A Novel Data Set for Semantic Parsing using SQL as a Formal Language

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Abstract

For a learning system, in order to understand a natural language sentence, it has to map the natural language sentence onto a computer-understandable meaning representation. A semantic parsing learner maps a natural language sentence onto its corresponding meaning representation in a formal language. It involves a deeper analysis of the natural language sentence. The formal language used in this work is Structured Query Language (SQL). SQL is a computer-understandable meaning representation. In this thesis, we have focused on testing a semantic parsing learner on a new domain. We have tested the semantic parsing learner on a new data set of queries concerning the game of Cricket. Cricket is a game similar to a base-ball game but the rules differ. We have implemented the semantic parsing learner using trained Support Vector Machine classifiers employing the String Subsequence Kernels. Our system takes natural language sentences paired with their formal meaning representations in SQL as training data. For each production in the formal language grammar, a Support-Vector Machine (SVM) classifier is trained using the String Subsequence Kernel. These trained SVM classifiers are used for testing the semantic parsing learner on the test data set. The semantic parsing learner follows an iterative learning process. The experimental results have shown that the accuracy of the system improves with the number of iteration and the amount of training data used for learning.
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1 Introduction

1.1 Thesis Overview

For a learning system, in order to understand a natural language sentence, it has to map the natural language sentence on to a computer-understandable meaning representation. This task of mapping a natural language sentence onto a computer-understandable meaning representation is called semantic parsing. A sentence can be analyzed at a level called shallow analysis or at a deeper level called semantic analysis. Shallow analysis involves analyzing the text to infer some simple properties or to represent the text in some intermediate representation. This intermediate representation is used for later processing. Tasks like semantic role labeling (Gildea and Jurafsky (2002); Carreras and Marquez (2004)) and information extraction (Cardie (1997); Califf (1999)) involve shallow analysis of text. Deeper analysis involves analyzing the meaning of a natural language sentence. The task of semantic parsing involves deeper analysis of the text.

A semantic parsing learner maps a natural language sentence onto its corresponding meaning representation in a formal language. It involves a deeper analysis of the natural language sentence. The formal language used in this work is Structured Query Language (SQL). SQL is a computer-understandable meaning representation.

Previously, semantic parsing learners were developed using rule-based learning methods (Zelle and Mooney (1996); Tang and Mooney (2001); Kate et al. (2005)). These learning systems lack robustness. These systems are generally domain specific,
that is, when these systems are tested on a different domain, they do not perform well. Recent learning systems for semantic parsing used statistical feature-based methods (Ge et al. (2005); Zettlemoyer and Collins (2005); Wong et al. (2006)). These systems have performed better than the domain-specific systems. In this thesis, we have used kernel-based statistical learning methods for implementing the semantic parsing leaner, which is based on the work of Kate and Mooney (2006). The semantic parser we have implemented is based on Kernel-based Robust Interpretation for Semantic Parsing, KRISP algorithm (Kate and Mooney, 2006). In this work, we have used Support Vector Machine classifiers (Boser et al., 1992) with String Subsequence Kernel (kernel function) (Lodhi et al., 2002) for learning the semantic parsing system.

1.2 Thesis Contribution

In this thesis, we have focused on testing a semantic parsing learner on a new data set. We have implemented the semantic parsing learner using trained Support Vector Machine classifiers employing the String Subsequence Kernel. We have tested the semantic parsing learner on a new data set of queries concerning the game of Cricket. Cricket is a game similar to a base-ball game but the rules differ. The data set is a collection of data points. Each data point is a combination of a natural language sentence in English, its meaning representation in SQL and a correct parse tree for the SQL statement in context free grammar. The entire data set is then randomly divided into a test set and a train set. The system uses the training set for learning a semantic parser and the semantic parsing learner is tested on the test data set. For learning the semantic parsing system, an SVM classifier is trained for each production in the context free grammar using the String Subsequence Kernel. These trained SVM classifiers are used for testing the semantic parsing learner on
the test data set. While testing the system, it takes a natural language sentence, a
set of trained SVM classifiers and a context free grammar as input and predicts the
most probable meaning representation parse tree as output for that natural language
sentence. A trained SVM classifier gives the probability of a production covering a
substring of a given natural language sentence.

1.3 Thesis Outline

We have organized the thesis as follows:

• **Background (Chapter 2):** In this chapter, we have presented the background
  knowledge used for understanding semantic parsing and the work in this thesis.
  Following the work of Kate and Mooney (2006), We have used KRISP algo-
  rithm for developing a semantic parsing learner. We have explained the KRISP
  algorithm in detail. We gave an overview of the data set used for testing this
  system.

• **Implementation (Chapter 3):** In this chapter, we have explained the String
  Subsequence Kernel used in Support Vector Machines (SVM). We have ex-
  plained the role of String Subsequence Kernels in this thesis. We have explained
  about how the probability is estimated by the system using the SVM classifiers.
  In the last section in this chapter, we have presented a detailed explanation
  about the type of queries used in our data set for training and testing the
  semantic parsing system.

• **Results (Chapter 4):** In this chapter, we presented a detailed explanation
  about the experiments conducted and the results obtained on testing the accu-
  racy of the system. We have conducted two N-fold cross-validation experiments
for testing the system. We have presented the statistical figure obtained in the each experiment conducted.

• **Conclusion (Chapter 5):** In this chapter, we have presented few ideas on the future research based on this work. We have concluded by presenting the contributions of this thesis.
2 Background

A semantic parser maps a natural language sentence onto its corresponding meaning in a formal language. In this work, our formal language is Structured Query Language (SQL). This chapter provides the background knowledge for understanding semantic parsing and the work in this thesis.

A natural language sentence can be analyzed at a level called shallow analysis or at a deeper level such as semantic analysis. Tasks like syntactic parsing (Collins (1997); Charniak (1997)), information extraction (Cardie (1997); Califf (1999)), and semantic role labeling (Gildea and Jurafsky (2002); Carreras and Marquez (2004)) involve shallow analysis of natural language sentences. These tasks involve shallow analysis of a sentence to infer simple properties or to represent a sentence in some intermediate representation. Shallow analysis extracts only a limited amount of syntactic information from natural language sentences. For example, semantic role labeling involves finding the subject and predicate of a sentence and classifying them into their roles. Consider the sentence "-John gave a pen to Mike." The verb "to give" is marked as the predicate, John is marked as "agent" and "Mike" is marked as recipient of the action. These roles assigned to the events and participants in the text are drawn from a predefined set of semantic roles. This type of semantic role labeling gives an intermediate semantic representation of a sentence, which indicates the semantic relations between the predicate and its associated participants. This intermediate semantic representation is used for latter processing. The long-distance relationships in the text can be captured well by deeper semantic analysis. Deeper
semantic analysis involves analyzing the meaning of a natural language sentence. Unlike shallow semantic analysis, deeper semantic analysis generates a representation of the meaning of sentence.

In this thesis, we have focused on the task of testing a semantic parsing learner on a new data set. The semantic parser maps a natural language sentence onto its corresponding meaning representation in SQL. For a learning system, in order to understand the natural language sentence, it has to map the natural language sentence onto a computer-understandable meaning representation. SQL is well known computer-understandable meaning representation. This task of developing a semantic parsing learner involves deeper semantic analysis as it deals with the meaning of the natural language sentence. The semantic parser we implemented is based on Kernel-based Robust Interpretation for Semantic Parsing, KRISP algorithm (Kate and Mooney, 2006).

2.1 Semantic Parsing

SQL (Structured Query Language) is a special purpose programming language used for interacting with relational database management systems. SQL is responsible for querying and editing information stored in a database management system. SQL is a widely-used database language. The general structure of an SQL query for retrieving data would be "select SB from FB where WB." Here, SB refers to the attributes we want to retrieve, FB refers to the database table from which the data has to be retrieved and WB refers to the condition in the query. The terms SB, FB and WB are parameters in the above general query. For example, "select coach from odiplayers where country = 'england' " is an SQL query to find the coach of the England cricket team. In this query, the attribute we are interested in is "coach,"
the table from which the data is retrieved is "odiplayers" and the phrase "country = 'england' " is the condition in the query. This is just a simple example of an SQL query. There can be very complex queries, which involve retrieving multiple attributes, accessing multiple tables, specifying multiple conditions and so on. This will be explained in detail in section 2.4.

The following are a few domains in which semantic parsers have been developed.

**Air Travel Information Services (Price, 1990)**: The domain of the ATIS system is a benchmark for speech recognition and understanding. The corpus for this system consists of spoken natural language sentences about air travel coupled with their meaning representation in SQL. The system was built by engaging the subjects in dialogs through speech in natural language sentences. It is Transformation-based learning for semantic parsing (Brill, 1995).

For example, a natural language sentence might be "Display the list of flights from Minneapolis to Chicago"
its corresponding query would be "SELECT flight_id FROM flight WHERE from_airport= 'Minneapolis' AND to_airport = 'Chicago' "
The above example shows the connection between natural language and the semantic meaning representation in the ATIS system.

**Geographical Database (Geobase)**: These databases consisted of data-related geographical facts. The meaning representation for this system is a query language Prolog, augmented with several meta-predicates (Zelle and Mooney, 1996). These queries were converted into a functional and variable-free query language (Kate and Mooney, 2006), which is more convenient for some semantic parsers.

\(^1\)https://pantherfile.uwm.edu/katerj/www/publications.html
For example, a natural language sentence might be "Which rivers run through the states bordering Texas?"

its corresponding query in Prolog would be "answer(A, (river(A), traverse(A,B), state(B), next_to(B,C), const(C, stateid('texas'))))"

in the functional query language this would be: "answer(traverse(next_to (stateid('texas'))))"

The above example shows the connection between natural language and semantic meaning representation. Each expression in the functional query returns a list of entities. The expression, stateid('texas') in the functional query returns a list containing a single constant, that is, the state 'texas.' The expression next_to(stateid(texas)) returns the list of all the states next to the state 'texas.' The expression traverse(next_to(stateid(texas))) returns the list of rivers which flow through the states next to Texas. This final list of rivers is returned as the answer. The expression traverse(next_to(stateid(texas))) in functional query language corresponds to the binary predicate traverse(A,B) of the query in Prolog. It is true if A flows through B. The remaining expressions can be correlated in a similar manner. Figure 2.1 shows the meaning representation for this set of data in the form of a parse tree that captures the semantic meaning.

**Non-terminals:** ANSWER, RIVER, TRAVERSE, STATE, NEXT TO, STATEID

**Terminals:** answer, traverse, next to, stateid, texas

**Productions:**

\[
\begin{align*}
\text{ANSWER} & \rightarrow \text{answer}(\text{RIVER}) \\
\text{RIVER} & \rightarrow \text{TRAVERSE}(\text{STATE}) \\
\text{STATE} & \rightarrow \text{NEXT TO}(\text{STATE}) \\
\text{STATE} & \rightarrow \text{STATEID}
\end{align*}
\]
Figure 2.1: Parse tree of the MR in functional query language

\[ TRAVERSE \rightarrow \text{traverse} \]
\[ NEXT\ TO \rightarrow \text{next to} \]
\[ STATEID \rightarrow \text{stateid 'texas'} \]

CLang (The RoboCup Coach Language): RoboCup (Adorni et al., 1999) is an international AI research initiative using robotic soccer as its primary domain. The Coach Competition (Kuhlmann et al., 2006) is one of the several competitions organized under it. The soccer teams compete on a simulated soccer field. The teams compete by receiving coaching advices from their agent (Coach) in a standard formal coaching language called CLang (Chen, 2003). CLang is a simple declarative language with a prefix notation similar to LISP. Figure 2.2 gives an example of a piece

\(^2\text{http://www.robocup.org/}\)
of coaching advice in natural language (NL) with its corresponding CLang meaning representation. In the meaning representation, *bowner* stands for ball owner and *UNUM* stands for uniform numbers (players 1 through 11).

The CLang corpus used in semantic parsing experiments was constructed by randomly selecting 300 pieces of coaching advices from the log files of 2003 RoboCup Coach Competition. These formal advice instructions were translated into English by one of four annotators. On average there were 22.5 words per sentence and 13.42 tokens per meaning representation in this corpus.

For example, a natural language sentence might be "If our player 4 has the ball, our player 4 should shoot."
its corresponding query in CLang would be "((bowner our 4) (do our 4 shoot))"

**Terminals:** bowner, our, 4, shoot

**Non-terminals:** RULE, CONDITION, DIRECTIVE, TEAM, UNUM, ACTION
Productions:

\[
\text{RULE} \rightarrow (\text{CONDITION DIRECTIVE})
\]

\[
\text{CONDITION} \rightarrow (\text{bowner TEAM UNUM})
\]

\[
\text{DIRECTIVE} \rightarrow (\text{do TEAM UNUM ACTION})
\]

\[
\text{TEAM} \rightarrow \text{our}
\]

\[
\text{UNUM} \rightarrow 4
\]

\[
\text{ACTION} \rightarrow \text{shoot}
\]

2.2 Learning Semantic Parsers

Some earlier learning systems (Zelle and Mooney (1996); Tang and Mooney (2001); Kate et al. (2005)) used rule-based learning methods for semantic parsing. These methods lacked robustness; these methods are domain specific. When these systems are tested on a different domain, they do not perform well. Recent learning systems for semantic parsing (Ge et al. (2005); Zettlemoyer and Collins (2005); Wong et al. (2006)) have used statistical feature-based methods. These systems have performed better than systems developed in the past.

In this thesis, we have used a kernel-based statistical method for semantic parsing based on the work of Kate and Mooney (2006). Over the last few years, kernel methods have become very popular and an integral part of machine learning. Kernel methods have been successfully applied to a number of real-world problems and are now used in various domains. The advantage of kernel methods over feature-based methods is that they can work with a large number of features. Features are the numerical representation of some objects related to the data under consideration. For example, when dealing with text data, methods such as the bag of words, part
of speech tags and frequency of term occurrences can be used as features related to that text.

Consider the task of word sense disambiguation (WSD) (Ide and Veronis, 1998). It is the task of assigning the correct sense to a word in a given context. For example, the word ”bass” may refer to a kind of fish or a musical instrument depending up on the context in which it is used. If a feature-based method is used for WSD, the first task is to select the features. These features may include collocational features or bag-of-words features. The collocational features give the information about the words surrounding the target word and their positions with respect to the target word. A typical set of collocation features includes the set of context words with their positions and their part-of-speech tags. For example, consider the sentence, "The guitar and bass player stand at the corner.”

The task is to assign a correct sense to the word ”bass” in the above sentence. A collocational feature-vector formed using two words on either side of the target word with their respective positions and their part-of-speech tags would look like,

\[\text{guitar}_2, \text{NN}, \text{and}_1, \text{CC}, \text{player}_1, \text{NN}, \text{stand}_2, \text{VB}\]

The above feature vector uses a window size of two words on either side of the target word. This feature vector is used as an input to a machine learning algorithm. Unlike the collocational features, the bag-of-words features do not give importance to the position of the words considered in the feature vector. These features capture the context of the word in which it is used. A feature vector includes a set of words, which are frequently used in the context with the target word. These context words are collected from the sentences that use the target word. For example, for the above sentence, the bag-of-words feature vector for the word ”bass” looks like,

\[\text{fishing}, \text{big}, \text{sound}, \text{player}, \text{fly}, \text{rod}, \text{pound}, \text{double}, \text{runs}, \text{playing}, \text{guitar}, \text{band}\]
These set of words are frequently used in the context with the word "bass." These words are collected from the sentences that have the word "bass" in them, from the Wall Street Journal corpus (Charniak and et al, 2000). Using the words in the above feature-vector, if we consider a window size of 3 on either side of the target word "bass" in the above sentence, a binary feature vector is formed as below

\[ [0,0,0,1,0,0,0,0,1,0] \]

A zero value indicates that the word in the above bag-of-words feature vector is not present in the given sentence for a given window size. For example, the first zero in the binary feature vector indicates that the word "fishing" is not present in the window of 3 words on the either side of the word "bass" in the sentence. The value one in the binary feature vector indicates the occurrence of a particular word in the window of the given sentence.

