimal to order a positive quantity and the

post-ordering inventory level must be greater than or equal to $s_{1,t}$, since otherwise it would violate the definition of $s_{1,t}$.

Let

$$\hat{S}_{i,t} = \min \arg \min_{y \geq s_{1,t}} \{H_t(y) + c_i y\},$$

i.e., $\hat{S}_{i,t}$ is the minimum of $H_t(y) + c_i y$ on $[s_{1,t}, \infty)$ (the existence of $\hat{S}_{i,t}$ is due to the continuity of H_t). Since $c_1 > c_2 > \dots > c_n$, it follows that $\hat{S}_{1,t} \leq \hat{S}_{2,t} \leq \dots \leq \hat{S}_{n,t}$. For $x \in (-\infty, s_{1,t})$, let

$$\begin{aligned} v_{i,t}(x) &= \min_{y_i \geq x} [c_i y_i + K_i \delta(y_i - x) + H_t(y_i)] - c_i x \\ &= \min \{ \min_{y_i > x} [c_i y_i + K_i + H_t(y_i)], c_i x + H_t(x) \} - c_i x. \end{aligned}$$

We have

$$\min_{y \geq x} \{c(y - x) + H_t(y)\} = \min_{i=1, \dots, n} \{v_{i,t}(x)\}.$$

It follows that for any starting inventory level $x \in (-\infty, s_{1,t})$ (note that it is optimal to order a positive quantity for such starting inventory level x), it must be optimal to order to one of the levels in $\{\hat{S}_{1,t}, \hat{S}_{2,t}, \dots, \hat{S}_{n-1,t}\}$. Ties can be broken by choosing the smallest solution.

Suppose that for small enough δ it is optimal to order up to $\hat{S}_{i_1,t}$ for some $i_1 \in \{1, \dots, n\}$ for starting inventory level $x \in [s_{1,t} - \delta, s_{1,t})$. If $i_1 = n$, then we are done. Otherwise, let $s_{2,t}$ be the smallest value such that it is optimal to order up to $\hat{S}_{i_1,t}$, i.e., $y_t(x) = \hat{S}_{i_1,t}$ for $x \in [s_{2,t}, s_{1,t})$. We define $S_{1,t} \equiv \hat{S}_{i_1,t}$. Since it is also optimal to order a positive quantity for $x < s_{2,t}$, by Lemma 10, we have $y_t(x) > \hat{S}_{i_1,t}$ for $x < s_{2,t}$. Again, suppose for some small enough δ it is optimal to order-up-to $\hat{S}_{i_2,t}$ for some $i_2 \in \{1, \dots, n\}$ for $x \in [s_{2,t} - \delta, s_{1,t})$. Obviously, we have

$i_2 > i_1$ by the conditional monotone property. We define $S_{2,t} \equiv \hat{S}_{i_2,t}$. If $S_{2,t} = \hat{S}_{n,t}$, then we are done. Otherwise by a similar argument, we can iteratively define $s_{i,t}$ and $S_{i,t}$ ($i > 3$) such that it is optimal to order-up-to $S_{i,t}$ for $x \in [s_{i+1,t}, s_i)$ until we have some $s_{m,t}$ and $S_{m,t} = \hat{S}_{n,t}$. Then it follows that if $x < s_{m,t}$, it is optimal to order up to $S_{m,t} = \hat{S}_{n,t}$. It is also clear that $m \leq n$. \square

To analyze the structure of the optimal policy, we can rewrite the ordering cost as follows: $c(x) = K_i + c_i x$ for $x \in [z_{i-1}, z_i], i = 2, \dots, n-1$, $c(x) = K_1 \delta(x) + c_1 x$ for $x \in [0, z_1]$ and $c(x) = K_n + c_n x$ for $x \geq z_{n-1}$, where $0 < z_1 < \dots < z_{n-1}$. Note that $c(x) - c_n x \geq 0$ for all $x \geq 0$. We first use the following transformation. We define $\bar{c}(x) \equiv c(x) - c_n x$. It is clear that $\bar{c}(x) = K_n$ for $x \geq z_{n-1}$. Then we can reformulate the dynamic recursion in (4.3) as follows.

$$\bar{v}_t^*(x) = \min_{y \geq x} E\{\bar{c}(y-x) + (1-\alpha)c_n[y-\xi_t] + c_n \xi_t + g_t(y-\xi_t) + \alpha E \bar{v}_{t+1}^*(y-\xi_t)\}, \quad (4.7)$$

with $\bar{v}_{T+1}(x) = c_n x$. Let

$$\bar{G}_t(y) = (1-\alpha)c_n[y - E\xi_t] + c_n E\xi_t + E g_t(y - \xi_t) + \alpha E \bar{v}_{t+1}^*(y - \xi_t),$$

and

$$\bar{L}_t(y) = (1-\alpha)c_n[y - E(\xi_t)] + c_n E\xi_t + E g_t(y - \xi_t).$$

Then we have

$$\bar{v}_t^*(x) = \min_{y \geq x} \{\bar{c}(y-x) + \bar{G}_t(y)\}. \quad (4.8)$$

We can show that $v_t^*(x) = \bar{v}_t^*(x) - c_n x$.

Next, we introduce a new generalized convexity notion to which we refer as c -convexity.

Definition 4. A function f is said to be c -convex for any nonnegative nondecreasing concave function c if for any $x_1 < x_2$ and $\theta \in (0, 1)$ the following inequality holds

$$\theta f(x_1) + (1 - \theta)[f(x_2) + c(x_2 - \theta x_1 - (1 - \theta)x_2)] \geq f(\theta x_1 + (1 - \theta)x_2).$$

One can view c -convexity as a generalization of K -convexity.

Based on the definition of c -convexity, we can show that the following lemma holds.

Lemma 11. c -convex functions have the following properties. Assuming $c^i, i = 1, 2$ are nonnegative nondecreasing concave functions.

1. Convexity is equivalent to 0-convexity, where 0 denotes that $c(x) \equiv 0$ for all $x \geq 0$.
2. If g is c -convex, then $g(x + a)$ is also c -convex for any a .
3. If g is c^1 -convex, then it is also c^2 -convex if $c^2(x) \geq c^1(x)$ for all $x \geq 0$.
4. If g_i is c^i -convex for $i = 1, 2$, then $a_1 g_1 + a_2 g_2$ is $a_1 c^1 + a_2 c^2$ -convex for all non-negative a_i .
5. If g is c -convex and $f(x) = E[g(x - \xi)] < \infty$, where ξ is a random variable, then $f(x)$ is also c -convex.