Machine learning algorithms use these features vectors for statistical analysis. But it is computationally impractical for the feature-based methods to handle all of the possible features. They consider only a predefined set of features for statistical analysis, which leads to a loss of information. For example, the feature vectors extracted in the above example used a limited window size of words for both the collocational feature vector and the bag-of-words feature vector. For the collocational feature vector, only two words on either side of the target word are considered in the above example. For bag-of-words feature vector, only three words on either side of the target word are considered in the above example. There may be some words in that context which are outside the predefined window size and still play a critical role in assigning a sense to the target word. It is computationally impractical to include all the words that occur in that context. Feature-based methods cannot capture the longer dependencies between the words in a given context. Unlike feature-based methods, kernel-based methods are capable of handling infinitely many features. Kernel-based
machine learning algorithms use different functions to compute similarity between the data points. The similarity between data points is estimated by the kernel functions.

Support Vector Machines (Boser et al., 1992) are a class of machine learning algorithms, which use kernels and a linear or quadratic programming method. Support Vector Machines (SVMs), when applied to text, can construct a hyperplane in a high-dimensional space which can be used for text classification. Semantic parsing is a specific form of data classification task. Given a set of training data examples and the class they belong to, an SVM classifier trained on this data is capable of assigning a new example to one of the two possible classes. SVMs use data to find the similarity between the data points. When two sets of data points consisting of positive and negative points are given, an SVM constructs a hyperplane which can separate these two classes of data points such that it maximizes the distance between the hyperplane and the closest data points on either side. When a new data point is given, the SVM can decide which side of the hyperplane it belongs to. In this work, we have used String Subsequence Kernel (kernel function) (Lodhi et al., 2002) for computing the similarity scores between data points. The String Subsequence Kernel is explained in detail in the next chapter.

In Figure 2.3, there are two classes of linearly separable data points represented by circles and triangles. There can be many hyperplanes that separate these two classes of data points. The task is to find a separating hyperplane, which has a maximum distance from the closest data points on either side. The best hyperplane is the one, which has the maximum margin between the two support vectors. In the figure, the separating hyperplane is shown by a red line and black lines show the support vectors. Optimization methods are used to find the best separating hyperplane. The data points shown in the Figure 2.3 are linearly separable in the input space. In this case, the SVM need not use any kernel function. But there can be data which cannot
be separated linearly in the input space. For example, the distribution of data points in Figure 2.4 is not linearly separable in the input space. In order to separate the two classes of data points, a non-linear curve is used as shown in the Figure 2.4. Rather than fitting non-linear curves to separate the data points, SVM handles this by using a kernel function to map the data points into a higher dimensional space called feature space, where a hyperplane can be used to do the separation. The kernel function used in our work is a String Subsequence Kernel. A String Subsequence Kernel considers all the common subsequences between the sentences as features. For example, consider the two sentences,

"how many fours are hit by the batsman"

"number of fours hit by the batsman"
These two sentences have some words in common. A String Subsequence Kernel estimates the similarity score between the two sentences by considering all the possible substrings that are common in both the sentences. These common substrings are called subsequences. All these possible common subsequences between sentences are considered as features by the kernel. By considering all the subsequences as features, the SVM kernel operates in the feature space as opposed to a linear SVM shown in Figure 2.3, which operates in the input space.

2.3 Kernel-Based Robust Interpretation for Semantic Parsing (Kate and Mooney, 2006)

We have divided this section into five subsections. The first subsection gives an overview of the KRISP training algorithm. It gives an overview of the steps involved in the KRISP training algorithm. In the next two subsections, we have
defined and explained the concepts of a most probable MR parse tree and constant productions. We have explained about the extended Earley’s parser in detail in the next subsection. The KRISP training algorithm invokes the extended Earley’s parser to get the most probable MR parse trees. In the last subsection, we have presented a detailed explanation about the KRISP training algorithm.

2.3.1 KRISP Training Algorithm Overview

For learning a semantic parser, the KRISP algorithm takes training data, which includes natural language sentences along with their formal representations of their meaning in SQL and their parse trees in a context free grammar. For each production in the parse tree, an SVM classifier is trained using a String Subsequence Kernel. A kernel is a function, which estimates the similarity between data points. Each of these classifiers can then estimate the probability of a production representing a substring of a natural language sentence. For a given natural language sentence, the semantic parser uses these classifiers to estimate the probabilities for all possible substrings to finally give the most probable meaning representation in SQL.

Figure 2.5 gives an overview of the KRISP algorithm. The learning process for the semantic parser is iterative in nature. In each iteration, the semantic parser attempts to improve its performance. For the given training data set, the KRISP first parses the meaning representation using the context free grammar provided. For each production in the grammar, it collects the positive and negative examples. Before starting the first iteration, the positive examples $POS(P_{s_k[i,j]})$ for a production $P$ are all those natural language sentences $s_k$, which use production $P$ in the parse tree of their meaning representation. The rest of the data points $s_l$ are the negative examples $NEG(P_{s_l[i,j]})$ for this production. The representation, $s[i,j]$ stands for the
substring of a sentence $s$, starting from the $i^{th}$ word to the $j^{th}$ word in that sentence. In the subscript, the variables $i$ and $j$ indicate the start and end index of the part of the sentence that is considered either as a positive example or negative example. In the first iteration, the entire sentence is considered as a positive or a negative example. So, the starting index would be zero and the ending index would be the length of that string. These indices would vary in the subsequent iterations. Using these positive and negative examples, a Support Vector Machine classifier $C(P)$ is trained for each production.
The SVM uses String Subsequence Kernel (Lodhi et al., 2002) to estimate the similarity between strings. The similarity score depends on how similar the two sentences are. For example, consider the two pairs of sentences below. The sentences in the first pair are more similar when compared to the sentences in the second pair. In the first pair, the two sentences share a good number of words. In the second pair, the two sentences do have many strings in common. In this example, the similarity score for the first pair would be greater than the similarity score for the second. The procedure to compute the similarity score is explained in detail in the next chapter.

"how many fours are hit by the batsman"
"number of fours hit by the batsman"

"how many fours are hit by the batsman"
"total number of wickets taken by the bowler"

Figure 2.6 is the parse tree for the meaning representation of a natural language sentence "who is the coach for the england cricket team." We will use this figure to explain the KRISP training algorithm.

A natural language sentence for the above parse tree might be: "who is the coach for england cricket team."
Its corresponding meaning representation (MR) in SQL would be: "select coach
In the first iteration, an extended Earley’s parser (Earley, 1970) is invoked to get the n most probable MR parse trees for a given natural language sentence using the trained SVM classifiers. The nodes of the parse tree are represented as \((P,(a,b))\), where \(P\) represents the production and \((a,b)\) represents the substring \(S[a,b]\) of the natural language sentence \(S\). For example, the natural language sentence in the above
figure is of length 8. In the parse tree, the node at the production \((WB \rightarrow B = WB)\) is represented as \((WB \rightarrow B = WB, (3, 6))\). The representation \((3, 6)\) indicates the substring from the third word to the sixth word of that string. It suggests that the production \((WB \rightarrow B = WB)\) covers the substring starting from the third word to the sixth word. The substrings covered by the child nodes should not overlap. The substrings covered by the child nodes form the substring covered by the production at their parent node. For example, all the child nodes of the node \((WB \rightarrow B = WB, (3, 6))\) should cover the substring from the third to sixth word of \(S\) and none of these substrings should overlap.

After the first iteration, the semantic parser may not return the correct parser trees. Before starting the next iteration, the algorithm collects the positive and negative examples for each production in the MR parse tree. The positive examples for the next iteration are collected from the most probable parse tree which gives the correct meaning representation. The negative examples are collected from the most probable parse trees which give an incorrect meaning representation. These positive and negative examples are more precise when compared to the positive and negative examples collected before first iteration. Using these positive and negative examples, the KRISP algorithm trains the semantic parser again, that is, it trains the SVM classifiers with the new sets of positive and negative examples. This newly trained semantic parser is an improved one over the previous one. The SVM classifiers are trained for a desired number of iterations using the positive and negative examples created at each iteration. The final set of classifiers is returned at the end of last iteration.
2.3.2 Most Probable MR Parse Tree

Let $P_\pi(S[i,j])$ be the probability of the production $\pi$ representing the substring $S[i,j]$. In Figure 2.6, the probabilities are not shown. The probability of an MR parse tree is defined as the product of the probabilities of all the productions in it covering their respective substrings. During the semantic parsing, some of the productions with probability less than a threshold value are ignored while generating the most probable parse trees. In order to find the most probable parse trees, the KRISP invokes the extended Earley’s parser to get the $x$ most probable parse trees which is discussed in the next section.

2.3.3 Constant Productions

Some productions may contain constant terms on the right-hand side. For example, consider the production $ODIPLAYERS \rightarrow 'Sachin.'$ It refers to some specific substring in the given natural language sentence. For such productions no classifier is trained. Such productions are considered as constant productions. When such constant productions are encountered, KRISP algorithm assigns a probability 1 for them. For example, consider a natural language sentence "number of fours hit by Sachin," which contains the word 'Sachin.' During the semantic parsing, the algorithm directly uses the production $(ODIPLAYERS \rightarrow 'Sachin')$ to represent the string 'Sachin' in the final parse tree and assigns it the probability 1. In the set of productions used in the Figure 2.6, the production $(WB \rightarrow 'england')$ is a constant production as it refers to a specific word in the natural language sentence.

A natural language sentence may be expressed in more than one way or the order of the terms used for expressing a natural language sentence may differ from user to user. In order to address this issue, the KRISP algorithm considers all the possible
permutations of non-terminals on the right hand side of each production. The next section explains about how the extended Earley’s parser derives the most probable parse trees.

### 2.3.4 Extended Earley’s Parser (Kate and Mooney, 2006) for Most Probable MR Parse Trees

The KRISP algorithm uses extended version of Earley’s parser (Earley, 1970) for finding the most probable parse tree. The basic Earley Parser is an algorithm that can parse a given string that can be represented in a given context free grammar. For a given string $s$, it does the parsing from left to right by filling an array called the chart which is of length $|s| + 1$ (length of the string plus 1). Each chart entry is a list of states, which represent the subtrees that are parsed so far. No duplicate states are allowed in the chart entries. When a state is completed, it is used by other subtrees which need them for their completion or changing their state. The parser reads the words from left to right and checks if there is any rule that is allowed to be used by these words. At each word entry, the parser stores a list of partially completed rules called states in the chart. At each chart entry position, the parser predicts new grammar rules that could be started for that word. The parser then determines if a partially completed rule needs the word at that position to complete itself further. If a rule gets completed, it is used for the completion of the other rules which may need it. This rule completion process is repeated. It follows a dynamic programming strategy. A modified version of Earley’s parser is used by KRISP for finding the most probable parse tree.
Extended Earley’s Parser (Kate and Mooney, 2006)

A chart entry consists of a list of states. A state consists of the information about

1) The root production of the subtree.
2) The position in the sentence where the subtree has started.
3) The non-terminal on the right hand side up to which the production is completed.
4) The position in the sentence at which the subtree ends.
5) The probability of the subtree completed so far.

All the above described information is represented as a production (state) as described below. The terms state and production are interchangeable. Consider an example of a state representing the above information, \((3 WB \rightarrow B \bullet_4 = WB, 0.76)\). In this production, the number 3 on the left hand side of the production indicates that this subtree starts at the 3rd word in the given natural language sentence. The dot (•) on the right hand side of the production after the non-terminal B indicates the subtree corresponding to the non-terminal B is completed. As the dot is not at the end of the production, it also indicates that the subtree (state) corresponding to the non-terminal WB on the left-hand side is only partially completed. The number 4 after the dot indicates that the subtree so far completed covers the substring starting from the third word to the 4th word in the sentence. The probability of the subtree completed so far is 0.76. This is said to be an incomplete state. Consider the production \((4 WB \rightarrow B = WB \bullet_6, 0.96)\). The subtree corresponding to the non-terminal WB on the left-hand side is said to be complete because the dot position on the right-hand side is after the last non-terminal. This is said to be a completed state. It also indicates
that all the subtrees corresponding the non-terminals on the right-hand side of the this production are also complete. The subtree corresponding to the non-terminal \(WB\) on the left-hand side covers the substring starting from the third word to the 6th word in the given sentence. The probability of the production \((4W B \rightarrow B = WB \bullet_6, 0.96)\) representing the substring starting from the third word to the sixth word is 0.96. The productions which do not contain any non-terminals on the right-hand side, are called base production or constant production. For example, the production \( (WB \rightarrow 'england')\) is a constant production. No classifier is learned for such productions. They are assigned a probability of 1.

In each iteration, the KRISP training algorithm invokes the extended version of the Earley’s parser by passing it the natural language sentence, the set classifiers and the meaning representation grammar (context free grammar). The parser can return a large number of parse trees. The parser does a beam search to return the \(x\) most probable parse trees. The variable \(x\) is a system parameter called beam width. The beam need not necessarily contain the correct parse trees for a given sentence. It only considers the \(x\) most probable parse trees for the given sentence.

From the work of Kate and Mooney (2006), we have used the pseudo code shown in the Figure 2.7 for implementing the extended Earley’s parser. In the pseudo code, the symbols \(\alpha\), \(\beta\), \(\gamma\) are used to represent the sequence of terminals and non-terminals. The capitalized words in the productions represent non-terminals. The extended Earley’s parser takes three parameters, the natural language sentence, the context free grammar \(G\) and the SVM classifiers \(P\). The parser starts by inserting a dummy state \((\text{NULL} \rightarrow \bullet_0 Q)\) into the first chart entry. We use \(Q\) as the start symbol for the context free grammar. The start symbol is on the right-hand side of this dummy state. The dot position and the subscripting number indicate that the parser has not yet parsed any substring in the sentence. It then starts parsing from left to right,
Function Extended_Earley(Sentence s, MRL grammar, Classifiers P)
INSERT((o NULL → ●0 start, 1), chart[0])
for i=0 to |s| do
  for each state in chart[0.....i-1] do
    if (BASE(state) and INCOMPLETE(state))
      then SCANNER(state, i)
  for each state in chart[i] do
    if (not BASE(state) and INCOMPLETE(state))
      then PREDICTOR(state)
    else if (BASE(state) and INCOMPLETE(state))
      then SCANNER(state, i)
    else COMPLETER(state)
return(chart)

Procedure PREDICTOR((A → α • j B β, p))
for each (B → γ) in MRL grammar do
  for each permutation γ ′ of γ do
    INSERT(( j B →● j γ ′, 1), chart[j])

Procedure SCANNER((A → ● i α, p), k)
if (p = PA→α(s[i..k]) ≥ θ) then
  INSERT((A → α • i+1 p), chart[k+1])

Procedure COMPLETER((B → γ • k, p))
for each ((A → α • j B β, q) in chart[j]) do
  if (INCOMPLETE((A → α B • k β, p * q)))
    INSERT((A → α B • k β, p * q), chart[k])
  else if (r = PA→α B β(s[i..k−1]) ≥ θ) then
    INSERT((A → α B • k β, p * q * r), chart[k])

Procedure INSERT(state, chart[j])
if (state is not already in chart[j]) then
  BEAM_PUSH(state, chart[j])

Figure 2.7: Pseudo code for extended Earley’s parser

starting from the first word in the sentence. There are three main routines in this parser. They are PREDICTOR, SCANNER and COMPLETER.

The PREDICTOR routine adds new states to the chart entries in order to find
the possible parse trees. For example, consider the production \((3 \text{WB} \rightarrow B = \bullet_5 \text{WB}, 0.86)\). When the PREDICTOR routine is called on this production, the routine adds a new state \((5 \text{WB} \rightarrow \bullet_5 'england,' 1)\) to the chart entry. It does this by considering the non-terminal on the right hand side of a production after the dot location. The PREDICTOR routine is called on a non-constant production and incomplete production. For examples, PREDICTOR cannot be called on the production \((5 \text{WB} \rightarrow 'england' \bullet_6, 1)\) as it is constant or base production. If the dot location is at the end of a production, then it indicates that the state is complete. In the above example, when the PREDICTOR routine is called on the production \((3 \text{WB} \rightarrow B = \bullet_5 \text{WB}, 0.86)\), it adds all the states that have the non-terminal "WB" on the left-hand side. The PREDICTOR routine is called on incomplete states. The probability assigned to a predicted state indicates the probability of that predicted state covering a specific substring. Out of all possible substrings a predicted state can cover, if there is no such substring for which the assigned probability is greater than a threshold value, then that predicted state is not included in the chart entry. The threshold value is a system parameter. This is done because we are looking for the \(x\) most probable parse trees. As explained previously, a parser can generate infinitely many parse trees with varying probabilities. The beam only includes the \(x\) most probable trees. The symbol \(\theta\) indicates the threshold value in the pseudo code (Figure 2.7). The PREDICTOR also looks for all the permutations of the non-terminals on the right-hand side of a state. For example, consider the state \((3 \text{WB} \rightarrow \bullet_3 B = \text{WB}, 1)\). In this state, the non-terminals on the right-hand side are "B" and "WB." The PREDICTOR adds the states \((3 \text{WB} \rightarrow \bullet_3 B = \text{WB}, 1)\) and \((3 \text{WB} \text{WB} \rightarrow \bullet_3 \text{WB} = B, 1)\) to the chart entry. The symbol "=" is not considered as a non-terminal. As explained previously, there can more than one natural language sentence that imply the same meaning representation. For example, "who is the coach for england cricket team"
and "for england cricket team, who is the coach" are two sentences which imply the same meaning. The order of the terms differs here. In order to address this issue, the PREDICTOR considers all the permutations of the non-terminals on the right-hand side of a state. The predicted states are initially assigned a probability of one, which later gets multiplied by actual probabilities as it gets completed.

The SCANNER routine is called on a state if it is a base production and is incomplete. For example, \((5WB \rightarrow \bullet_5 'england,' 1)\) is a base production but is an incomplete state as the dot location is not at the end of the production. The SCANNER looks for that specific word on the right-hand side of the production in the natural language sentence. In this case, it looks for the word "england" in the natural language sentence and assigns a probability of 1 for that. It then adds a completed state \((5WB \rightarrow "england" \bullet_6, 1)\) to the chart next entry if the probability of this state is greater than the threshold value.

The COMPLETER routine is called on a complete state. Some of the previously added states in the chart may need this completed state for their completion or for moving into a new state. For example, consider the production \((3WB \rightarrow B = \bullet_5 WB, 0.86)\), the PREDICTOR is called on this state to add the new state \((5WB \rightarrow \bullet_5 'england,' 1)\). This is a base state and an incomplete state. The SCANNER routine is called on this state to add a complete state \((5WB \rightarrow "england" \bullet_6, 1)\). The COMPLETER routine looks for all the states in the chart which may need this completed state. In this example, the state \((3WB \rightarrow B = \bullet_5 WB, 0.86)\) needs the completed state \((5WB \rightarrow "england" \bullet_6, 1)\) for its completion. So, when the COMPLETER routine is called on the completed state \((5WB \rightarrow "england" \bullet_6, 1)\), a new state \((3WB \rightarrow B = WB \bullet_6, 0.86)\) is added to the chart entry. The probability of this new state is the product of the probability of previous state \((p)\) and the probability of the state on which the COMPLETER is called \((q)\). In this case, \(p\) is
the probability of the state \((3WB \rightarrow B = \bullet_5 \ WB, 0.86)\) and \(q\) is the probability of the completed state \((5WB \rightarrow \text{"england"} \bullet_6, 1)\). So, the probability of the new state \((3WB \rightarrow B = WB \bullet_6, 0.86)\) would be \(p \times q\), that is, 0.8*1.

On calling the COMPLETER routine, if the new state added is also a completed state (as in the above example), then it gets a probability \(k\) assigned by an SVM classifier as well. For example, assume that the new state obtained after calling the COMPLETER routine on the completed state \((5WB \rightarrow \text{"england"} \bullet_6, 1)\) is \((3WB \rightarrow B = \bullet_5 \ WB, 0.91)\). In this case, the state \((3WB \rightarrow B = \bullet_5 \ WB, 0.86)\) whose probability is \(p\) has used the completed state \((5WB \rightarrow \text{"england"} \bullet_6, 1)\) whose probability is \(q\) to give a new completed state \((3WB \rightarrow B = WB \bullet_6, 0.91)\) with probability \(k\). As this new state is a completed state, the probability \(k\) (0.91) is assigned by the SVM classifier. The probability \(k\) is the probability of that state covering a substring from third word to the sixth word of the given sentence and it has to be greater than the threshold value for it be considered further. This probability \(k\), is not its final probability. This probability \(k\) is multiplied by the probabilities \(p\) and \(q\) to give its final probability, that is, \(k \times p \times q\). In this case, the probability of the new state \((3WB \rightarrow B = WB \bullet_6)\) would be 0.91*0.86*1.

The common operation in the above three routines is to add the new states to the chart. This is done by calling an INSERT routine, which adds new states to the chart. INSERT routine does not add duplicate states. The extended Earley’s parser returns the \(x\) most probable parser trees.
2.3.5 Krisp Training Algorithm (Kate and Mooney, 2006) in Detail

This section presents a detailed explanation about the KRISP training algorithm. The training data set consists of the natural language sentences along with their meaning representations in SQL and the context free grammar. From the work of Kate and Mooney (2006), we have used the pseudo code shown in the Figure 2.8 for the KRISP training algorithm. As explained previously, KRISP follows an iterative learning process. The SVM classifiers learned in each iteration are an improvement over the classifiers learned in the previous iteration. Before starting the first iteration, for each production in the context free grammar, positive and negative examples are collected. The positive examples $P$ for a production $p$ are all those natural language sentences which use the production $p$ in the parse trees of their meaning representations. And the rest of sentences are taken as negative examples $N$ for that production. For example, the parse tree in Figure 2.6 uses the production $SB \rightarrow coach$. So, the natural language sentence "who is the coach for england cricket team." is considered as a positive example for the production $SB \rightarrow coach$. This complete sentence is taken as a positive example for all the productions used in the Figure 2.6.

Similarly, all the natural language sentences which does not use the production $SB \rightarrow coach$ in the parse of their meaning representation are considered as negative examples for $SB \rightarrow coach$. After collecting the positive and negative examples for each production as explained above, an SVM classifier $C$ is trained for each production using a String Subsequence Kernel. A String Subsequence Kernel is a function that estimates the similarity score between two strings. The similarity score depends on the number of subsequences those two strings share. A more detailed explana-

\[\text{http://www.csie.ntu.edu.tw/~cjlin/libsvm/}\]
Function TRAIN_KRISP( Training corpus \(\{(s_i, m_i)\}_{i=1..N}\), MRL grammar G)

// collect all the positive and negative examples for all productions
for each \(\pi \in G\)
  for \(i = 1\) to \(N\) do
    if \(\pi\) is used in parse\((m_i)\) then
      include \(s_i\) in \(P(\pi)\)
    else
      include \(s_i\) in \(N(\pi)\)
  
for iteration = 1 to MAX_ITER do
  for each \(\pi \in G\) do
    \(C_\pi = \text{train}\_\text{SVM}(P(\pi), N(\pi))\) // training SVM classifiers
  for each \(\pi \in G\)
    \(P(\pi) = \emptyset\) // remove the positive examples
  for \(i = 1\) to \(N\) do
    \(d^* = \text{Extended}\_\text{Earley}(s_i, G, P, m_i)\)
    // first call to the parser to get most probable correct tree
    // \(d^*\) is the most probable parse tree
    // \(m_i\) is the correct MR parse tree for the sentence \(s_i\)
    \(D = \text{Extended}\_\text{Earley}(s_i, G, P)\)
    // second call to get most probable parse trees
    // \(D\) represents \(x\) most probable parse trees (correct or incorrect)
    // Now, collect positive examples from most probable correct parse
    // tree obtained in the first call to Extended Earley's parser
    COLLECT_POSITIVES\((d^*)\)
    // collect negative examples from incorrect parse trees we got in second call
    // to Extended Earley's parser
    for each \(d \in D\) do
      if \(P(d) > P(d^*)\) and MR_parse\((d) \neq \text{parse}\((m_i)\))
        COLLECT_NEGATIVES\((d, d^*)\)
  
return classifiers \(C = \{C_\pi | \pi \in G\}\)

Figure 2.8: Pseudo code for KRISP training algorithm
tion about the String Subsequence Kernels and how similarity scores are estimated will be presented in the next chapter. If the two substrings share maximum possible subsequences, then the similarity score will be high and vice versa. The String Subsequence Kernels were previously used in natural language processing for text classification (Lodhi et al., 2002) and in relation information extraction (Bunescu and Mooney, 2005). The advantage in using String Subsequence Kernel is that they can use infinitely many subsequences as features.

The trained SVM classifiers decide the class to which a test data point (sentence) belongs. The probability of a production representing a substring is obtained by mapping the distance of the data point from the separating hyperplane constructed by SVM to a range of [0,1] using learned sigmoid function (Platt, 1999). The KRISP algorithm starts its first iteration by invoking the extended Earley’s parser two times for two purposes. In the first call, the extended Earley’s parser uses the trained SVM classifiers, context free grammar and a natural language sentence to give the $x$ most probable correct parse trees for a given sentence. Here, the extended Earley’s parser takes the correct parse tree for the given sentence as an additional parameter. Using this additional parameter, it finds the $x$ most probable parse trees. In this first call to the extend Earley’s parser, the emphasis is to find the most probable correct parse tree which need not be the most probable parse tree. The extended Earley’s parser does this by making sure the parse trees it derives contains only the subtrees of the correct parse tree. In the second call, the KRISP algorithm invokes the extended Earley’s parser without any additional parameter (without the correct parse for the given sentence) as in the first call. Here, the extended Earley’s parser uses the probability of the most probable correct parse tree returned in the first call as the threshold value. This will return $x$ most probable parse trees which need not necessarily be the correct parse trees. The emphasis in the second call is the find
the most probable parse trees which need not be the correct parse trees but have the
probability greater than the correct parse tree we got in the first call. After the first
iteration, the semantic parser may not give the correct parse trees for a given sentence.
Using the parse trees generated in the two calls to the extended Earley’s parser, the
KRISP algorithm collects more refined sets of positive and negative examples for the
next iteration.

The positive examples are collected from the most probable correct parse tree
returned on the first call to the extended Earley’s parser. Figure 2.6 is an example
for the correct parse tree for the sentence ”who is the coach for england cricket team.”
In the correct parse tree, the substring covered at each node is taken as a positive
element for the production at that node. For example, in the Figure 2.6, the substring
from the 3th word to the 6th word is taken as a positive example for production
$(_3WB \rightarrow B = WB \bullet_6)$. Negative examples are collected from the incorrect parse
trees returned when the extended Earley’s parser is called for the second time. The
parse trees returned on the second call have probability greater than that of the most
probable correct parse tree returned in the first call.

The KRISP collects the negative examples from the incorrect parse trees as de-
scribed below. The Figure 2.9 shows the incorrect parse tree. The procedure starts by traversing the correct and the incorrect parse trees simultaneously from the root using a breadth first traversal. The first node at which these two parse trees differ is found and all the words covered by the productions at these two nodes are noted. Figure 2.6 is the correct MR parse tree and the Figure 2.9 is the incorrect MR parse tree for the sentence "who is the coach for england cricket team." In the figures for correct and incorrect parse trees, they differ at the node \((WB \rightarrow B = WB)(3,6)\) and \((WB \rightarrow 'england')(5,6)\). The union of the words covered by both these production are from 3 to 6(coach for england). Now, it collects all the productions which cover any of these noted words (3 to 6) which are in the incorrect parse tree but not in the correct parse tree. For example, the production \(FB \rightarrow ODIPLAYERS (2,5)\) covers the noted words "coach" and "for" in the incorrect parse tree but in the correct parse tree it does not cover any of the noted words. In this case, the production \(FB \rightarrow ODIPLAYERS (2,5)\) is considered as a negative example, that is, the string "the coach for" is a negative example for the production \(FB \rightarrow ODIPLAYERS\). In the next iteration, the probability of this production representing the string "the coach for" is reduced. This reduces the overall probability of this incorrect tree to fall below the probability of the correct parse tree.

At the end of each iteration, the positive examples are erased and a new set of positive examples is collected. But the negative examples in each iteration are retained in order to get better accuracy. Also, we are considering only \(x\) most probable parse tree. We are missing a lot of incorrect parse trees. So retaining the negative examples across the iterations can train the system well and improve and accuracy. In order to further increase the number of negative examples, all the positive examples for a production \(P\) are taken as negative examples for the productions which have the same left-hand side non-terminal as \(P\). This avoids the incorrect parse trees generated
by use one production for the other which share the same left-hand side non-terminal.

Remember, the positive and negative examples collected now are more refined than those collected at the beginning of the first iteration. Initially, The production \( FB \rightarrow ODIPLAYERS \) considers the entire sentence "who is the coach for england cricket team." as a positive example. In the Figure 2.6 for the correct MR parse tree, the substring \( S[2,3] \), that is, the word "the" is taken as a positive example for the production \( FB \rightarrow ODIPLAYERS \) at the end of the first iteration. After collecting the positive and negative examples, the SVM classifiers are again trained. Using these newly trained SVM classifiers, the KRISP starts the second iteration and the entire processes is repeated again. For the subsequent iterations, it collects the positive and negative examples as explained above. The number of iterations is a system parameter. After the desired number of iterations, the SVM classifiers are returned. These classifiers are used on the test data set. The testing is done by invoking the extended Earley’s parser, which returns the \( x \) most probable parse trees. It takes the natural language sentence, the trained SVM classifiers and the meaning representational grammar as input.

As explained previously, KRISP does not learn classifiers for constant productions. In our work, the constant productions are differentiated from other productions by the usage of a dashed arrow in the grammar for representing them. For example, the production \( WB \rightarrow 'england’ \) is a constant substring. These productions represent a specific substring in the given natural language sentence. The parser looks for this substring in the given natural language sentence. Such productions are assigned a probability 1.
2.4 Data set

In this thesis, the meaning presentation used is Structured Query Language for the data set of queries concerning the game Cricket.

2.4.1 Cricket

Cricket is a game played in most of the countries around the globe. It is a game played between two teams with each team having 11 players. The basic idea of cricket is similar to a baseball game but the rules differ. It is usually played in an elliptical ground called the field. The most common terms used in cricket are bat, ball, stumps, fours, sixes, fifty, century, double century, wickets, bowler, batsman, all rounder, umpire, venues, man of the match, captain, coach, over etc. Some of the terms are self-explanatory.

In the center of the ground, there is 22 yard specially prepared pitch, which aligns along one of the axes of the ground. There are three wooden poles called stumps placed at both ends of the pitch. Cricket ball is a hard cork ball with leather covering and cricket bat is made of willow wood used to hit the ball. One team does the batting while the other does the bowling. An over is a set of six bowls bowled from one end of the pitch. The bowling team has to bowl a specified number of overs with changing ends. Each team gets to bat and bowl once. The target of the team batting first is to score as many runs as possible. The target of the team bowling first is to restrict the batting team’s score to a low as possible. A run is the unit of score. A batsman can score runs either by running from one end of the pitch to the other end or by hitting boundaries in fours or sixes. If the ball bounces at least once inside the field before it reaches the boundary of the field, then it is termed as four runs. If the ball reaches the boundary of the field without bouncing in the field, then it is called
a six. The most common way to get a batsman out is either by hitting the stumps or by catching the ball with in the field fully after the batsman hitting the ball with the bat. The target of the team which bowls second is to defend their score. The target of the team batting second is to score just more than the other team has scored. The player who performs well in all aspects of the game is awarded man of the match. A bowler is one who is good at bowling. A batsman is one who is good at batting and one who is good at both bowling and batting is called an all rounder.

2.4.2 Data Points

The data set for this thesis is formed from cricket. The data set consists of data points. Each data point is combination of a natural language sentence in English, its meaning representation in SQL and a correct parse tree for the SQL statement in context free grammar. The KRISP reads the data points and parse their corresponding SQL statements initially. The entire data set is then randomly divided into test set and train set. KRISP uses the training set for training the semantic parser. The data set includes more than one query of each type. For example, for the query "who is the coach for england cricket team," there is more than one query of similar structure with constants ('england' in this case) differing. This ensures the queries of similar structure are present in both training data set and test data set. As explained previously, different users can express the same query in more than one way. We have included all the possible ways a query can be expressed but with the same meaning representation and parse tree. For example, the above query can be expressed as "for england, who is the coach." Both these sentences differ in terms of the order of the words used but they have the same meaning. KRISP captures this by taking the permutations of the non-terminals on the right-hand side of the productions. Also
the sentences may differ in terms of the words used describe it but share the same meaning. For example, the above sentence can be stated as "who is the guide for england cricket team," "who is the mentor for england cricket team." We tried to include many possible ways a natural language sentence can be expressed.