The proof of the lemma is straightforward and hence omitted.

Next, we show that the value function is indeed \bar{c} -convex where $\bar{c}(x) = c(x) - c_n x$.

Lemma 12. $\bar{v}_t^*(x)$ is \bar{c} -convex for all t .

Proof. We use induction to show that $\bar{v}_t^*(x)$ is \bar{c} -convex for all t . It is clear that \bar{v}_T^* is \bar{c} -convex since it is a linear function. Suppose that \bar{v}_{t+1} is \bar{c} -convex. Since \bar{L}_t is convex, then $\bar{G}_t(x) = \bar{L}_t(x) + \alpha E\bar{v}_{t+1}^*(x - \xi_t)$ must be \bar{c} -convex according to Lemma 11. For any $x_1 < x_2$ and $\theta \in (0, 1)$, we differentiate two possible cases: (1) $y_t(x_1) \geq \theta x_1 + (1 - \theta)x_2$ and (2) $y_t(x_1) < \theta x_1 + (1 - \theta)x_2$ (recall that $y_t(x)$ is the optimal order-up-to level for x).

For case (1), we have

$$\begin{aligned}
& \theta \bar{v}_t^*(x_1) + (1 - \theta)[\bar{v}_t^*(x_2) + \bar{c}(x_2 - \theta x_1 - (1 - \theta)x_2)] \\
&= \theta[\bar{G}_t(y_t(x_1)) + \bar{c}(y_t(x_1) - x_1)] + (1 - \theta)[\bar{G}_t(y_t(x_2)) + \bar{c}(y_t(x_2) - x_2) \\
&\quad + \bar{c}(x_2 - \theta x_1 - (1 - \theta)x_2)] \\
&\geq \theta[\bar{G}_t(y_t(x_1)) + \bar{c}(y_t(x_1) - x_1)] + (1 - \theta)[\bar{G}_t(y_t(x_2)) + \bar{c}(y_t(x_2) - \theta x_1 - (1 - \theta)x_2)] \\
&\geq \theta[\bar{G}_t(y_t(x_1)) + \bar{c}(y_t(x_1) - \theta x_1 - (1 - \theta)x_2)] + (1 - \theta)[\bar{G}_t(y_t(x_2)) \\
&\quad + \bar{c}(y_t(x_2) - \theta x_1 - (1 - \theta)x_2)] \\
&\geq \bar{v}_t^*(\theta x_1 + (1 - \theta)x_2).
\end{aligned}$$

The first inequality is due to the subadditivity of \bar{c} since \bar{c} is concave. The second inequality is due to the fact that \bar{c} is nondecreasing and $y_t(x_1) \geq \theta x_1 + (1 - \theta)x_2$.

For case (2), we differentiate two subcases: (2a) $\bar{G}_t(y_t(x_1)) + \bar{c}(y_t(x_1) - x_1) > \bar{G}_t(y_t(x_2)) + \bar{c}(y_t(x_2) - x_2) + \bar{c}(x_2 - \theta x_1 - (1 - \theta)x_2)$ and (2b) $\bar{G}_t(y_t(x_1)) + \bar{c}(y_t(x_1) - x_1) \leq \bar{G}_t(y_t(x_2)) + \bar{c}(y_t(x_2) - x_2) + \bar{c}(x_2 - \theta x_1 - (1 - \theta)x_2)$.

First, we consider subcase (2a). In this subcase, we have

$$\begin{aligned}
& \theta \bar{v}_t^*(x_1) + (1 - \theta)[\bar{v}_t^*(x_2) + \bar{c}(x_2 - \theta x_1 - (1 - \theta)x_2)] \\
&= \theta[\bar{G}_t(y_t(x_1)) + \bar{c}(y_t(x_1) - x_1)] + (1 - \theta)[\bar{G}_t(y_t(x_2)) + \bar{c}(y_t(x_2) - x_2) \\
&\quad + \bar{c}(x_2 - \theta x_1 - (1 - \theta)x_2)]
\end{aligned}$$

$$\begin{aligned}
&> \theta[\bar{G}_t(y_t(x_2)) + \bar{c}(y_t(x_2) - x_2) + \bar{c}(x_2 - \theta x_1 - (1 - \theta)x_2)] \\
&\quad + (1 - \theta)[\bar{G}_t(y_t(x_2)) + \bar{c}(y_t(x_2) - \theta x_1 - (1 - \theta)x_2)] \\
&\geq \bar{G}_t(y_t(x_2)) + \bar{c}(y_t(x_2) - \theta x_1 - (1 - \theta)x_2) \\
&\geq \bar{v}_t^*(\theta x_1 + (1 - \theta)x_2).
\end{aligned}$$

The first inequality is due to the subadditivity of \bar{c} and the assumption of subcase (2a) and the second inequality is due to the subadditivity of \bar{c} .

Next, we consider subcase (2b). Since $y_t(x_1) < \theta x_1 + (1 - \theta)x_2$ and $\bar{G}_t(y_t(x_1)) + \bar{c}(y_t(x_1) - x_1) \leq \bar{G}_t(x_1)$, it follows that there exists some \hat{x}_1 such that $x_1 \leq \hat{x}_1 \leq y_t(x_1)$ and $\bar{G}_t(\hat{x}_1) = \bar{G}_t(y_t(x_1)) + \bar{c}(y_t(x_1) - x_1)$ since \bar{G}_t is continuous. Then there exists $1 > \rho \geq \theta$ such that $\rho \hat{x}_1 + (1 - \rho)y_t(x_2) = \theta x_1 + (1 - \theta)x_2$. In this subcase we have