The data set includes data points with varying complexity. The general structure of the SQL statement in context free grammar is represented as 

"select SB from FB where WB." The words represented in capitalized letter are non-terminals and the rest of the words are terminals. The non-terminals SB refers to the attributes the user wants to retrieve from a table. The non-terminal FB refers to the table from which the data is accessed. The non-terminal WB refers to the condition in the query. The example discussed in the above paragraph is a query selecting a single attribute "coach." In this case, \( SB \rightarrow coach \) is the production used in the parse tree.

A user can select multiple attributes from a table. For example, "what is the name and age of the coach for england cricket team." Its corresponding SQL statement would be "select name,age from teams where country = 'england'." In this example, the productions \((SB \rightarrow SB, SB)\), \((SB \rightarrow age)\) and \((SB \rightarrow coach)\) are the used in its parse tree. We have included queries that select multiple attributes from the table.

The non-terminal FB corresponds to the table from which the data is accessed. Again, we have included queries accessing data from single table and queries accessing data from multiple tables. For example, the query "select coach from teams where country = 'england'" access data from a single table "teams." In this case the production corresponding to FB is \( FB \rightarrow teams \). Here is an example of query in our data set which access two table:

A natural language sentence might be :"Obtain player name from australia who play both in odis and tests"
Its corresponding MR in SQL would be: "select names from odiplayers, testplayers
where odiplayers.nation = 'australia' and odiplayers.names = testplayer.names"

In the above example, the SQL query uses two tables odiplayers and testplayers
for finding data. In this case the productions corresponding to the non-terminal
FB would be (FB → FB , FB), (FB → odiplayers) and (FB → testplayers). We
have included queries operating on single table and multiple tables in the database.
The non-terminal WB is meant to specify the conditions in the query. We have in-
cluded queries that take a single condition and multiple conditions. For example, the
query "select coach from teams where country = 'england' " is an example
of a query that takes a single condition. The condition in this query is "country = 'england'". Its corresponding productions are (WB → B = WB), (B → country)
and (WB → 'england'). The example query shown above is an example of a query
having two conditions in it. The productions specifying multiple condition would be
something like, (WB → WB and WB), (WB → FB dot NATION = WB) and (WB
→ FB dot NATION = WB). The data set includes queries with different combina-
tions of the attributes, tables and conditions in them. For each English sentence, we
tried to include all the possible ways in which the same sentence can be expressed
which have the same meaning in SQL. For example,

who is the coach for canada cricket team
list the coach for australia cricket team
who is the trainer for india cricket team
who is the guide for australia cricket team
find the coach for australia cricket team
so on..
All the above English sentences represent different ways of expressing the same question. They all share a common meaning representation in SQL and parse tree. We have included at least 8 data points of each type of sentence in the data set, which differ by the constant string in them. Below is the set of sentences for the first type of above sentences which differ by constant terms.

- who is the coach for canada cricket team
- who is the coach for india cricket team
- who is the coach for australia cricket team
- who is the coach for bangladesh cricket team
- who is the coach for pakistan cricket team
- who is the coach for usa cricket team
- who is the coach for canada cricket team
- who is the coach for srilanka cricket team

We have presented a more detailed explanation in the next chapter about the type of queries included in our data set.
3 Implementation

This chapter presents a detailed explanation about the tasks we have implemented as part of developing a complete semantic parser. In the first section we discuss the string subsequence kernel used in this work. In the next section, we explain about the type of queries used in the data set. We have used the Java programming language to implement the system.

3.1 String Subsequence Kernel (Lodhi et al., 2002)

As explained in the previous chapter, in our thesis work, we use a string subsequence kernel function as the kernel function in SVM. A string subsequence kernel estimates the similarity scores between natural language sentences. The similarity score between two sentences depends up on the number of subsequences the two sentences share.

From the framework of Lodhi et al. (2002), a string kernel between two strings is defined as the number of subsequences they share. The subsequences Lodhi et al. used in their framework are characters, while in this work our subsequences are words in the strings. The similarity score is high when the two sentences share a good number of subsequences and vice versa.

From the formal notation of Rousu and Shawe-Taylor (2005), let $\Sigma$ be a finite alphabet, a string is a finite sequence of elements from $\Sigma$, and the set of all strings is denoted by $\Sigma^*$. We denote the length of a string $s$ by $|s|$ and the string is represented
as \( s = s_1 s_2 s_3 \ldots s_{|s|} \). Here \( 1, 2, 3, \ldots |s| \) represent the indices of the words in the string \( s \). For example, "number_1 of_2 fours_3 hit_4 by_5 sachie_6" is a string with indices starting from 1 to 6 where the word of index 1 is "number," 2 is "of," etc. A substring of the string \( s \) is represented as \( s[i..j] \), where \( i \) and \( j \) are the indices of the words in the sentences. The words in a substring of a sentence are contiguous, that is, the word indices in a substring form a contiguous sequence of numbers. A string \( u \) is said to be a subsequence of the string \( s \) if there exists a sequence of word indices \( j= (j_1 j_2 j_3 \ldots j_{|u|}) \) in \( u \) such that \( 1 \leq j_1 < j_2 < \ldots < j_{|u|} \leq |s| \). It is represented as \( u = s[j] \). A subsequence of a sentence need not form a contiguous set of words of a string. For example, the string "fours_3 by_5 sachie_6" is a subsequence of the string "number_1 of_2 fours_3 hit_4 by_5 sachie_6." The word index sequence for the subsequence is \( j = (3 5 6) \) and it follows the above condition, \( 1 \leq 3 < 5 < 6 \leq 6 \). The word index sequence of this subsequence is not contiguous. The distance between the first index and the last index of a subsequence is called the span of the subsequence, \( \text{span}(j) \). Formally it is given by, \( \text{span}(j) = j_{|u|} - j_1 + 1 \). The span of the subsequence "fours_3 by_5 sachie_6" is, \( \text{span}(j) = 6 - 3 + 1 = 4 \).

There can be multiple subsequences with a unique set of word indices \( j \) for a given sentence. We define \( \Phi_u(s) \) as number of such subsequences possible for a given sentence, that is, \( \Phi_u(s) = |\{j \mid s[j] = u\}| \) where \( u \) and \( s[j] \) denote a subsequence of the sentence \( s \). This definition of finding the count of subsequences does not consider the gaps that may be present in the word indices of a subsequence. As explained in the above paragraph, a subsequence need not form a contiguous set of words in the sentence. For example, the subsequence "fours_3 by_5 sachie_6" has a gap between the words "fours_3" and "by_5." The word "hit" is missing in the subsequence, which makes it a non-contiguous sequence. We redefine the above definition as Equation 3.1 by introducing a gap penalty factor \( \lambda \). When estimating the similarity scores, the gap penalty has to be considered if the subsequence does not form a contiguous
In the above equation, \( \lambda \) is the gap penalty factor which takes a value in \((0,1]\). It is a system parameter. In Equation \(3.1\), the factor \(\lambda^{|u|}\) in the denominator is the normalization factor. It ensures that the gap penalty is applied to only the subsequences which have gaps in their index sequence. If \( \lambda \) is 1, that is, when there is no gap penalty in the subsequence, then the above formula is same as the original one.

We define the kernel \( K(s,t) \) between two string \( s \) and \( t \) as

\[
K(s,t) = \sum_{u \in \Sigma^*} \Phi_u(s) \Phi_u(t) \tag{3.2}
\]

The kernel defined by Equation \(3.2\) considers all the subsequences between two string \( s \) and \( t \) to compute the similarity score between them. The kernel \( K(s,t) \) represents the similarity score between the string \( s \) and \( t \).

Consider two sentences "number_1 of_2 fours_3 hit_4 by_5 the_6 batsman_7" and "the_1 fours_2 by_3 batsman_4." Let these sentences be \( s \) and \( t \). These two sentences have multiple subsequences in common. The table below shows all the subsequences that are shared between these two sentences. The words that are common in both the sentences are "fours," "by," "the," and "batsman."

Table 3.1 is an example showing how the subsequence kernel is computed for two strings \( s \) and \( t \). The first column in the table shows the list of all the subsequences that are shared by the two string \( s \): "number_1 of_2 fours_3 hit_4 by_5 the_6 batsman_7" and \( t \): "the_1 fours_2 by_3 batsman_4." All subsequences that are common between the two strings are considered for computing the kernel. Subsequences of all possible lengths are taken into consideration. It starts with the possible common string subsequences of length one, two and so on. The subscript numbers in strings \( s \) and \( t \) are the indices of the words in both the strings. The second column in the table 3.1 shows the indices.
of the words in the corresponding subsequence of the sentence $s$. For example, $\{(3, 5), 3\}$ in the second column indicates that the subsequence "fours by" has indices (3 5) in the string $s$ and its span is 5-3+1, that is, 3. The third column in the table is the list of indices along with the span of the corresponding subsequences of the sentence $t$. The fourth column computes the number of such unique index sequences in the string $s$ using the equation $\Phi_u(s) = 1/|u| \sum_{i:s[i]=u} \lambda^{\text{span}(i)}$.

| Subsequences (u) | $\{(j, \text{span}(j)) | s[j] = u\}$ | $\{(j, \text{span}(j)) | t[j] = u\}$ | $\Phi_u(s)$ | $\Phi_u(t)$ | $\Phi_u(s) \ast \Phi_u(s)$ |
|------------------|-------------------------------------|-------------------------------------|-------------|-------------|------------------|
| fours            | $\{(3, 1)\}$                        | $\{(2, 1)\}$                       | 1           | 1           | 1                |
| by               | $\{(5, 1)\}$                        | $\{(3, 1)\}$                       | 1           | 1           | 1                |
| the              | $\{(6, 1)\}$                        | $\{(1, 1)\}$                       | 1           | 1           | 1                |
| batsman          | $\{(7, 1)\}$                        | $\{(4, 1)\}$                       | 1           | 1           | 1                |
| fours by         | $\{(3 5, 3)\}$                      | $\{(2 3, 2)\}$                     | $\lambda$  | 1           | $\lambda$       |
| fours batsman    | $\{(3 7, 5)\}$                      | $\{(2 4, 3)\}$                     | $\lambda^3$| $\lambda$  | $\lambda^4$     |
| by batsman       | $\{(5 7, 3)\}$                      | $\{(3 4, 2)\}$                     | $\lambda$  | 1           | $\lambda$       |
| the batsman      | $\{(6 7, 2)\}$                      | $\{(1 4, 4)\}$                     | 1           | $\lambda^2$| $\lambda^2$     |
| fours by batsman| $\{(3 5 7, 5)\}$                    | $\{(2 3 4, 3)\}$                   | $\lambda^2$| 1           | $\lambda^2$     |

$K(s,t) = 4 + 2\lambda + 2\lambda^2 + \lambda^4$

Table 3.1: An example for computing subsequence kernel

The symbol $\lambda$ denotes the gap penalty in the subsequence. The Equation 3.1 penalize a subsequence if there is a gap in its index sequence. For example, the subsequence "fours by" in the sentence $s$ has the indices (3 5). The index 4 is missing in the subsequence. So the Equation 3.1 for this particular subsequence would be

$$\Phi_u(s) = \lambda^{\text{span}(u)} / \lambda^{|u|}.$$
The length, \(|u|\) of the subsequence "fours by" is 2 and its span is 3.

So, \(\Phi_u(s) = \lambda^3 / \lambda^2 = \lambda\)

If the subsequence does not have any gap between its index sequence, then \(\Phi_u(s)\) would be 1. For example, the subsequence "the batsman" in the sentence \(s\) does not have any gap in its index sequence. For such subsequence, \(|u| = \text{span}(u)\). So, \(\Phi_u(s) = 1\). Column five of table 3.1 shows the number of such unique index sequences in the string \(t\) using Equation 3.1. The last column in table 3.1 shows the product of corresponding entries in column four and five. The final kernel between the two strings \(s\) and \(t\) is then computed using the Equation 3.2, which is the sum of all the entries in the last column of the table 3.1.

The kernel is then normalized to have the value in \([0,1]\) by using the equation,

\[
K_{\text{norm}}(s,t) = K(s,t) / \sqrt{K(s,s) \times K(t,t)}
\]

The kernel is normalized to remove any bias that might be introduced due to the lengths of the strings.

### 3.1.1 Probability Estimation

All possible common subsequences between the strings are taken as features in our kernel. The kernel implicitly uses the feature space to compute the dot product. KRISP invokes an extended Earley’s parser by passing it a natural language sentence, a set of classifiers and the context free grammar. The extended Earley’s parser starts parsing the given sentence from the left to right. It associates each substring of the sentence with all productions in the grammar. It invokes the classifier associated with a production in order to find the probability of that production representing a particular substring. The classifier uses a string subsequence kernel to find the similarity score. For example, when a substring "fours by" is associated with a production
SB → odiplayers, the parser invokes the trained classifier $C$ for this production. The SVM classifier $C$ uses the string subsequence kernel to compute the similarity scores between the string "fours by" and every string on which the classifier $C$ is trained. So, if the classifier $C$ is trained on ten strings previously, then the similarity score is computed between the string "fours by" and the ten strings on which $C$ is trained. These similarity scores are computed as explained in the previous section.

The similarity scores computed form a feature vector $V$, which is used in computing the dot product. Each dimension in the feature vector $V$ represents a similarity score computed between the string "fours by" and the strings on which the classifier $C$ is trained previously. The SVM uses a decision function to decide which side of the hyperplane the string "fours by" belongs to. Remember, this is a two-class classification. The string may be a positive example or a negative example for a given classifier. The positive examples and negative examples are separated by a hyperplane. Figure 2.4 shows an example of a representation of a hyperplane separating two classes of data. The decision function uses the sum of the dot products between the feature vector $V$ and all the support vectors to decide which side of the hyperplane the string "fours by" belongs to. The sign of the sum of the dot product decides whether the string "fours by" belongs to the set of positive examples or negative examples. After mapping the string on one side of the hyperplane, its the distance from the hyperplane is used to estimate the probability of the associated production representing that particular string. In our example, the mapping distance of the data point "fours by" from the hyperplane is used to compute the probability of the production $SB → odiplayers$ representing the string "fours by." If the probability is greater than the threshold value, it is retained. If the probability is below the threshold value, it is ignored. A detailed explanation is presented in the previous chapter about how the cumulative probability of the entire parse tree is computed.
3.2 Data Set

As explained in the previous chapter, the data set for the system is collected from a game called cricket. The cricketing terminology is introduced in the previous chapter. The system reads the data set from a file and randomly divides the entire data set into a test set and a train set. The classifiers are trained using the training data set and the system is tested on the test data set. The data set includes queries of varying complexity.

The data set includes queries retrieving data from a single table and from multiple tables. The following are few English sentences with their corresponding SQL statements that retrieve data from a single table:

Who is the coach for canada cricket team?
select coach from nations where team = 'canada'

Number of fours hit by gambhir in odis?
select fours from odiplayers where names = 'gambhir'

Total count of matches played at mcg?
select played from venues where ground = 'mcg'

What are the names of westindies players who can just bowl?
select names from odiplayers where country = 'westindies' and role = 'bowl'

What are the names of westindies players who can just bat?
select names from odiplayers where country = 'westindies' and role = 'bat'

Name teams that participate in superseries?
select teams from odistats where tour_name = 'superseries'

All test tournaments where pakistan was victorious?
select tour_name from teststats where teststats . champions = 'pakistan'

All odi clashes where england was winner?
select tour_name from odistats where odistats . champions = 'england'

All of the above listed queries retrieve data from a single table. For example, the query corresponding to the natural language sentence "Total count of matches played at mcg?" is "select played from venues where ground = 'mcg'." The SQL statement retrieves the attribute "played" from the table "venues" in the database. The natural language sentence corresponding to the above SQL statement is flexible in terms of the order of the words or the type of words. There can be more than one natural language sentence that expresses the same meaning or users may type different English sentences to search the same entity in the system. In our data set, we have included many possible natural language sentences for the same meaning representation, that is, in SQL. For example, the following are the list of some of the natural language sentences for the SQL statement mentioned above.

Give count for matches played at mcg.
Display count of matches hosted at mcg.
Display the count of matches played at mcg.

How many matches played at mcg?

Get count of matches played at mcg.

What is the count of matches played at mcg?

Number of matches played at mcg?

For all the above listed natural languages sentences, the corresponding SQL statement is "select played from venues where ground = 'mcg'." For each sentence in the above set, we have included at least six sentences that have a similar structure but with a different constant term in them. For example, for the query "What is the count of matches played at mcg," the constant term in it is the word "mcg." We have included the following set of queries having similar structure.

What is the count of matches played at mcg?

What is the count of matches played at buffalo?

What is the count of matches played at adeleide?

What is the count of matches played at vacca?

What is the count of matches played at uppal?

What is the count of matches played at feroshahkotla?

What is the count of matches played at eadengarden?

What is the count of matches played at chinnaswamy?

As explained previously, the data set is randomly divided into test set and the train set. We expect that the above set of queries are present in both the train set and the test set so that the system can be tested on what it has learned. The queries in the above set differ in the constant terms. The following are the set of productions
for the above set of queries that retrieve single attribute accessing a single table. An SQL statement would be "select played from venues where ground = 'mcg'." The constant term in this query is 'mcg,' but it could be any of the names depending on the query.

In the general structure of the SQL query, the production with non-terminal $SB$ on the left hand side refers to the attribute(s) selected by the query. In the above set of queries, it refers to the attribute 'played.' The production with non-terminal $FB$ on the left hand side refers to the FROM clause, that is, the table(s) from which the data is retrieved. The production with non-terminal $WB$ on the left hand side refers to the WHERE clause, that is, the condition(s) specified in the query. In the above set of productions, the production $SB \rightarrow \text{played}$ refers to the attribute 'played' of the SQL statement. The production $FB \rightarrow \text{venues}$ refers to the table 'venues' of
the database. The production $WB \rightarrow B = STADIUM$ refers to the condition in the SQL statement. In this production, the non-terminals on the right hand side ’B’ and ’STADIUM’ are part of the where clause. All the productions are given id numbers. As explained in the previous chapter, a data point is a combination of the English sentence, its corresponding SQL statement and SQL parse tree in the context free grammar. All the data points are listed in one file and the all the productions are listed in another file. An SQL parse tree is represented by a sequence of numbers that indicates the production number in the file containing all the productions. For example, the following is the representation of two data points.

1) What is the count of matches played at mcg?

   "select played from venues where ground = 'mcg' "
   0 1 2 3 4 5

2) What is the count of matches played at mcg?

   "select played from venues where ground = 'buffalo' "
   0 1 2 3 4 6

Figure 3.1 is the parse tree for the SQL statement in the first data point. The number at each node is the production number. In a data point, the first sentence is a query in natural language sentence, the second statement is its corresponding SQL statement and the third sentence is the sequence of numbers representing the parse of the SQL statement. The numbers indicate the productions in the SQL parse. The parse starts with the zeroth production $Q \rightarrow \text{return} \text{select} SB \text{from} FB \text{where} WB$. The capitalized words represent the non-terminals.

All of the above queries retrieve a single attribute from a table. We have also included queries that retrieve multiple attributes from the tables. An SQL statement
which retrieves multiple attributes from a table would be something like "select experience, country from the nations where coach = 'fletcher'." The English query for this SQL statement might be "What is the experience and country for which fletcher is working as the coach." This query retrieves two attributes, 'experience' and 'country.' This query shares almost the same structure as the above set of queries but differs in the number of attributes. So, the production corresponding the attributes would change to $SB \rightarrow SB, SB$. The non-terminals on the right hand side refer to the attributes of the query. In this example, the two non-terminals on the right hand side would refer to the attributes 'experience' and 'country' respectively.

We have included queries that retrieve data from multiple tables. For example, the following is an example of a data point in which the SQL statement retrieves data
Natural language (NL):

Obtain player’s name from australia who play both in odis and tests.

SQL:

```
select names from odiplayers, testplayers where odiplayers . nation = 'australia' and odiplayers . names = odiplayers . names
```

For such queries, the production corresponding to the FROM clause would be $FB \rightarrow FB, FB$. The two non-terminals on the right hand side of the production refer to the two tables in database. In the above example, the two tables are "odiplayer" and "testplayers." It can be generalized to access more than two tables. The above query is also an example of a query which has more than one condition in the WHERE clause. The production corresponding to the WHERE clause would be $WB \rightarrow WB and WB$. The two non-terminals on the right hand side of the production refer to two conditions in the SQL statement. In the above example, the two conditions are "odiplayers . nation = 'australia' " and "odiplayers . names = odiplayers . names." Again, for all these data points we have included at least six queries with the same structure and also considered all the possible ways in which the same natural language sentence of a data point can be expressed. The following are some of the English sentences along with their corresponding SQL queries in our data set which have multiple conditions in them.

NL: Obtain player’s name from india who play both in odis and tests

SQL: select names from odiplayers, testplayers where odiplayers . nation = 'india' and odiplayers . names = testplayers . names
What are the names of Pakistan players who can just bowl?

SQL: `select names from odiplayers where country = 'pakistan' and role = 'bowl'`

Which ODI team won Bordergavaskartrophy in the year 2006?

SQL: `select champions from odiStats where tour_name = 'bordergavaskartrophy' and year = '2006'`

Which test team won Royalseries in the year 2006?

SQL: `select champions from testStats where tour_name = 'royalseries' and year = '2006'`

Show ODI mom for first match of Bordergavaskartrophy in year 2006

SQL: `select mom from odiStats where tour_name = 'bordergavaskartrophy' and year = '2006' and matchnumber = 'first'`

What are the names of Pakistan players who can just bat?

SQL: `select names from odiplayers where country = 'pakistan' and role = 'bat'`

What are the names of India players who can bat and bowl?

SQL: `select names from odiplayers where country = 'india' and role = 'alrounder'`
The following is the list of productions for queries accessing two tables and queries which specify two condition:

\[
\begin{align*}
0 \ Q & \rightarrow \ \text{select } SB \ \text{from } FB \ \text{where } WB \\
1 \ SB & \rightarrow \ \text{names} \\
2 \ FB & \rightarrow \ FB, \ FB \\
3 \ FB & \rightarrow \ \text{odiplayers} \\
4 \ FB & \rightarrow \ \text{testplayers} \\
5 \ WB & \rightarrow \ WB \ and \ WB \\
6 \ WB & \rightarrow \ FB . \ B = \text{NATION} \\
7 \ WB & \rightarrow \ FB . \ SB = FB . \ SB \\
8 \ B & \rightarrow \ \text{nation} \\
9 \ NATION & \rightarrow \ 'india'
\end{align*}
\]

Using the above list of productions, the parse sequence for the SQL query "\textbf{select} names from odiplayers, testplayers where odiplayers. nation = 'india' and odiplayers. names = testplayers. names" would be

\[0 \ 1 \ 2 \ 3 \ 4 \ 5 \ 6 \ 3 \ 8 \ 9 \ 7 \ 3 \ 1 \ 4 \ 1\]

\textbf{Note:} The period(.) and equals(=) symbols are not non-terminals.

The system first reads the two files, one containing the data set and other containing the list of productions used in the parse of the SQL statements in the data points. After reading the data set, all the data points are randomly divided into to train set and test set. The total number of data points in our data set is 262. All the data points in the data set are numbered from 0 to 261. We have included six queries of each type. The six queries for each type of query would have a similar structure.
They differ in terms of the constants used in those queries. The reason for including six queries of each type is to make sure that the test set and the train contains at least one query of a each type. This would make the training and testing easier. For example, if three queries of a particular type are in test set and the remaining three queries of that type are in train set, then the system can learn the structure of that particular type of query using the three queries in train set and it can be tested on the other three queries in the test set. If the number of queries for a particular type are low, then the chances of having the queries of a particular type in both test set and train set are low. For example, consider only two queries of each type. In the worst case, if those two queries are in test set and no query of that type in the train set, then the system would predict them incorrectly as the train set does not include any query of that type. As the number of queries of a particular type is high, then the chances of having the queries of a particular type in both train set and test set are high. We have considered to go with six queries of each type.

We have not included any duplicate data points in our data set. We have used libsvm.jar for the SVM support. Libsvm (Chang et al., 2011) is a library for SVM (Boser et al., 1992). It supports multi-class classification as well. We integrated the libsvm.jar library into our system. The libsvm only deals with numerical data. So, the data representation has to be converted into numerical format before it is fed to the libsvm library. That is the reason why we have used numerical representations for productions and index sequences for strings.
4 Results

This chapter presents the experiments we have conducted using our new data set on semantic parsing system to test its accuracy.

4.1 Experiments

As explained in the previous chapters, the system divides the entire data set into a test set and training set. The train data set is used for learning the semantic parsing system and the test data set is used to test the accuracy of the system trained on the train data set.

4.1.1 N-fold Cross-Validation Experiment 1

We have conducted N-fold cross-validation experiments to test the accuracy of the semantic parsing system using our new data set. The variable $N$ in an N-fold cross-validation experiment is a system parameter. We repeat an experiment $N$ times and each such repetition is termed as fold in an N-fold cross-validation experiment.

In each fold of an N-fold cross-validation experiment, the entire data set is randomly divided into a test set and training set. The system is trained on training set and tested on test set. The system is trained for $K$ iterations before it is tested on the test data set. The variable $K$ is a system parameter. Before starting the first iteration, the system collects positive and negative examples for each production. Each iteration involves two major steps. In the first step, all the SVM classifiers
corresponding to all their productions are trained on their respective sets of positive
and negative examples collected. In the second step, the system is trained on the
training data set to collect positive and negative examples for the next iteration. The
procedure to collect the positive and negative examples is explained in the second
chapter. If the system is tested at this point, it may not predict the correct MR
parse trees for all the natural language sentences in the test set. This concludes the
first iteration. The system then starts the second iteration and repeats the above two
steps. In the first step, it trains all the SVM classifiers with their respective sets of
positive and negative examples collected in the second step of the previous iteration.
These classifiers are more refined than the previous set of classifiers used in the first
iteration. These newly trained classifiers are used by the system in second step to
collect positive and negative examples for the next iteration. The two steps explained
above are repeated in each iteration on the same training set. After the $K^{th}$ iteration,
the final set of SVM classifiers are returned. These final set of SVM classifiers are used
by the system to test on the test data set. The results obtained at this point gives
the accuracy of the system. A fold in an N-fold cross-validation experiment concludes
with testing the system on test data set using the final set of SVM classifiers returned
after the $K^{th}$. In the next fold, the data set is again randomly divided into training
set and test set. The data set is divided in such a way that, the test set in each fold
of an experiment is unique, that is, the data points which are part of the test set in
one fold would not be part of a test set in other fold. We repeat the same procedure
for $N$ times in an N-fold cross-validation experiment and record the accuracy of the
system obtained in each fold. The Figure 4.1 below, shows the training and test sets
in each fold of a 3-fold cross-validation experiment on a data set consisting of 20 data
points.

The example shown in Figure 4.1 considers a data set with twenty data points.
Figure 4.1: Training and test sets in each fold of a 3-fold cross-validation experiment.

The data points are numbered from 0 to 19 as shown in the first array of the Figure 4.1. The system reads all the data points from one file and all the productions from another file. In the example shown in Figure 4.1, let us assume we train the system for 3 iterations in each fold. According to our experiment, initially, all the data points are shuffled using a random number generator as shown in the second array of numbers.
In Figure 4.1. In the first fold of the 3-fold cross-validation experiment, the data set is first divided into training set and test set. The test data set would contain the data points numbered 2,15,12,13,8,7 and 17 as shown the figure. The training set holds the remaining data points. In each iteration, the two steps that are explained in the above paragraph are executed. The third iteration in this example concludes the first fold of our 3-fold experiment. After the third iteration, we test the system on the test data set and record the accuracy of the system. In the next fold, we create a new training data set and test data set. According to our implementation, the test set in the second fold would contain the data points numbered 18,19,1,0,10,4 and 11 as shown in the figure. The training set holds the remaining data points. Again, the system is trained for three iteration and then tested. In the last fold, the test set would contain the data points 5,16,9,6,3 and 14 as shown in the figure. If the test sets in all the three folds are combined, we get the entire data set. The accuracy of the system is recorded in each fold. The table below shows the results obtained after testing the system on our entire data set using a 10-fold cross-validation experiment with 3 iterations in each fold. We have recorded the accuracy of the system after each iteration in each fold of the N-fold cross-validation experiment. The first row in the table below shows the folds and the remaining rows show the percentage accuracy obtained after each iteration of a fold. Figure 4.2 shows the accuracy of the system at each iteration of a 5-fold cross-validation experiment.

<table>
<thead>
<tr>
<th>Fold</th>
<th>1</th>
<th>2</th>
<th>3</th>
<th>4</th>
<th>5</th>
<th>6</th>
<th>7</th>
<th>8</th>
<th>9</th>
<th>10</th>
</tr>
</thead>
<tbody>
<tr>
<td>Iteration 1</td>
<td>0%</td>
<td>0%</td>
<td>0%</td>
<td>0%</td>
<td>0%</td>
<td>0%</td>
<td>0%</td>
<td>0%</td>
<td>0%</td>
<td>0%</td>
</tr>
<tr>
<td>Iteration 2</td>
<td>74%</td>
<td>90%</td>
<td>62%</td>
<td>85%</td>
<td>90%</td>
<td>87%</td>
<td>88%</td>
<td>86%</td>
<td>90%</td>
<td>81%</td>
</tr>
<tr>
<td>Iteration 3</td>
<td>92%</td>
<td>91%</td>
<td>90%</td>
<td>92%</td>
<td>91%</td>
<td>95%</td>
<td>89%</td>
<td>91%</td>
<td>94%</td>
<td>93%</td>
</tr>
</tbody>
</table>
4.1.2 N-fold Cross-Validation Experiment 2

In this experiment, we tried to draw a relationship between the size of the training set used for training the system and the accuracy of that system. This experiment is almost same as the N-fold cross-validation experiment described in the previous section but with some minor changes. In N-fold cross-validation experiment described in the previous section, we have created new training and test data sets in each fold. But, in this experiment, we create a training set and test set only once at the start of the N-fold cross-validation experiment. The test set is fixed for the entire experiment but the training set in each fold is incrementally selected from the original training set.
set that is created at the start of this experiment. We have adopted the following procedure to conduct this experiment.

First, we have divided the entire data set into test set $T_t$ and the training set $T_r$. The test set is fixed for all the folds in the N-fold cross-validation experiment. For the first fold of this experiment, the training set includes only 10% of the original training set ($T_r$). So, for the first fold, the training set is 10% of original training set $T_r$ and test set is the original test set $T_t$. We have set the number of iteration to 5. After the fifth iteration, the system is tested on the test set $T_t$ and its accuracy is recorded. For the next fold, the training set would be 20% of the original training set $T_r$ and the test set would be the same test set $T_t$. The training set in each fold includes the training set of the previous folds. For example, the 20% of training set ($T_r$) used in the second fold includes the 10% of training set ($T_t$) used in the first fold. The same procedure is repeated until we use 100% of the original training set $T_r$, that is, we perform a 10 fold experiment by incrementing the size of the training set by 10% in each fold but with a fixed test set $T_t$. The accuracy of the system is recorded in each fold. This experiment shows dependency of the accuracy of the system on the amount of training data used for the system.

We have repeated the above experiment 10 times. Figure 4.3 shows the graph that is plotted between the amount of training data and accuracy of the system. We considered the mean of the values gained in all 10 experiments to plot the graph.

The curve in the graph 4.3 indicates that the accuracy of the system increases with the amount of training data used for the system. Figure 4.4 shows the same graph with error bars.