$$\begin{aligned}
&\theta \bar{v}_t^*(x_1) + (1 - \theta)[\bar{v}_t^*(x_2) + \bar{c}(x_2 - \theta x_1 - (1 - \theta)x_2)] \\
&= \theta[\bar{G}_t(y_t(x_1)) + \bar{c}(y_t(x_1) - x_1)] + (1 - \theta)[\bar{G}_t(y_t(x_2)) + \bar{c}(y_t(x_2) - x_2) \\
&\quad + \bar{c}(x_2 - \theta x_1 - (1 - \theta)x_2)] \\
&= \theta[\bar{G}_t(\hat{x}_1)] + (1 - \theta)[\bar{G}_t(y_t(x_2)) + \bar{c}(y_t(x_2) - x_2) + \bar{c}(x_2 - \theta x_1 - (1 - \theta)x_2)] \\
&\geq \rho[\bar{G}_t(\hat{x}_1)] + (1 - \rho)[\bar{G}_t(y_t(x_2)) + \bar{c}(y_t(x_2) - x_2) + \bar{c}(x_2 - \theta x_1 - (1 - \theta)x_2)] \\
&\geq \rho[\bar{G}_t(\hat{x}_1)] + (1 - \rho)[\bar{G}_t(y_t(x_2)) + \bar{c}(y_t(x_2) - \rho \hat{x}_1 - (1 - \rho)y_t(x_2))] \\
&\geq \bar{G}_t(\rho \hat{x}_1 + (1 - \rho)y_t(x_2)) \\
&= \bar{G}_t(\theta x_1 + (1 - \theta)x_2) \\
&\geq \bar{v}_t^*(\theta x_1 + (1 - \theta)x_2),
\end{aligned}$$

where the first inequality is due to $\rho \geq \theta$ and $\bar{G}_t(y_t(x_1)) + \bar{c}(y_t(x_1) - x_1) \leq \bar{G}_t(y_t(x_2)) + \bar{c}(y_t(x_2) - x_2) + \bar{c}(x_2 - \theta x_1 - (1 - \theta)x_2)$. The second inequality is due

to the subadditivity of \bar{c} . The third inequality is due to \bar{G}_t being \bar{c} -convex. This completes the inductive proof. \square

Remark 1. *In contrast to the proof of K -convexity in Scarf (1959), our proof of the \bar{c} -convexity of the value function does not rely on any structural properties of the optimal policy.*

Based on the \bar{c} -convexity of the value function, we can characterize the structure of the optimal policy as follows. Define

$$\hat{S}_{i,t} = \min \arg \min_y [\bar{G}_t(y) + (c_i - c_n)y].$$

Let

$$\hat{s}_{n,t} = \max\{x | \bar{G}_t(x) > \bar{G}_t(\hat{S}_{n,t}) + K_n, \hat{S}_{n,t} - x \geq z_{n-1}\},$$

i.e., $\hat{s}_{n,t}$ is the largest value such that ordering up to $\hat{S}_{n,t}$ is preferable to not ordering (there always exists such $\hat{s}_{n,t}$ since \bar{G}_t is \bar{c} -convex and $L_t(x) \rightarrow \infty$ as $x \rightarrow -\infty$). Also let $\underline{s}_{0,t}$ be the maximum z such that it is optimal to order a positive quantity for all $x \in [\hat{s}_{n,t}, z]$ in period t .

Theorem 11. *The optimal policy has the following properties.*

- (1) *The generalized (s, S) policy is optimal for $x < \underline{s}_{0,t}$.*
- (2) *It is optimal not to order for $x \geq \hat{S}_{n,t}$.*
- (3) *If it is optimal to order for $x \in (\underline{s}_{0,t}, \hat{S}_{n,t})$, then its optimal order-up-to level is less than $\hat{S}_{n,t}$.*

Thus, except for the interval $(\underline{s}_{0,t}, \hat{S}_{n,t})$, the generalized (s, S) policy is optimal.

Proof. (1) First, we show that if $\bar{G}_t(x) > \bar{G}_t(\hat{S}_{n,t}) + \bar{c}(\hat{S}_{n,t} - x)$ and $\hat{S}_{n,t} - x \geq z_{n-1}$, then we must have $\bar{G}_t(x - \delta) > \bar{G}_t(\hat{S}_{n,t}) + \bar{c}(\hat{S}_{n,t} - x + \delta)$ for all $\delta > 0$. We show this by contradiction. Suppose that $\bar{G}_t(x_0) > \bar{G}_t(\hat{S}_{n,t}) + \bar{c}(\hat{S}_{n,t} - x_0)$ with $\hat{S}_{n,t} - x_0 \in (z_{n-1}, \infty)$ but $\bar{G}_t(x_0 - \delta) \leq \bar{G}_t(\hat{S}_{n,t}) + \bar{c}(\hat{S}_{n,t} - x_0 + \delta)$ for some $\delta > 0$. There exists some $\rho \in (0, 1)$ such that $\rho(x_0 - \delta) + (1 - \rho)\hat{S}_{n,t} = x_0$. Note that

$$\bar{c}(\hat{S}_{n,t} - x_0) = \rho\bar{c}(\hat{S}_{n,t} - x_0 + \delta) + (1 - \rho)\bar{c}(\hat{S}_{n,t} - x_0),$$

since $\bar{c}(\hat{S}_{n,t} - x_0 + \delta) = \bar{c}(\hat{S}_{n,t} - x_0) = K_n$ for $\hat{S}_{n,t} - x_0 \geq z_{n-1}$. Thus, we have

$$\begin{aligned} & \bar{G}_t(\hat{S}_{n,t}) + \bar{c}(\hat{S}_{n,t} - x_0) \\ &= \rho[\bar{G}_t(\hat{S}_{n,t}) + \bar{c}(\hat{S}_{n,t} - x_0 + \delta)] + (1 - \rho)[\bar{G}_t(\hat{S}_{n,t}) + \bar{c}(\hat{S}_{n,t} - x_0)] \\ &\geq \rho[\bar{G}_t(x_0 - \delta)] + (1 - \rho)[\bar{G}_t(\hat{S}_{n,t}) + \bar{c}(\hat{S}_{n,t} - x_0)] \\ &\geq \bar{G}_t(x_0). \end{aligned}$$

This implies that $\bar{G}_t(\hat{S}_{n,t}) + \bar{c}(\hat{S}_{n,t} - x_0) \geq \bar{G}_t(x_0)$, which contradicts the fact that $\bar{G}_t(\hat{S}_{n,t}) + \bar{c}(\hat{S}_{n,t} - x_0) < \bar{G}_t(x_0)$.

Since

$$\hat{s}_{n,t} = \max\{x \mid \bar{G}_t(x) > \bar{G}_t(\hat{S}_{n,t}) + K_n, \hat{S}_{n,t} - x \geq z_{n-1}\},$$

it follows that it is optimal to order a positive quantity for $x < \hat{s}_{n,t}$ based on the above result. By the definition of $\underline{s}_{0,t}$, it is optimal to order a positive quantity for all $x \in [\hat{s}_{n,t}, \underline{s}_{0,t})$ in period t . It follows that the generalized (s, S) policy is optimal on $(-\infty, \underline{s}_{0,t})$ based on Theorem 10.