The Tables 4.1 and 4.2 below show the accuracy of the system with the percentage of the training data used in each experiment. The data in the first column in both the tables is the percentage of training data used in each fold of an experiment. The
data in the remaining columns in both the tables is the percentage accuracy obtained in each fold of the corresponding experiments. In each experiment we have set the number of iteration to 5.
Figure 4.4: Accuracy vs %Training Data with error bars
The experimental results show that the accuracy of the semantic parsing system improves as the amount of the training data grows.
5 Conclusions

5.1 Contribution

In this thesis we have tested a semantic parsing learner which is based on the system, Kernel-based Robust Interpretation for Semantic Parsing, KRISP algorithm (Kate and Mooney, 2006) on a new data set. The data set we have used is created from the game called Cricket. Our data set includes queries of varying complexity. The system uses the training data set for learning the semantic parser and it is tested on the test data set. We have conducted the two experiments for testing the accuracy of the semantic parsing system. In the first N-fold cross-validation experiments, the results have shown that the accuracy of the system improves with the number of iteration used for learning the system. For example, if the system is trained for just one iteration and tested, the results obtained may not be good. It may not predict correct meaning representation parse trees for all the sentences in the test data set. But if the system is trained for 4 or 5 iterations and then tested, the results obtained are better than the results obtained in the previous case. The system may predict correct meaning representation parse trees for most of sentences in the test data set. In this case, the system learns from the positive and negative examples collected in each iteration and generates more refined SVM classifiers in each iteration.

The second N-fold cross-validation experiment concludes that the accuracy of the semantic parsing system depends on the amount of training data used with a fixed test data set. The accuracy of the system grows as the amount of training data used
5.2 Future Work

5.2.1 Extending to Other Domains

In this thesis, we have tested the semantic parsing system on the data set related to a specific domain. The semantic parsing system can be extended to other domains as well. For example, it can be tested on movie databases or geographical databases. The data sets would contain the data specific to those domains on which it is tested. It can be extended to open domains as well. Developing a semantic parser which can work with a data set from an open domain would be a challenging task for the future research. Such system can be used as language translation system or a better question and answering system. The difficult task is to come up with a single meaning representational language for an open domain data. The meaning representation can be narrowed down to a set of rules by interpreting the actions the semantic parser is supposed to take in response to a given natural language sentence. The semantic parsing learning system may have to do some preprocessing using some shallow analysis techniques like word sense disambiguation (Ide and Veronis, 1998) in order to handle the ambiguities in the natural language sentences in the training data.

5.2.2 Testing With Other Semantic Parsing Methods

We have used Support Vector Machine classifiers with String Subsequence Kernel to implement the semantic parsing learner. Some earlier learning systems (Zelle and Mooney (1996); Tang and Mooney (2001); Kate et al. (2005)) used rule-based
learning methods for semantic parsing. Recent learning systems for semantic parsing (Ge et al. (2005); Zettlemoyer and Collins (2005); Wong et al. (2006)) have used statistical feature-based methods. It would be interesting to test our Cricket data set on these systems. The results obtained would give a good comparison between the rule-based methods, feature-based methods and the kernel-based methods used for implementing a semantic parsing learner.

5.2.3 Data Collection

The data set we have collected is a limited set. We have collected the data set by assuming several ways a natural language sentence could be expressed. There are efficient ways of collecting the data set. One better way to collect a data set is by providing a user-interface where a user can type and submit a natural language sentence, possibly with an answer. This would cover a broad range of queries and this would also give the more possible ways a natural language sentence can be expressed by different users. When users from different domains submit their queries, a lot of ambiguous queries can be covered. This would make a broader and challenging data set for the system.

5.2.4 Testing on Complex queries

The data set we have collected includes queries of moderate complexity. We have not considered complex queries which includes nested queries and queries which use predefined functions. It would be interesting to see how the system performs when it is tested on more complex and lengthy queries. As the length of a query grows, the time taken by the semantic parser to parse the sentence would increase as well. For a nested query, the productions used to specify the meaning representation parse tree
would be more complex when compared to a simple query.
6 Bibliography


7 Appendix

7.1 Productions

Each production is given a number, which is used to specify the parse tree for an SQL query.

Q

0 Q → select ?SB from ?FB where ?WB

1 SB → fours

2 SB → coach

3 SB → played

4 B → names

5 B → team

6 B → ground

7 FB → ODIPLAYERS

8 FB → NATIONS

9 FB → VENUES

10 WB → ?B ?EQ ?PLAYER


13 EQ → =

14 PLAYER → ' sachin ' 

15 PLAYER → ' gambhir '
<table>
<thead>
<tr>
<th>Player</th>
<th>Country</th>
<th>Venue</th>
</tr>
</thead>
<tbody>
<tr>
<td>yadav</td>
<td>canada</td>
<td>mcg</td>
</tr>
<tr>
<td>sehag</td>
<td>pakistan</td>
<td>chinnaswamy</td>
</tr>
<tr>
<td>rohit</td>
<td>india</td>
<td>buffalo</td>
</tr>
<tr>
<td>kohli</td>
<td>southafrica</td>
<td>adeleide</td>
</tr>
<tr>
<td>raina</td>
<td>australia</td>
<td>lords</td>
</tr>
<tr>
<td>dhoni</td>
<td>newzeland</td>
<td>feroshahkotla</td>
</tr>
<tr>
<td>praveen</td>
<td>bangladesh</td>
<td>edengarden</td>
</tr>
<tr>
<td>yuvi</td>
<td>srilanka</td>
<td>trentbridge</td>
</tr>
</tbody>
</table>
42 VENUE --> ' uppal '
43 VENUE --> ' vacca '
44 WB -> ?W1 and ?W2
45 R1 --> ' bowl '
46 R -> COUNTRY
47 SB -> names
48 W1 --> ?R ?EQ ?NATION
49 W2 --> ?SPC ?EQ ?R1
50 SPC -> ROLE
51 R1 --> ' bat '
52 VENUE --> ' vcc '
53 B -> nation
54 FB -> ODIPLAYER , ?F2
55 F2 -> TESTPLAYERS
56 W1 --> ?FB . ?B = ?NATION
57 W2 --> ?FB . names = ?F2 . names
58 SB -> teams
59 B -> tour_name
60 FB -> ODISTATS
62 TOUR_NAME --> ' bordergavaskartrophy '  
63 TOUR_NAME --> ' worldcup '  
64 TOUR_NAME --> ' championstrophy '  
65 TOUR_NAME --> ' unitycup '  
66 TOUR_NAME --> ' superseries '  
67 TOUR_NAME --> ' dlfcup '
68 TOUR_NAME → 'miniworldcup'
69 TOUR_NAME → 'sharjahcup'
70 TOUR_NAME → 'asiacup'
71 TOUR_NAME → 'ipl'
72 SB → tour_name
73 B → champions
74 WB → ?FB dot ?B = ?NATION
75 FB → TESTSTATS
76 SB → champions
77 W1 → ?B ?EQ ?TOUR_NAME
79 B → year
80 TOUR_NAME → 'bordergavaskartrophy2'
81 TOUR_NAME → 'worldcup2'
82 TOUR_NAME → 'championstrophy2'
83 TOUR_NAME → 'unitycup2'
84 TOUR_NAME → 'superseries2'
85 TOUR_NAME → 'dlfcup2'
86 TOUR_NAME → 'miniworldcup2'
87 TOUR_NAME → 'sharjahcup2'
88 TOUR_NAME → 'asiacup2'
89 TOUR_NAME → 'ipl2'
90 TOUR_NAME → 'ashes'
91 TOUR_NAME → 'royalseries'
92 TOUR_NAME → 'natwestseries'
93 TOUR_NAME → 'hondacup'
7.2 Data Points

There are 262 data points described below. Each data point has a number (first line), its English sentence, the corresponding query in SQL and two lines showing parses of the SQL query. The sequence of numbers in the last two lines of data point represent the corresponding productions listed above. The parse tree for an SQL query can be obtained from these sequence of numbers. For example, for the first data point defined below, the parse tree for its SQL query can be obtained from the sequence 0 2 8 11 5 13 24. These numbers represent their corresponding productions in the above list. The parse tree starts with the first number 0 which represents the production $Q \rightarrow \text{select } ?SB \text{ from } ?FB \text{ where } ?WB$. The next number in the sequence is 2 which represents the production $SB \rightarrow \text{coach}$. The number 8 represents
the production $FB \rightarrow NATIONS$, the number 11 represents the production $WB \rightarrow \text{?B ?EQ ?NATION}$, the number 5 represents the production $B \rightarrow \text{team}$, the number 13 represents the production $EQ \rightarrow =$ and the number 24 represents the production $NATION \rightarrow 'canada'$. All these productions represent the parse tree for the SQL query “select coach from NATIONS where team = ‘canada’.”

262
0
who is the coach for canada cricket team

select coach from NATIONS where team = 'canada'
0 2 8 11 5 13 24
0 2 8 11 5 13 24
1
who is the coach for pakistan cricket team

select coach from NATIONS where team = 'pakistan'
0 2 8 11 5 13 25
0 2 8 11 5 13 25
2
Obtain player’s name from india who play both in odis and tests

select names from ODIPLAYERS , TESTPLAYERS where ODIPLAYERS . nation = 'india' and ODIPLAYERS . names = TESTPLAYERS . names
0 47 54 55 44 56 7 53 26 57 7 55
0 47 54 55 44 56 7 53 26 57 7 55
3
Obtain player’s name from pakistan who play both in odis and tests
select names from ODIPLAYERS, TESTPLAYERS where ODIPLAYERS.nation = 'pakistan' and ODIPLAYERS.names = TESTPLAYERS.names

Obtain player’s name from srilanka who play both in odis and tests
select names from ODIPLAYERS, TESTPLAYERS where ODIPLAYERS.nation = 'srilanka' and ODIPLAYERS.names = TESTPLAYERS.names

what are the names of pakistan players who can just bowl
select names from ODIPLAYERS where COUNTRY = 'pakistan' and ROLE = 'bowl'

what are the names of srilanka players who can just bowl
select names from ODIPLAYERS where COUNTRY = 'srilanka' and ROLE = 'bowl'

what are the names of australia players who can just bowl
select names from ODIPLAYERS where COUNTRY = 'australia' and ROLE = 'bowl'
who is the coach for india cricket team
select coach from NATIONS where team = 'india'
0 2 8 11 5 13 26
0 2 8 11 5 13 26
9

who is the coach for southafrica cricket team
select coach from NATIONS where team = 'southafrica'
0 2 8 11 5 13 27
0 2 8 11 5 13 27
10

who is the coach for australia cricket team
select coach from NATIONS where team = 'australia'
0 2 8 11 5 13 28
0 2 8 11 5 13 28
11

list the coach for australia cricket team
select coach from NATIONS where team = 'australia'
0 2 8 11 5 13 28
0 2 8 11 5 13 28
12

list the coach for india cricket team
select coach from NATIONS where team = 'india'
0 2 8 11 5 13 26
0 2 8 11 5 13 26
13

who is the trainer for canada cricket team
select coach from NATIONS where team = 'canada'
0 2 8 11 5 13 24
0 2 8 11 5 13 24
14
who is the trainer for pakistan cricket team
select coach from NATIONS where team = 'pakistan'
0 2 8 11 5 13 25
0 2 8 11 5 13 25
15
who is the trainer for india cricket team
select coach from NATIONS where team = 'india'
0 2 8 11 5 13 26
0 2 8 11 5 13 26
16
who is the guide for india cricket team
select coach from NATIONS where team = 'india'
0 2 8 11 5 13 26
0 2 8 11 5 13 26
17
who is the guide for canada cricket team
select coach from NATIONS where team = 'canada'
0 2 8 11 5 13 24
0 2 8 11 5 13 24
18
who is the guide for canada cricket team
select coach from NATIONS where team = 'canada'
who is the guide for australia cricket team
select coach from NATIONS where team = 'australia'

Number of fours hit by sachin in odis
select fours from ODIPLAYERS where names = 'sachin'

Number of fours hit by gambhir in odis
select fours from ODIPLAYERS where names = 'gambhir'

Number of fours hit by yadav in odis
select fours from ODIPLAYERS where names = 'yadav'

Number of fours hit by sehwag in odis
select fours from ODIPLAYERS where names = 'sehwag'
0 1 7 10 4 13 17
24
Number of fours hit by yadav in odis
select fours from ODIPLAYERS where names = 'yadav'
0 1 7 10 4 13 16
0 1 7 10 4 13 16
25
how many fours hit by sachin in odis
select fours from ODIPLAYERS where names = 'sachin'
0 1 7 10 4 13 14
0 1 7 10 4 13 14
26
how many fours hit by gambhir in odis
select fours from ODIPLAYERS where names = 'gambhir'
0 1 7 10 4 13 15
0 1 7 10 4 13 15
27
how many fours hit by yadav in odis
select fours from ODIPLAYERS where names = 'yadav'
0 1 7 10 4 13 16
0 1 7 10 4 13 16
28
how many fours hit by sehwag in odis
select fours from ODIPLAYERS where names = 'sehwag'
0 1 7 10 4 13 17
0 1 7 10 4 13 17

84
29

how many fours hit by sehwag in odis

select fours from ODIPLAYERS where names = 'sehwag'

0 1 7 10 4 13 17
0 1 7 10 4 13 17

30

Number of boundaries hit by sachin in odis

select fours from ODIPLAYERS where names = 'sachin'

0 1 7 10 4 13 14
0 1 7 10 4 13 14

31

Number of boundaries hit by gambhir in odis

select fours from ODIPLAYERS where names = 'gambhir'

0 1 7 10 4 13 15
0 1 7 10 4 13 15

32

Number of boundaries hit by yadav in odis

select fours from ODIPLAYERS where names = 'yadav'

0 1 7 10 4 13 16
0 1 7 10 4 13 16

33

Number of boundaries hit by sehwag in odis

select fours from ODIPLAYERS where names = 'sehwag'

0 1 7 10 4 13 17
0 1 7 10 4 13 17

34
Number of boundaries hit by Sehwag in ODIs

select fours from ODIPLAYERS where names = 'sehwag'

0 1 7 10 4 13 17
0 1 7 10 4 13 17
35

give total fours hit by Sachin in ODIs

select fours from ODIPLAYERS where names = 'sachin'

0 1 7 10 4 13 14
0 1 7 10 4 13 14
0 1 7 10 4 13 14 36

36

give total fours hit by Sachin in ODIs

select fours from ODIPLAYERS where names = 'sachin'

0 1 7 10 4 13 14
0 1 7 10 4 13 14
0 1 7 10 4 13 14 36

37

give total fours hit by Gambhir in ODIs

select fours from ODIPLAYERS where names = 'gambhir'

0 1 7 10 4 13 15
0 1 7 10 4 13 15
0 1 7 10 4 13 15 38

38

give total fours hit by Yadav in ODIs

select fours from ODIPLAYERS where names = 'yadav'

0 1 7 10 4 13 16
0 1 7 10 4 13 16
0 1 7 10 4 13 16 39

39

give total fours hit by Sehwag in ODIs

86
select fours from ODIPLAYERS where names = 'sehwag'
0 1 7 10 4 13 17
0 1 7 10 4 13 17
40
Give count for matches played at mcg
select played from VENUES where ground = 'mcg'
0 3 9 12 6 13 34
0 3 9 12 6 13 34
41
Give count for matches played at chinnaswamy
select played from VENUES where ground = 'chinnaswamy'
0 3 9 12 6 13 35
0 3 9 12 6 13 35
42
Give count for matches played at buffalo
select played from VENUES where ground = 'buffalo'
0 3 9 12 6 13 36
0 3 9 12 6 13 36
43
Give count for matches played at adeleide
select played from VENUES where ground = 'adeleide'
0 3 9 12 6 13 37
0 3 9 12 6 13 37
44
Give count for matches played at adeleide
select played from VENUES where ground = 'adeleide'

87
0 3 9 12 6 13 37

45

Total count of matches played at mcg

select played from VENUES where ground = 'mcg'

0 3 9 12 6 13 34

46

Total count of matches played at mcg

select played from VENUES where ground = 'mcg'

0 3 9 12 6 13 34

47

Total count of matches played at chinnaswamy

select played from VENUES where ground = 'chinnaswamy'

0 3 9 12 6 13 35

48

Total count of matches played at buffalo

select played from VENUES where ground = 'buffalo'

0 3 9 12 6 13 36

49

Total count of matches played at adeleide

select played from VENUES where ground = 'adeleide'