(2) We show this result by contradiction. Suppose that $y_t(x_0) > x_0$ for some $x_0 > \hat{S}_{n,t}$. Then we must have

$$\bar{G}_t(x_0) > \bar{G}_t(y_t(x_0)) + \bar{c}(y_t(x_0) - x_0).$$

But we know that for any $x < x_0$ and $\theta \in (0, 1)$ such that $\theta x + (1 - \theta)y_t(x_0) = x_0$, we have

$$\theta \bar{G}_t(x) + (1 - \theta)[\bar{G}_t(y_t(x_0)) + \bar{c}(y_t(x_0) - x_0)] \geq \bar{G}_t(x_0),$$

since \bar{G}_t is \bar{c} -convex. It follows that $\bar{G}_t(x) > \bar{G}_t(x_0)$ and hence $\bar{G}_t(\hat{S}_{n,t}) > \bar{G}_t(x_0)$, which contradicts the fact that $\hat{S}_{n,t}$ minimizes $\bar{G}_t(x)$.

(3) is due to Theorem 9. □

Remark 2. *In general, the generalized (s, S) policy may not be optimal over the interval $(\underline{s}_{0,t}, \hat{S}_{n,t})$. In the Appendix, we provide a counter example.*

We conclude this section by noting that the results regarding the structure of the optimal policy extend to the case where the ordering cost $c(x)$ in each period is a general increasing concave function. This follows from the fact that we can approximate, with arbitrary accuracy, an increasing concave function by a piecewise linear concave function.

4.4 Further Characterization of the Optimal Policy

In this section, we further characterize the optimal policy by showing that (1) the region over which the generalized (s, S) policy may not be optimal can be further reduced, (2) this region is increasing in $c_1 - c_n$, and (3) providing bounds on the optimal order-up-to levels $\hat{S}_{i,t}$.

First, let

$$\eta_{i,t} = \min \arg \min_y [\bar{L}_t(y) + (c_i - c_n)y],$$

i.e., $\eta_{i,t}$ is the global minimum of $\bar{L}_t(y) + (c_i - c_n)y$ for $i = 1, \dots, n$ assuming it exists. It is clear that $\eta_{1,t} \leq \eta_{2,t} \leq \dots \leq \eta_{n,t}$ since $c_1 > c_2 > \dots > c_n$.

Assumption 2. $\eta_{n,1} \leq \eta_{n,2} \leq \dots \leq \eta_{n,T}$.

Assumption 2 is satisfied if demands and costs are stationary.

First, we state a lemma which is useful in the further characterization of the optimal policy.

Lemma 13. *Under Assumption 2,*

(1) $\bar{v}_t^*(x)$ is nonincreasing in x for any $x \leq \eta_{n,t}$, and

(2) $\bar{G}_t(x_2) - \bar{G}_t(x_1) + (c_i - c_n)(x_2 - x_1) \leq \bar{L}_t(x_2) - \bar{L}_t(x_1) + (c_i - c_n)(x_2 - x_1) \leq 0$
for $x_1 \leq x_2 \leq \eta_{i,t}$.

Proof. (1) We show this result by induction. Observe that it is true for period T . Assume that it is true for period $t + 1$. It is clear that $\beta E[\bar{v}_{t+1}^*(y - \xi_t)]$ is nonincreasing in y for any $y \leq \eta_{n,t+1}$ by the inductive assumption since $\xi_t \geq 0$. Since $\eta_{n,t} \leq \eta_{n,t+1}$, it follows that $\bar{G}_t(y)$ is also nonincreasing in y for $y \leq \eta_{n,t}$. Let $x_1 < x_2 \leq \eta_{n,t}$. Define $a \vee b = \max\{a, b\}$ for some real numbers a, b . Let $y_t(x)$ be the optimal order-up-to level under state x . We have

$$\begin{aligned} \bar{v}_t^*(x_1) &= \bar{G}_t(y_t(x_1)) + \bar{c}(y_t(x_1) - x_1) \\ &\geq \bar{G}_t(y_t(x_1) \vee x_2) + \bar{c}(y_t(x_1) \vee x_2 - x_2) \\ &\geq \min_{y \geq x_2} [\bar{G}_t(y) + \bar{c}(y - x_2)] \\ &= \bar{v}_t^*(x_2). \end{aligned}$$

The first inequality is due to the fact that $\bar{G}_t(y)$ is nonincreasing in y for $y \leq \eta_{n,t}$ and the fact that $\bar{c}(y_t(x_1) - x_1) \geq \bar{c}(y_t(x_1) \vee x_2 - x_2)$ for any $x_2 \geq x_1$. The second inequality is due to the fact that $y_t(x_1) \vee x_2 \geq x_2$.

(2) For $x_1 \leq x_2 \leq \eta_{i,t}$, we have

$$\begin{aligned} & \bar{G}_t(x_2) - \bar{G}_t(x_1) + (c_i - c_n)(x_2 - x_1) \\ &= \bar{L}_t(x_2) - \bar{L}_t(x_1) + (c_i - c_n)(x_2 - x_1) + \beta E[\bar{v}_{t+1}^*(x_2 - \xi_t) - \bar{v}_{t+1}^*(x_1 - \xi_t)] \\ &\leq \bar{L}_t(x_2) - \bar{L}_t(x_1) + (c_i - c_n)(x_2 - x_1) \leq 0. \end{aligned}$$

The first inequality is due to the fact that $\bar{v}_{t+1}^*(x)$ is nonincreasing in x for any $x \leq \eta_{n,t+1}$. The second inequality is due to the fact that $\bar{L}_t(x) + (c_i - c_n)$ is nonincreasing in x for any $x \leq \eta_{i,t}$ and the fact that $\eta_{i,t} \leq \eta_{n,t}$. \square

Proposition 5. *Under Assumption 2, the optimal policy can be further characterized as follows.*

(1) *The generalized (s, S) policy is optimal for $x < \eta_{1,t}$.*

(2) *It is optimal not to order for $x > \eta_{n,t}$.*

(3) *$\eta_{1,t} - \eta_{n,t}$ is increasing in $c_1 - c_n$.*

(4) *$\hat{S}_{i,t} \geq \eta_{i,t}$ for $i = 1, \dots, n$.*

Proof. (1) We show that if it is optimal to order at some $x_2 < \eta_{1,t}$ it must be optimal to order at any x_1 such that $x_1 < x_2$. Then according to Theorem 10, the generalized (s, S) policy is optimal for $x < \eta_{1,t}$. Suppose that $c(y_t(x_2) - x_2) = K_i + c_i(y_t(x_2) - x_2)$. Then we have

$$\bar{G}_t(y_t(x_2)) + \bar{c}(y_t(x_2) - x_1) \leq \bar{G}_t(y_t(x_2)) + K_i + c_i(y_t(x_2) - x_1) - c_n(y_t(x_2) - x_1)$$

$$\begin{aligned}
&= \bar{G}_t(y_t(x_2)) + \bar{c}(y_t(x_2) - x_2) + (c_i - c_n)(x_2 - x_1) \\
&< \bar{G}(x_2) + (c_i - c_n)(x_2 - x_1) \\
&\leq \bar{G}(x_1).
\end{aligned}$$

The first inequality is due to the fact that $\bar{c}(x) = \min_i \{K_i + c_i x\} - c_n x$. The second inequality is due to the fact that $\bar{G}_t(y_t(x_2)) + \bar{c}(y_t(x_2) - x_2) < \bar{G}(x_2)$, i.e., it is optimal to order at x_2 . The last inequality is due to $\bar{G}_t(x) + (c_i - c_n)x$ being nonincreasing for $x < \eta_{1,t}$. Thus, it is optimal to order at x_1 .

(2) First, we show $\bar{v}_t^*(x_2) + \bar{c}(x_2 - x_1) \geq \bar{v}_t^*(x_1)$ for $x_2 > x_1$. Note that

$$\begin{aligned}
\bar{v}_t^*(x_2) + \bar{c}(x_2 - x_1) &= \bar{G}_t^*(y_t(x_2)) + \bar{c}(y_t(x_2) - x_2) + \bar{c}(x_2 - x_1) \\
&\geq \bar{G}_t(y_t(x_2)) + \bar{c}(y_t(x_2) - x_1) \\
&\geq \bar{v}_t^*(x_1).
\end{aligned}$$

The first inequality is due to the subadditivity of \bar{c} . Since $\bar{v}_t^*(x_2) + \bar{c}(x_2 - x_1) \geq \bar{v}_t^*(x_1)$ and $\bar{L}_t(x_2) \geq \bar{L}_t(x_1)$ for $x_2 > x_1 \geq \eta_{n,t}$, as a result we must have $\bar{G}_t(x_2) + \bar{c}(x_2 - x_1) \geq \bar{G}_t(x_1)$ for $x_2 > x_1 \geq \eta_{n,t}$. Hence, it is optimal to order nothing for $x \geq \eta_{n,t}$.

(3) This result follows directly from the definition of $\eta_{i,t}$.

(4) This result is due to property (2) of Lemma 13. \square

Note that results (2) and (3) of the above proposition hold even without Assumption 2.

As we can see, under Assumption 2, the generalized (s, S) policy is optimal except for the interval $(\eta_{1,t}, \eta_{n,t})$ and the width of this interval is increasing in $c_1 - c_n$, implying that when $c_1 - c_n$ is small, the width of $(\eta_{1,t}, \eta_{n,t})$ is also small.

For $K_1 = 0$, we can further reduce the interval over which the generalized (s, S) policy may not be optimal. Let $\bar{c}_1(x) \equiv c(x) - c_1x$,

$$\bar{L}_{1,t}(y) = (1 - \alpha)c_1[y - E(\xi_t)] + c_1E\xi_t + Eg_t(y - \xi_t),$$

and

$$\bar{G}_{1,t}(y) = \bar{L}_{1,t}(y) + \alpha E\bar{v}_{1,t+1}^*(y - \xi_t).$$

Then we can reformulate the dynamic recursion as follows.

$$\begin{aligned} \bar{v}_{1,t}^*(x) &= \min_{y \geq x} E\{\bar{c}_1(y - x) + (1 - \alpha)c_1[y - \xi_t] + c_1\xi_t + g_t(y - \xi_t) \\ &\quad + \alpha\bar{v}_{1,t+1}^*(y - \xi_t)\} \\ &= \min_{y \geq x} \{\bar{c}_1(y - x) + \bar{G}_{1,t}(y)\}, \end{aligned}$$

with $\bar{v}_{1,T+1}(x) = c_1x$. Note that $\bar{c}_1(x) \leq 0$ for all $x \geq 0$ and

$$\hat{S}_{1,t} = \min \arg \min_y [\bar{G}_t(y) + (c_1 - c_n)y] = \min \arg \min_y \bar{G}_{1,t}(y).$$

This leads to the following proposition.

Proposition 6. *If $K_1 = 0$, then the region over which we cannot fully characterize the structure of the optimal policy can be reduced to $(\hat{S}_{1,t}, \eta_{n,t})$.*

Proof. We first show that the generalized (s, S) policy is optimal for $x < \hat{S}_{1,t}$.

Note that the optimal ordering decision is given by

$$\min_{y \geq x} [\bar{G}_{1,t}(y) + \bar{c}_1(y - x)].$$

Since $\bar{c}_1(y - x) \leq 0$ and $\bar{G}_{1,t}(x) > \bar{G}_{1,t}(\hat{S}_{1,t})$ for $x < \hat{S}_{1,t}$, we have

$$\bar{G}_{1,t}(\hat{S}_{1,t}) + \bar{c}_1(\hat{S}_{1,t} - x) < \bar{G}_{1,t}(x)$$

for $x < \hat{S}_{1,t}$. Hence it is optimal to order a positive quantity under x for all $x < \hat{S}_{1,t}$. In turn, this implies, based on Theorem 10, that the generalized (s, S) policy is optimal for $x < \hat{S}_{1,t}$. By a similar argument as in Proposition 5, we can show that it is optimal not to order for $x > \eta_{n,t}$. \square

Note that since $\hat{S}_{1,t} \geq \eta_{1,t}$, the region in which the general (s, S) policy may not be optimal is reduced from $(\eta_{1,t}, \eta_{n,t})$ to $(\hat{S}_{1,t}, \eta_{n,t})$. Also, note that if $\hat{S}_{1,t} \geq \eta_{n,t}$, then a generalized (s, S) policy is optimal over the entire state space. Finally, note that we do not need Assumption 2 for Proposition 6.