0 3 9 12 6 13 37
display count of matches hosted at mcg
select played from VENUES where ground = 'mcg'

display count of matches hosted at mcg
select played from VENUES where ground = 'mcg'

display count of matches hosted at chinnaswamy
select played from VENUES where ground = 'chinnaswamy'

display count of matches hosted at buffalo
select played from VENUES where ground = 'buffalo'

display count of matches hosted at adelaide
select played from VENUES where ground = 'adelaide'
55
get count of matches played at mcg
select played from VENUES where ground = 'mcg'
0 3 9 12 6 13 34
0 3 9 12 6 13 34
56
get count of matches played at mcg
select played from VENUES where ground = 'mcg'
0 3 9 12 6 13 34
0 3 9 12 6 13 34
57
get count of matches played at chinnaswamy
select played from VENUES where ground = 'chinnaswamy'
0 3 9 12 6 13 35
0 3 9 12 6 13 35
58
get count of matches played at buffalo
select played from VENUES where ground = 'buffalo'
0 3 9 12 6 13 36
0 3 9 12 6 13 36
59
get count of matches played at adeleide
select played from VENUES where ground = 'adeleide'
0 3 9 12 6 13 37
0 3 9 12 6 13 37
60
list the coach for canada cricket team
select coach from NATIONS where team = 'canada'
0 2 8 11 5 13 24
0 2 8 11 5 13 24
61
list the coach for pakistan cricket team
select coach from NATIONS where team = 'pakistan'
0 2 8 11 5 13 25
0 2 8 11 5 13 25
62
list the coach for southafrica cricket team
select coach from NATIONS where team = 'southafrica'
0 2 8 11 5 13 27
0 2 8 11 5 13 27
63
what are the names of westindies players who can just bowl
select names from ODIPLAYERS where COUNTRY = 'westindies' and ROLE = 'bowl'
0 47 7 44 48 46 13 32 49 50 13 45
0 47 7 44 48 46 13 32 49 50 13 45
64
what are the names of westindies players who can just bowl
select names from ODIPLAYERS where COUNTRY = 'westindies' and ROLE = 'bowl'
0 47 7 44 48 46 13 32 49 50 13 45
0 47 7 44 48 46 13 32 49 50 13 45
find only names of pakistan players who can bowl
select names from ODIPLAYERS where COUNTRY = 'pakistan' and ROLE = 'bowl'
0 47 7 44 48 46 13 25 49 50 13 45
0 47 7 44 48 46 13 25 49 50 13 45

find only names of srilanka players who can bowl
select names from ODIPLAYERS where COUNTRY = 'srilanka' and ROLE = 'bowl'
0 47 7 44 48 46 13 31 49 50 13 45
0 47 7 44 48 46 13 31 49 50 13 45

find only names of australia players who can bowl
select names from ODIPLAYERS where COUNTRY = 'australia' and ROLE = 'bowl'
0 47 7 44 48 46 13 28 49 50 13 45
0 47 7 44 48 46 13 28 49 50 13 45

find only names of westindies players who can bowl
select names from ODIPLAYERS where COUNTRY = 'westindies' and ROLE = 'bowl'
0 47 7 44 48 46 13 32 49 50 13 45
0 47 7 44 48 46 13 32 49 50 13 45

find only names of westindies players who can bowl
select names from ODIPLAYERS where COUNTRY = 'westindies' and ROLE = 'bowl'
0 47 7 44 48 46 13 32 49 50 13 45
show only names of pakistan players who can bowl
select names from ODIPLAYERS where COUNTRY = 'pakistan' and ROLE = 'bowl'

show only names of srilanka players who can bowl
select names from ODIPLAYERS where COUNTRY = 'srilanka' and ROLE = 'bowl'

show only names of australia players who can bowl
select names from ODIPLAYERS where COUNTRY = 'australia' and ROLE = 'bowl'

show only names of westindies players who can bowl
select names from ODIPLAYERS where COUNTRY = 'westindies' and ROLE = 'bowl'

show only names of westindies players who can bowl
select names from ODIPLAYERS where COUNTRY = 'westindies' and ROLE = 'bowl'
Display only the names of Pakistan players who can just bowl
select names from ODIPLAYERS where COUNTRY = 'pakistan' and ROLE = 'bowl'

Display only the names of Sri Lanka players who can just bowl
select names from ODIPLAYERS where COUNTRY = 'srilanka' and ROLE = 'bowl'

Display only the names of Australia players who can just bowl
select names from ODIPLAYERS where COUNTRY = 'australia' and ROLE = 'bowl'

Display only the names of West Indies players who can just bowl
select names from ODIPLAYERS where COUNTRY = 'westindies' and ROLE = 'bowl'
'bowl'  
0 47 7 44 48 46 13 32 49 50 13 45  
0 47 7 44 48 46 13 32 49 50 13 45  
80  
Name teams that participate in bordergavaskartrophy  
select teams from ODISTATS where tour.name = 'bordergavaskartrophy'  
0 58 60 61 59 13 62  
0 58 60 61 59 13 62  
81  
Name teams that participate in worldcup  
select teams from ODISTATS where tour.name = 'worldcup'  
0 58 60 61 59 13 63  
0 58 60 61 59 13 63  
82  
Name teams that participate in championstrophy  
select teams from ODISTATS where tour.name = 'championstrophy'  
0 58 60 61 59 13 64  
0 58 60 61 59 13 64  
83  
Name teams that participate in unitycup  
select teams from ODISTATS where tour.name = 'unitycup'  
0 58 60 61 59 13 65  
0 58 60 61 59 13 65  
84  
Name teams that participate in superseries  
select teams from ODISTATS where tour.name = 'superseries'
List out teams participating in bordergavaskartrophy
select teams from ODISTATS where tour_name = 'bordergavaskartrophy'

List out teams participating in worldcup
select teams from ODISTATS where tour_name = 'worldcup'

List out teams participating in championstrophy
select teams from ODISTATS where tour_name = 'championstrophy'

List out teams participating in unitycup
select teams from ODISTATS where tour_name = 'unitycup'

List out teams participating in superseries
select teams from ODISTATS where tour_name = 'superseries'
which teams are playing in bordergavaskartrophy
select teams from ODIStats where tour_name = 'bordergavaskartrophy'

which teams are playing in worldcup
select teams from ODIStats where tour_name = 'worldcup'

which teams are playing in championstrophy
select teams from ODIStats where tour_name = 'championstrophy'

which teams are playing in unitycup
select teams from ODIStats where tour_name = 'unitycup'

which teams are playing in superseries
select teams from ODIStats where tour_name = 'superseries'
List out nations that compete in bordergavaskartrophy
select teams from ODISTATS where tour_name = 'bordergavaskartrophy'

List out nations that compete in worldcup
select teams from ODISTATS where tour_name = 'worldcup'

List out nations that compete in championstrophy
select teams from ODISTATS where tour_name = 'championstrophy'

List out nations that compete in unitycup
select teams from ODISTATS where tour_name = 'unitycup'

List out nations that compete in superseries
select teams from ODISTATS where tour_name = 'superseries'
all odi series where India was victorious

select tour_name from ODISTATS where ODISTATS dot champions = 'india'
0 72 60 74 60 73 26
0 72 60 74 60 73 26
101

all odi series where Pakistan was victorious

select tour_name from ODISTATS where ODISTATS dot champions = 'pakistan'
0 72 60 74 60 73 25
0 72 60 74 60 73 25
102

all odi series where Srilanka was victorious

select tour_name from ODISTATS where ODISTATS dot champions = 'srilanka'
0 72 60 74 60 73 31
0 72 60 74 60 73 31
103

all odi series where Australia was victorious

select tour_name from ODISTATS where ODISTATS dot champions = 'australia'
0 72 60 74 60 73 28
0 72 60 74 60 73 28
104

all odi series where Westindies was victorious

select tour_name from ODISTATS where ODISTATS dot champions = 'westindies'
0 72 60 74 60 73 32
0 72 60 74 60 73 32
105

all odi clashes where Southafrica was winner

99
<table>
<thead>
<tr>
<th>Page</th>
<th>Query</th>
<th>Details</th>
</tr>
</thead>
<tbody>
<tr>
<td>106</td>
<td><code>select tour_name from ODISTATS where ODISTATS.dot.champions = 'southafrica';</code></td>
<td>072 60 74 60 73 27</td>
</tr>
<tr>
<td>107</td>
<td><code>select tour_name from ODISTATS where ODISTATS.dot.champions = 'newzeland';</code></td>
<td>072 60 74 60 73 29</td>
</tr>
<tr>
<td>108</td>
<td><code>select tour_name from ODISTATS where ODISTATS.dot.champions = 'bangladesh';</code></td>
<td>072 60 74 60 73 30</td>
</tr>
<tr>
<td>109</td>
<td><code>select tour_name from ODISTATS where ODISTATS.dot.champions = 'canada';</code></td>
<td>072 60 74 60 73 24</td>
</tr>
<tr>
<td>110</td>
<td><code>select tour_name from ODISTATS where ODISTATS.dot.champions = 'england';</code></td>
<td>072 60 74 60 73 33</td>
</tr>
<tr>
<td></td>
<td>all odi clashes where newzeland was winner</td>
<td>all odi clashes where bangladesh was winner</td>
</tr>
<tr>
<td></td>
<td><code>select tour_name from ODISTATS where ODISTATS.dot.champions = 'india';</code></td>
<td>all odi clashes where canada was winner</td>
</tr>
<tr>
<td></td>
<td><code>select tour_name from ODISTATS where ODISTATS.dot.champions = 'england';</code></td>
<td>all test tournaments where india was victorious</td>
</tr>
<tr>
<td></td>
<td><code>select tour_name from TESTSTATS where TESTSTATS.dot.champions = 'india';</code></td>
<td></td>
</tr>
</tbody>
</table>
all test tournaments where pakistan was victorious
select tour_name from TESTSTATS where TESTSTATS dot champions = 'pakistan'

all test tournaments where srilanka was victorious
select tour_name from TESTSTATS where TESTSTATS dot champions = 'srilanka'

all test tournaments where australia was victorious
select tour_name from TESTSTATS where TESTSTATS dot champions = 'australia'

all test tournaments where westindies was victorious
select tour_name from TESTSTATS where TESTSTATS dot champions = 'westindies'

all test cups where southafrica was winner
select tour_name from TESTSTATS where TESTSTATS dot champions = 'southafrica'
all test cups where New Zealand was winner

```
select tour_name from TESTSTATS where TESTSTATS.dot champions = 'New Zealand'
```

all test cups where Bangladesh was winner

```
select tour_name from TESTSTATS where TESTSTATS.dot champions = 'Bangladesh'
```

all test cups where Canada was winner

```
select tour_name from TESTSTATS where TESTSTATS.dot champions = 'Canada'
```

all test cups where England was winner

```
select tour_name from TESTSTATS where TESTSTATS.dot champions = 'England'
```

Which ODI team won Border Gavaskar Trophy in the year 2006

```
select champions from ODI STATS where tour_name = 'Border Gavaskar Trophy' and year = '2006'
```
Which ODI team won Worldcup in the year 2006

```
select champions from ODIstats where tour_name = 'worldcup' and year = '2006'
```

Which ODI team won Championstrophy in the year 2006

```
select champions from ODIstats where tour_name = 'championstrophy' and year = '2006'
```

Which ODI team won Unitycup in the year 2006

```
select champions from ODIstats where tour_name = 'unitycup' and year = '2006'
```

Which ODI team won Superseries in the year 2006

```
select champions from ODIstats where tour_name = 'superseries' and year = '2006'
```

Which ODI team got DLFCup in the year 2006

```
select champions from ODIstats where tour_name = 'dlfcup' and year = '2006'
```
Which odi team got miniworldcup in the year 2006
select champions from ODISTATS where tour_name = 'miniworldcup' and year = '2006'

Which odi team got sharjahcup in the year 2006
select champions from ODISTATS where tour_name = 'sharjahcup' and year = '2006'

Which odi team got asiacup in the year 2006
select champions from ODISTATS where tour_name = 'asiacup' and year = '2006'

Which odi team got ipl in the year 2006
select champions from ODISTATS where tour_name = 'ipl' and year = '2006'

Which test team won ashes in the year 2006
select champions from TESTSTATS where tour_name = 'ashes' and year = '2006'
Which test team won natwestseries in the year 2006
select champions from TESTSTATS where tour_name = 'natwestseries' and year = '2006'

Which test team won royalseries in the year 2006
select champions from TESTSTATS where tour_name = 'royalseries' and year = '2006'

Which test team won hondacup in the year 2006
select champions from TESTSTATS where tour_name = 'hondacup' and year = '2006'

Which test team won ashes in the year 2001
select champions from TESTSTATS where tour_name = 'ashes' and year = '2001'

Which test team got ashes in the year 2007
select champions from TESTSTATS where tour_name = 'ashes' and year = '2007'
Which test team got ashes in the year 2008
select champions from TESTSTATS where tour_name = 'ashes' and year = '2008'

Which test team got ashes in the year 2009
select champions from TESTSTATS where tour_name = 'ashes' and year = '2009'

Which test team got ashes in the year 2010
select champions from TESTSTATS where tour_name = 'ashes' and year = '2010'

Which test team got ashes in the year 2011
select champions from TESTSTATS where tour_name = 'ashes' and year = '2011'

Show odi mom for first match of bordergavaskartrophy in year 2006
select mom from ODISTATS where tour_name = 'bordergavaskartrophy' and year = '2006' and matchnumber = ' first '
Show odi mom for first match of worldcup in year 2006
select mom from ODISTATS where tour_name = 'worldcup' and year = '2006' and
matchnumber = 'first'

Show odi mom for first match of championstrophy in year 2006
select mom from ODISTATS where tour_name = 'championstrophy' and year = '2006'
and matchnumber = 'first'

Show odi mom for first match of unitycup in year 2006
select mom from ODISTATS where tour_name = 'unitycup' and year = '2006' and
matchnumber = 'first'

Show odi mom for first match of superseries in year 2006
select mom from ODISTATS where tour_name = 'superseries' and year = '2006' and
matchnumber = 'first'
Find odi mom for first match of dlfcup of year 2006
select mom from ODISTATS where tour_name = 'dlfcup' and year = '2006' and
matchnumber = ' first '
0 104 60 105 77 59 13 67 78 79 13 94
0 104 60 105 77 59 13 67 78 79 13 94

Find odi mom for first match of miniworldcup of year 2006
select mom from ODISTATS where tour_name = 'miniworldcup' and year = '2006'
and matchnumber = ' first '
0 104 60 105 77 59 13 68 78 79 13 94
0 104 60 105 77 59 13 68 78 79 13 94

Find odi mom for first match of sharjahcup of year 2006
select mom from ODISTATS where tour_name = 'sharjahcup' and year = '2006' and
matchnumber = ' first '
0 104 60 105 77 59 13 69 78 79 13 94
0 104 60 105 77 59 13 69 78 79 13 94

Find odi mom for first match of asiacup of year 2006
select mom from ODISTATS where tour_name = 'asiacup' and year = '2006' and
matchnumber = ' first '
0 104 60 105 77 59 13 70 78 79 13 94
0 104 60 105 77 59 13 70 78 79 13 94

Find odi mom for first match of ipl of year 2006
select mom from ODI_STATS where tour_name = 'ipl' and year = '2006' and match-number = 'first'

Show test mom for first match of ashes in year 2006
select mom from TESTSTATS where tour_name = 'ashes' and year = '2006' and matchnumber = 'first'

Show test mom for first match of natwestseries in year 2006
select mom from TESTSTATS where tour_name = 'natwestseries' and year = '2006' and matchnumber = 'first'

Show test mom for first match of royalseries in year 2006
select mom from TESTSTATS where tour_name = 'royalseries' and year = '2006' and matchnumber = 'first'

Show test mom for first match of hondacup in year 2006
select mom from TESTSTATS where tour_name = 'hondacup' and year = '2006' and matchnumber = 'first'
Show test mom for first match of ashes in year 2006
select mom from TESTSTATS where tour_name = 'ashes' and year = '2006' and matchnumber = 'first'

Find test mom for first match of ashes of year 2007
select mom from TESTSTATS where tour_name = 'ashes' and year = '2007' and matchnumber = 'first'

Find test mom for first match of ashes of year 2008
select mom from TESTSTATS where tour_name = 'ashes' and year = '2008' and matchnumber = 'first'

Find test mom for first match of ashes of year 2009
select mom from TESTSTATS where tour_name = 'ashes' and year = '2009' and matchnumber = 'first'
Find test mom for first match of ashes of year 2010

```
select mom from TESTSTATS where tour_name = 'ashes' and year = '2010' and matchnumber = 'first'
```