4.5 The Optimality of the Generalized (s, S) Policy

In this section, we show that a generalized (s, S) policy is optimal for all regions of the state space if the single period inventory cost satisfies the following assumption.

Assumption 3. *The following inequality holds for all x_1, x_2 such that $|x_1 - x_2| \geq 1$ and $\theta \in (0, 1)$*

$$\theta \bar{L}_t(x_1) + (1 - \theta) \bar{L}_t(x_2) \geq \bar{L}_t(\theta x_1 + (1 - \theta)x_2) + \theta(1 - \theta)(c_1 - c_n)|x_2 - x_1|.$$

Remark 3. *This assumption is related to the concept of strong convexity. A function f is called strongly convex with parameter $m > 0$ if*

$$\theta f(x_1) + (1 - \theta)f(x_2) \geq f(x_1 + (1 - \theta)x_2) + \frac{1}{2}m|x_1 - x_2|^2,$$

which is equivalent to $f(x) - \frac{1}{2}mx^2$ being convex (See Rockafellar (2015)). It is clear that Assumption 3 is stronger than regular convexity but weaker than strong convexity with parameter $2(c_1 - c_n)$ since $|x_2 - x_1|^2 \geq |x_2 - x_1|$ for $|x_2 - x_1| \geq 1$.

One example that satisfies Assumption 3 is the case in which $g_t(x) = h_t \cdot (x^+)^2 + b_t \cdot (x^-)^2$, where $h_t \geq c_1 - c_n$ and $b_t \geq c_1 - c_n$.

Theorem 12. *If Assumption 3 is true and the minimum ordering quantity is always larger than or equal to 1, then the generalized (s, S) policy is optimal.*

Proof. We show this result by contradiction. Suppose that it is optimal to order under state x_0 and it is not optimal to order under state $x_0 - \delta$ for some $\delta > 0$. Let y be the optimal order-up-to level under state x_0 . Suppose that $y - x_0 \in (z_{i-1}, z_i]$ for some i . Then we have $\bar{G}_t(y) + \bar{c}(y - x_0) < \bar{G}_t(x_0)$ and $\bar{G}_t(y) + \bar{c}(y - x_0 + \delta) \geq \bar{G}_t(x_0 - \delta)$. There exists some $\rho_1 \in (0, 1)$ such that $\rho_1(x_0 - \delta) + (1 - \rho_1)y = x_0$. Note that for this ρ_1 , we have

$$\begin{aligned} & \bar{c}(y - x_0 + \delta) - (1 - \rho_1)(c_i - c_n)(y - x_0 + \delta) \\ & \leq K_i + (c_i - c_n)(y - x_0 + \delta) - (1 - \rho_1)(c_i - c_n)(y - x_0 + \delta) \\ & = K_i + (c_i - c_n)(y - x_0) \\ & = \bar{c}(y - x_0), \end{aligned}$$

which implies that

$$\bar{c}(y - x_0) \geq \rho_1[\bar{c}(y - x_0 + \delta) - (1 - \rho_1)(c_i - c_n)(y - x_0 + \delta)] + (1 - \rho_1)[\bar{c}(y - x_0)].$$

As a result, we have

$$\bar{G}_t(y) + \bar{c}(y - x_0)$$

$$\begin{aligned}
&\geq \rho_1[\bar{G}_t(y) + \bar{c}(y - x_0 + \delta) - (1 - \rho_1)(c_i - c_n)(y - x_0 + \delta)] \\
&\quad + (1 - \rho_1)[\bar{G}_t(y) + \bar{c}(y - x_0)] \\
&\geq \rho_1[\bar{G}_t(x_0 - \delta) - (1 - \rho_1)(c_i - c_n)(y - x_0 + \delta)] + (1 - \rho_1)[\bar{G}_t(y) + \bar{c}(y - x_0)] \\
&\geq \rho_1[\bar{L}_t(x_0 - \delta) - (1 - \rho_1)(c_1 - c_n)(y - x_0 + \delta)] + (1 - \rho_1)\bar{L}_t(y) \\
&\quad + \alpha E[\rho_1 \bar{v}_{t+1}^*(x_0 - \delta - \xi_t) + (1 - \rho_1)(\bar{v}_{t+1}^*(y - \xi_t) + \bar{c}(y - x_0))] \\
&\geq \bar{L}_t(x_0) + \alpha E \bar{v}_{t+1}^*(x_0 - \xi_t) \\
&= \bar{G}_t(x_0).
\end{aligned}$$

The third inequality is due to the fact that $c_1 \geq c_i$ for all i . The last inequality is due to Assumption 3 and the \bar{c} -convexity of \bar{v}_{t+1}^* . This contradicts the fact that $\bar{G}_t(y) + \bar{c}(y - x_0) < \bar{G}_t(x_0)$.

As a result, if it is optimal to order under x_2 , then it must be optimal to order under x_1 for any $x_1 < x_2$. Hence, from Theorem 10, the generalized (s, S) policy is optimal. \square

The assumption that the minimum ordering quantity is larger than or equal to 1 applies to the discrete demand case (all related proofs can be modified to accommodate the discrete demand case) and to the case in which K_1 is sufficiently large.

Note that under Assumption 2, Assumption 3 can be relaxed as follows: for all $x_1, x_2 \in (\eta_{1,t}, \eta_{n,t})$ such that $|x_1 - x_2| \geq 1$ and $\theta \in (0, 1)$, the following inequality holds

$$\theta \bar{L}_t(x_1) + (1 - \theta) \bar{L}_t(x_2) \geq \bar{L}_t(\theta x_1 + (1 - \theta)x_2) + \theta(1 - \theta)(c_1 - c_n)|x_2 - x_1|.$$

This is due to the fact that under Assumption 2, we only need to check whether Condition 1 holds for the interval $(\eta_{1,t}, \eta_{n,t})$.

4.6 Extensions to Other Settings

In this section, we briefly explain how our approach can be extended to time-varying ordering costs case, the lost sales case, and the non-zero leadtime case with backordering.

First, our result can be extended to the following time varying piecewise concave ordering costs: $c_t(x) = \min_i \{K_{i,t}\delta(x) + c_{i,t}x\}$ with $\alpha\bar{c}_{t+1}(x) \leq \bar{c}_t(x)$ for all $x \geq 0$ and all t , where subscript t denotes the dependency of the cost parameters on period t . By similar arguments as in the stationary ordering cost case, we can characterize the structure of the optimal policy and show that, except for a bounded region, it is a generalized (s, S) policy.

Next, we consider the case where unfulfilled demand is lost instead of backordered. Let p_t be the unit lost sales cost and h_t be the unit holding cost in period t . Let $v_t^*(x)$ be the value function in period t with starting inventory level x . Then, the function G_t can be modified as follows

$$G_t(y) = E(h_t[y - \xi]^+ + p_t[\xi - y]^+) + \alpha E v_{t+1}^*([y - \xi]^+); \quad (4.9)$$

and the optimality equation rewritten as

$$v_t^*(x) = \min_{y \geq x} [c(y - x) + G_t(y)]. \quad (4.10)$$

Using similar analysis to the one for the backordering case, we can show here too that the structure of the optimal policy, except for a bounded region, is also a generalized (s, S) policy.

Finally, consider the case with fixed leadtime l and backordering, where an order placed in period t is delivered in period $t + l$. Let x be the current inventory

in stock, and x_i be the amount of inventory delivered i periods later, where $i = 1, \dots, l-1$. By a standard transformation as in Zipkin (2000), we can show that the value function in each period only depends on the starting aggregate inventory level $x + x_1 + \dots + x_{l-1}$. Then, we can carry out similar analysis to the one for zero leadtime case and show that the optimal policy is again a generalized (s, S) except for a bounded region.

4.7 Appendix: A Counter Example

Here we provide a counterexample illustrating that the generalized (s, S) policy may not be optimal over the entire state space. In particular, we show that in this example there exist x_1 and x_2 such that $x_1 < x_2$ and it is optimal to order at x_2 but not optimal to order at x_1 .

We consider a 2-period problem with the ordering cost $c(x) = \min\{2x, x+192\}$, i.e., $c_1 = 2$, $c_2 = 1$, $K_1 = 0$ and $K_2 = 192$. Let

$$g_t(x - \xi_t) = \begin{cases} (x - \xi_t), & \text{if } x \geq \xi_t, \\ 3(\xi_t - x), & \text{otherwise,} \end{cases}$$

i.e., the unit holding cost is 1 and the unit backorder cost is 3 in every period. Let $\alpha = 1$. Demands in different periods are *i.i.d.* with density function ϕ given

as follows:

$$\phi(x) = \begin{cases} \frac{1}{1280}(x - 270), & \text{if } 270 < x \leq 280, \\ \frac{1}{1280}[-\frac{8}{5}(x - 280) + 10], & \text{if } 280 < x \leq 285, \\ \frac{1}{640}, & \text{if } 285 < x \leq 405, \\ \frac{1}{1280}[2(x - 405) + 2], & \text{if } 405 < x \leq 410, \\ \frac{1}{1280}[-2(x - 410) + 12], & \text{if } 410 < x \leq 415, \\ \frac{1}{640}, & \text{if } 415 < x \leq 540, \\ \frac{1}{1280}[-\frac{3}{20}(x - 540) + 2], & \text{if } 540 < x \leq 550, \\ \frac{1}{1280}[\frac{2}{5}(x - 550) + \frac{1}{2}], & \text{if } 550 < x \leq 555, \\ \frac{5}{2560}, & \text{if } 555 < x \leq 803, \text{ and} \\ 0 & \text{otherwise.} \end{cases}$$

It is clear that ϕ is not a Polya density function. Denote the distribution function of demand by Φ . Then $\Phi^{-1}(0.25) = 405$, $\Phi^{-1}(0.5) = 540$ and $\Phi^{-1}(0.75) = 675$. It is easy to check that in the second period, the optimal policy is a generalized (s, S) policy with $s_1 = S_1 = 405$, $s_2 = 270$ and $S_2 = 540$. Then we have

$$v_2^*(x) = \begin{cases} 192 + (540 - x) + E(540 - \xi_t), & \text{if } x \leq 270, \\ 2(405 - x) + E(450 - \xi_t), & \text{if } 270 < x \leq 450, \\ E(x - \xi_t) & \text{otherwise.} \end{cases}$$

In the first period, we can check that for any $y \geq 540$, $H_1(540) \leq c(y - 540) + H_1(y)$ and $H(550) > c(555 - 550) + H_1(555)$. Thus, it is not optimal to order at $x_1 = 540$

and it is optimal to order at $x_2 = 550$. Therefore, the optimal policy in the first period is not a generalized (s, S) policy.

In this example, we can show that the optimal policy in the first period is

$$y_1^*(x) = \begin{cases} 717, & \text{if } x \leq 443, \\ 540, & \text{if } 443 < x \leq 540, \\ x, & \text{if } 540 < x \leq 545, \\ 555, & \text{if } 545 < x \leq 555, \\ x, & \text{otherwise.} \end{cases}$$

Thus, the generalized (s, S) policy is optimal for $x < 540$ and it is optimal to order nothing for $x > 555$. This is consistent with our statement about the optimal policy in Theorem 11. Finally, we note that a counter example for $K_1 > 0$ can also be found with the same demand density function.

Chapter 5

Conclusions and Other Research Projects

In this chapter, we provide conclusions and future research directions on the work presented in Chapters 2 and 3. We also briefly discuss other research projects.

5.1 On the Impact of Input Price Variability and Correlation in Stochastic Inventory Systems

In Section 2, we examined the impact of input price variability on expected cost in inventory systems with stochastic demand and stochastic input prices. For a general class of such systems, we showed that higher input price variability leads to lower expected cost. We showed that this is true for a wide range of assumptions regarding price evolution, including i.i.d. prices and prices that evolve according to a Markovian martingale. We also showed that this is true for systems with both

single and multiple periods. We described how the impact of price variability on expected cost can be traced to the concavity of the cost function in input price, which is itself a consequence of the flexibility in adjusting the order quantity as prices vary. In addition, we examined the impact of price correlation over time and across inputs. We found that expected cost is increasing in price correlation over time and decreasing in price correlation across components. Numerical results suggest that higher correlation of prices over time diminishes the benefit derived from price variability while higher correlation of prices across components enhances it.

There are several avenues for future research. It would be useful to extend the analysis to broader classes of systems, including systems with multiple production stages where different components may be needed at different stages. In particular, it would be of interest to investigate how the position of a component in the production process affects the benefit derived from the variability in its input price (e.g., is price variability more beneficial for components that are upstream in the production process or is it more so for components that are downstream?). It would also be useful to consider settings in which there is variability in both the input purchase price and the output selling price. For example, a firm may purchase input from one spot market and sell output to another, with the firm observing both input and output prices at the beginning of each period and then deciding on how much input to buy and how much output to produce and sell. Lastly, it would be valuable to extend our analysis to settings where the firm may not be risk neutral and to account for its attitude toward risk by studying a decision criterion other than expected value.

5.2 Optimal Control of an Inventory System with Stochastic and Independent Leadtimes

In Section 3, we studied an inventory system with stochastic and independent leadtimes. For the case of exponentially distributed leadtimes, we resolved the open question regarding the structure of the optimal policy. In particular, we showed that the optimal policy is specified by a threshold function that is non-increasing in the inventory level. We showed that once the threshold function starts to decrease it continues to do so at a rate that is greater than or equal to one. This implies that the threshold function can be fully described by at most m parameters. Taking advantage of this structure, we provided an efficient algorithm for computing these parameters and the corresponding optimal cost. Also, inspired by the structure of the optimal policy, we investigated two plausible heuristics, as alternatives to the optimal, and examined their performance for a wide range of parameter values. We showed that the heuristics can perform poorly for certain parameter values. Finally, we extended our analysis to systems with lost sales and to systems where order cancellations are possible.

There are several possible avenues for future research. It would be of interest to extend the results to more general settings with respect to the distribution of demand and leadtime. It would also be of interest to extend the results to settings where leadtimes are not identical, as in systems with heterogeneous production facilities or delivery modes. We expect the analysis to be much more difficult in those cases, but there may be special cases for which at least partial characterization of the optimal policy is possible.

5.3 Other Research Projects

5.3.1 Managing Stochastic Inventory Systems with Scarce Resources

We consider a production-inventory system where the input material is scarce and its consumption is subject to a limit over a specified *compliance* period. Examples of such settings are many and include those where limits are imposed on the harvesting of forest products, the hunting and fishing of wild life, and the mining of rare minerals and metals. They also include settings where limits are imposed on the consumption of water or the emission of harmful pollution. In such cases, the amount that can be produced over the compliance period, which may consist of multiple production periods, cannot exceed the specified limit. Imposing such a limit introduces capacity dependencies across production periods, absent from traditional models where capacity constraints are imposed on individual periods. In particular, capacity in each period depends on the production decisions in previous periods and affects the capacity available in future periods. The objective of the system manager, in the face of stochastic demand, is to minimize the sum of inventory holding and shortage costs over a planning horizon consisting of one or more compliance review periods.

We formulated the problem as a stochastic dynamic program with a two-dimensional state space: on-hand inventory level and remaining capacity. We considered an extended state-space version of the problem and showed that this modified version of the problem reduces to a one-dimensional problem. We described various properties of the optimal policy for the modified version of the

problem and then showed that these properties also hold for the original problem. We then used these properties to characterize the structure of the optimal policy for the original problem. In particular, we showed that the optimal ordering policy is specified by dynamic thresholds that depend on both the on-hand inventory level and the remaining capacity but only via the sum of these two quantities. In addition, we characterized the impact of the capacity constraint and showed that the expected optimal cost is convex with respect to the remaining capacity, implying that there is diminishing value to capacity. We provided numerical results that examine the tradeoff between the expected optimal cost and the expected cumulative amount ordered, and discussed how both are affected by problem parameters. We evaluated the performance of three plausible heuristics that are simpler to compute and implement. We showed that, although the heuristics can be quite effective under some settings, they can also perform poorly under others. We then considered the problem of jointly optimizing for capacity and inventory control and showed that the associated total cost is convex in capacity and, therefore, the optimal capacity can be computed easily. We also showed that the optimal capacity can be quite sensitive to the price of capacity initially, with even modest prices leading to a significant reduction in the capacity purchased. Finally, we considered various extensions to the original model and show that the optimal policy of the extended models has similar structure.

There are several possible avenues for future research. It would be useful to generalize the results to a broader class of systems, including multi-stage systems where each stage may have its own cumulative ordering/production capacity constraint. It would also be useful to study systems with both cumulative

and period capacity constraints, where the period constraint may be due to production capacity limits while the cumulative constraint due to limits on input material availability or negative environmental externalities. Moreover, it would be interesting to compare systems where the cumulative amount ordered over the planning horizon is limited via an explicit constraint (as considered in this paper) to systems where this is achieved via imposing a penalty (or a tax) on ordering, or to systems where there are both a reward and penalty with production depending on whether the cumulative quantity falls below or over a specified threshold. For more details, please refer to Benjaafar et al. (2015b).

5.3.2 Stochastic Inventory Systems with Discount-driven Backorders

Stockouts are quite common for consumer products due to the variability in demand. Most of the inventory literature assumes either backorders or lost sales when stockouts occur. Existing literature that considers both backorders and lost sales assumes that when the on-hand inventory is not available to fulfill current demand, the inventory manager could decide whether to backlog or to reject some or all the demand. However, in practice, when stockouts occur, it is the customer herself who decides whether to wait for the product or walk away. The seller can offer a price discount to incentivize customers to wait for the product in the case of a stockout in order to mitigate lost sales.

In this study, we consider a multi-period stochastic inventory systems with both backorders and lost sales. In the case of a stockout, customers may choose to either wait for the product which corresponds to backorders, or walk away which

corresponds to lost sales. We assume that a fraction of the unfulfilled customers are willing to wait and this fraction depends on a discount the seller offers. The higher the discount is, the higher this fraction is, i.e., the more customers are willing to wait. We show that for a given discount, the optimal policy is a base stock policy. The optimal cost is convex in the discount and therefore the optimal discount can be computed easily. We also consider a continuous version of this problem. In this continuous review model, the probability that a customer would wait for the product is increasing in the discount. Again, we characterize the structure of the optimal policy and provide some managerial insights. These results are similar to those in the periodic review model.

We are currently extending the periodic review model by assuming that the backordering process is probabilistic, i.e., there is a range of possible outcomes, including with some positive probability that no customer would be willing to be backordered. In other words, given a discount, there is a distribution for the number of customers who are willing to be backordered. To do so, we introduce a notion of customer valuation of waiting (backordering), which is a random variable. If this value is less than the discount, the customer waits; otherwise the customer does not wait. This allows us to endogenize the probability of backordering and to study how the distribution of valuations affects the optimal discount and the corresponding base stock level. We are also extending the analysis to settings where lost sales in one period affect demand in future periods and those where a customer's probability of waiting is affected by current backorder levels. For more details, please refer to Chen et al. (2015b).

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