0 104 75 105 77 59 13 90 78 79 13 102
0 104 75 105 77 59 13 90 78 79 13 102

Find test mom for first match of ashes of year 2011

```
select mom from TESTSTATS where tour_name = 'ashes' and year = '2011' and matchnumber = 'first'
```

0 104 75 105 77 59 13 90 78 79 13 103
0 104 75 105 77 59 13 90 78 79 13 103

who is the trainer for australia cricket team

```
select coach from NATIONS where team = 'australia'
```

0 2 8 11 5 13 28
0 2 8 11 5 13 28

who is the trainer for southafrica cricket team

```
select coach from NATIONS where team = 'southafrica'
```

0 2 8 11 5 13 27
0 2 8 11 5 13 27

who is the guide for canada cricket team

```
select coach from NATIONS where team = 'canada'
```

0 2 8 11 5 13 24
Obtain player’s name from australia who play both in odis and tests
select names from ODIPLAYERS, TESTPLAYERS where ODIPLAYERS . nation = 'australia' and ODIPLAYERS . names = TESTPLAYERS . names

Obtain player’s name from westindies who are both in odis and tests
select names from ODIPLAYERS, TESTPLAYERS where ODIPLAYERS . nation = 'westindies' and ODIPLAYERS . names = TESTPLAYERS . names

Obtain player’s name from southafrica who are both in odis and tests
select names from ODIPLAYERS, TESTPLAYERS where ODIPLAYERS . nation = 'southafrica' and ODIPLAYERS . names = TESTPLAYERS . names

Obtain player’s name from newzeland who are both in odis and tests
select names from ODIPLAYERS, TESTPLAYERS where ODIPLAYERS . nation = 'newzeland' and ODIPLAYERS . names = TESTPLAYERS . names
Obtain player’s name from bangladesh who are both in odis and tests
select names from ODIPLAYERS , TESTPLAYERS where ODIPLAYERS . nation = 'bangladesh' and ODIPLAYERS . names = TESTPLAYERS . names

Obtain player’s name from canada who play are in odis and tests
select names from ODIPLAYERS , TESTPLAYERS where ODIPLAYERS . nation = 'canada' and ODIPLAYERS . names = TESTPLAYERS . names

Obtain player’s name from england who play are in odis and tests
select names from ODIPLAYERS , TESTPLAYERS where ODIPLAYERS . nation = 'england' and ODIPLAYERS . names = TESTPLAYERS . names

what are the names of pakistan players who can just bat
select names from ODIPLAYERS where COUNTRY = 'pakistan' and ROLE = 'bat'

what are the names of srilanka players who can just bat
select names from ODIPLAYERS where COUNTRY = 'srilanka' and ROLE = 'bat'
what are the names of australia players who can just bat
select names from ODIPLAYERS where COUNTRY = 'australia' and ROLE = 'bat'

what are the names of westindies players who can just bat
select names from ODIPLAYERS where COUNTRY = 'westindies' and ROLE = 'bat'

find only names of southafrica players who can bat
select names from ODIPLAYERS where COUNTRY = 'southafrica' and ROLE = 'bat'

find only names of bangladesh players who can bat
select names from ODIPLAYERS where COUNTRY = 'bangladesh' and ROLE = 'bat'

find only names of newzeland players who can bat
select names from ODIPLAYERS where COUNTRY = 'newzeland' and ROLE = 'bat'
0 47 7 44 48 46 13 29 49 50 13 51
0 47 7 44 48 46 13 29 49 50 13 51

177

find only names of canada players who can bat
select names from ODIPLAYERS where COUNTRY = 'canada' and ROLE = 'bat'
0 47 7 44 48 46 13 24 49 50 13 51
0 47 7 44 48 46 13 24 49 50 13 51

178

find only names of england players who can bat
select names from ODIPLAYERS where COUNTRY = 'england' and ROLE = 'bat'
0 47 7 44 48 46 13 33 49 50 13 51
0 47 7 44 48 46 13 33 49 50 13 51

179

find only names of india players who can bat
select names from ODIPLAYERS where COUNTRY = 'india' and ROLE = 'bat'
0 47 7 44 48 46 13 26 49 50 13 51
0 47 7 44 48 46 13 26 49 50 13 51

180

show only names of bangladesh players who can bat
select names from ODIPLAYERS where COUNTRY = 'bangladesh' and ROLE = 'bat'
0 47 7 44 48 46 13 30 49 50 13 51
0 47 7 44 48 46 13 30 49 50 13 51

181
show only names of newzeland players who can bat
select names from ODIPLAYERS where COUNTRY = 'newzeland' and ROLE = 'bat'
0 47 7 44 48 46 13 29 49 50 13 51
0 47 7 44 48 46 13 29 49 50 13 51

show only names of canada players who can bat
select names from ODIPLAYERS where COUNTRY = 'canada' and ROLE = 'bat'
0 47 7 44 48 46 13 24 49 50 13 51
0 47 7 44 48 46 13 24 49 50 13 51

show only names of england players who can bat
select names from ODIPLAYERS where COUNTRY = 'england' and ROLE = 'bat'
0 47 7 44 48 46 13 33 49 50 13 51
0 47 7 44 48 46 13 33 49 50 13 51

show only names of india players who can bat
select names from ODIPLAYERS where COUNTRY = 'india' and ROLE = 'bat'
0 47 7 44 48 46 13 26 49 50 13 51
0 47 7 44 48 46 13 26 49 50 13 51

Display only the names of bangladesh players who can just bat
select names from ODIPLAYERS where COUNTRY = 'bangladesh' and ROLE = 'bat'
0 47 7 44 48 46 13 30 49 50 13 51
0 47 7 44 48 46 13 30 49 50 13 51
Display only the names of New Zealand players who can just bat
select names from ODIPLAYERS where COUNTRY = 'New Zealand' and ROLE = 'bat'

Display only the names of Canada players who can just bat
select names from ODIPLAYERS where COUNTRY = 'Canada' and ROLE = 'bat'

Display only the names of England players who can just bat
select names from ODIPLAYERS where COUNTRY = 'England' and ROLE = 'bat'

Display only the names of India players who can just bat
select names from ODIPLAYERS where COUNTRY = 'India' and ROLE = 'bat'

who is the coach for Canada cricket team
select coach from NATIONS where team = 'Canada'
who is the coach for canada cricket team

```sql
select coach from NATIONS where team = 'canada'
```

0 2 8 11 5 13 24
0 2 8 11 5 13 24

list the coach for canada cricket team

```sql
select coach from NATIONS where team = 'canada'
```

0 2 8 11 5 13 24
0 2 8 11 5 13 24

list the coach for canada cricket team

```sql
select coach from NATIONS where team = 'canada'
```

0 2 8 11 5 13 24
0 2 8 11 5 13 24

list the coach for canada cricket team

```sql
select coach from NATIONS where team = 'canada'
```

0 2 8 11 5 13 24
0 2 8 11 5 13 24

list the coach for canada cricket team

```sql
select coach from NATIONS where team = 'canada'
```

0 2 8 11 5 13 24
0 2 8 11 5 13 24
who is the trainer for canada cricket team
select coach from NATIONS where team = 'canada'
0 2 8 11 5 13 24
0 2 8 11 5 13 24
197

who is the trainer for canada cricket team
select coach from NATIONS where team = 'canada'
0 2 8 11 5 13 24
0 2 8 11 5 13 24
198

who is the trainer for canada cricket team
select coach from NATIONS where team = 'canada'
0 2 8 11 5 13 24
0 2 8 11 5 13 24
199

who is the guide for canada cricket team
select coach from NATIONS where team = 'canada'
0 2 8 11 5 13 24
0 2 8 11 5 13 24
200

who is the guide for canada cricket team
select coach from NATIONS where team = 'canada'
0 2 8 11 5 13 24
0 2 8 11 5 13 24
201

who is the guide for canada cricket team
select coach from NATIONS where team = 'canada'
0 2 8 11 5 13 24
0 2 8 11 5 13 24
202

Number of fours hit by yadav in odis
select fours from ODIPLAYERS where names = 'yadav'
0 1 7 10 4 13 16
0 1 7 10 4 13 16
203

Number of fours hit by sehwag in odis
select fours from ODIPLAYERS where names = 'sehwag'
0 1 7 10 4 13 17
0 1 7 10 4 13 17
24

Number of fours hit by yadav in odis
select fours from ODIPLAYERS where names = 'yadav'
0 1 7 10 4 13 16
0 1 7 10 4 13 16
204

how many fours hit by sachin in odis
select fours from ODIPLAYERS where names = 'sachin'
0 1 7 10 4 13 14
0 1 7 10 4 13 14
205

how many fours hit by gambhir in odis
select fours from ODIPLAYERS where names = 'gambhir'
120
how many fours hit by yadav in odis
select fours from ODIPLAYERS where names = 'yadav'

Number of boundaries hit by sachin in odis
select fours from ODIPLAYERS where names = 'sachin'

Number of boundaries hit by gambhir in odis
select fours from ODIPLAYERS where names = 'gambhir'

Number of boundaries hit by yadav in odis
select fours from ODIPLAYERS where names = 'yadav'

give total fours hit by sachin in odis
select fours from ODIPLAYERS where names = 'sachin'
give total fours hit by sachin in odis
select fours from ODIPLAYERS where names = 'sachin'

give total fours hit by gambhir in odis
select fours from ODIPLAYERS where names = 'gambhir'

Give count for matches played at mcg
select played from VENUES where ground = 'mcg'

Give count for matches played at chinnaswamy
select played from VENUES where ground = 'chinnaswamy'

Give count for matches played at buffalo
select played from VENUES where ground = 'buffalo'
<p>| | | | | | | |</p>
<table>
<thead>
<tr>
<th></th>
<th></th>
<th></th>
<th></th>
<th></th>
<th></th>
<th></th>
</tr>
</thead>
<tbody>
<tr>
<td>0</td>
<td>3</td>
<td>9</td>
<td>12</td>
<td>6</td>
<td>13</td>
<td>34</td>
</tr>
<tr>
<td>0</td>
<td>3</td>
<td>9</td>
<td>12</td>
<td>6</td>
<td>13</td>
<td>34</td>
</tr>
</tbody>
</table>

Total count of matches played at mcg

select played from VENUES where ground = 'mcg'

0 3 9 12 6 13 34
0 3 9 12 6 13 34

Total count of matches played at chinnaswamy

select played from VENUES where ground = 'chinnaswamy'

0 3 9 12 6 13 35
0 3 9 12 6 13 35

display count of matches hosted at mcg

select played from VENUES where ground = 'mcg'

0 3 9 12 6 13 34
0 3 9 12 6 13 34

display count of matches hosted at chinnaswamy

select played from VENUES where ground = 'chinnaswamy'

0 3 9 12 6 13 35
0 3 9 12 6 13 35
display count of matches hosted at buffalo
select played from VENUES where ground = 'buffalo'
0 3 9 12 6 13 36
0 3 9 12 6 13 36

222
Name teams that participate in bordergavaskartrophy
select teams from ODISTATS where tour_name = 'bordergavaskartrophy'
0 58 60 61 59 13 62
0 58 60 61 59 13 62

223
Name teams that participate in worldcup
select teams from ODISTATS where tour_name = 'worldcup'
0 58 60 61 59 13 63
0 58 60 61 59 13 63

224
Name teams that participate in championstrophy
select teams from ODISTATS where tour_name = 'championstrophy'
0 58 60 61 59 13 64
0 58 60 61 59 13 64

225
List out teams participating in bordergavaskartrophy
select teams from ODISTATS where tour_name = 'bordergavaskartrophy'
0 58 60 61 59 13 62
0 58 60 61 59 13 62

226
List out teams participating in worldcup
select teams from ODISTATS where tour_name = 'worldcup'
0 58 60 61 59 13 63
0 58 60 61 59 13 63
227
List out teams participating in championstrophy
select teams from ODISTATS where tour_name = 'championstrophy'
0 58 60 61 59 13 64
0 58 60 61 59 13 64
228
which teams are playing in bordergavaskartrophy
select teams from ODISTATS where tour_name = 'bordergavaskartrophy'
0 58 60 61 59 13 62
0 58 60 61 59 13 62
229
which teams are playing in worldcup
select teams from ODISTATS where tour_name = 'worldcup'
0 58 60 61 59 13 63
0 58 60 61 59 13 63
230
which teams are playing in championstrophy
select teams from ODISTATS where tour_name = 'championstrophy'
0 58 60 61 59 13 64
0 58 60 61 59 13 64
231
List out nations that compete in bordergavaskartrophy
select teams from ODISTATS where tour_name = 'bordergavaskartrophy'
125
List out nations that compete in worldcup

```sql
select teams from ODISTATS where tour_name = 'worldcup'
```

List out nations that compete in championstrophy

```sql
select teams from ODISTATS where tour_name = 'championstrophy'
```

all odi series where srilanka was victorious

```sql
select tour_name from ODISTATS where ODISTATS.dot champions = 'srilanka'
```

all odi series where australia was victorious

```sql
select tour_name from ODISTATS where ODISTATS.dot champions = 'australia'
```

all odi series where westindies was victorious

```sql
select tour_name from ODISTATS where ODISTATS.dot champions = 'westindies'
```
all odi clashes where southafrica was winner
select tour_name from ODISTATS where ODISTATS dot champions = 'southafrica'

all odi clashes where newzeland was winner
select tour_name from ODISTATS where ODISTATS dot champions = 'newzeland'

all odi clashes where bangladesh was winner
select tour_name from ODISTATS where ODISTATS dot champions = 'bangladesh'

all test tournaments where srilanka was victorious
select tour_name from TESTSTATS where TESTSTATS dot champions = 'srilanka'

all test tournaments where australia was victorious
select tour_name from TESTSTATS where TESTSTATS dot champions = 'australia'
all test tournaments where westindies was victorious
select tour_name from TESTSTATS where TESTSTATS dot champions = 'westindies'
0 72 75 74 75 73 32
0 72 75 74 75 73 32

all test cups where southafrica was winner
select tour_name from TESTSTATS where TESTSTATS dot champions = 'southafrica'
0 72 75 74 75 73 27
0 72 75 74 75 73 27

all test cups where newzealand was winner
select tour_name from TESTSTATS where TESTSTATS dot champions = 'newzealand'
0 72 75 74 75 73 29
0 72 75 74 75 73 29

all test cups where bangladesh was winner
select tour_name from TESTSTATS where TESTSTATS dot champions = 'bangladesh'
0 72 75 74 75 73 30
0 72 75 74 75 73 30

Which odi team won unitycup in the year 2006
select champions from ODISTATS where tour_name = 'unitycup' and year = '2006'
0 76 60 44 77 59 13 65 78 79 13 94
0 76 60 44 77 59 13 65 78 79 13 94
Which ODI team won superseries in the year 2006
select champions from ODISTATS where tour_name = 'superseries' and year = '2006'
0 76 60 44 77 59 13 66 78 79 13 94
0 76 60 44 77 59 13 66 78 79 13 94

Which ODI team got DLFCup in the year 2006
select champions from ODISTATS where tour_name = 'dlfcup' and year = '2006'
0 76 60 44 77 59 13 67 78 79 13 94
0 76 60 44 77 59 13 67 78 79 13 94

Which ODI team got miniworldcup in the year 2006
select champions from ODISTATS where tour_name = 'miniworldcup' and year = '2006'
0 76 60 44 77 59 13 68 78 79 13 94
0 76 60 44 77 59 13 68 78 79 13 94

Which test team won hondacup in the year 2006
select champions from TESTSTATS where tour_name = 'hondacup' and year = '2006'
0 76 75 44 77 59 13 93 78 79 13 94
0 76 75 44 77 59 13 93 78 79 13 94

Which test team won ashes in the year 2001
select champions from TESTSTATS where tour_name = 'ashes' and year = '2001'
0 76 75 44 77 59 13 90 78 79 13 95
0 76 75 44 77 59 13 90 78 79 13 95
Which test team got ashes in the year 2007
select champions from TESTSTATS where tour_name = 'ashes' and year = '2007'
0 76 75 44 77 59 13 90 78 79 13 99
0 76 75 44 77 59 13 90 78 79 13 99

Which test team got ashes in the year 2008
select champions from TESTSTATS where tour_name = 'ashes' and year = '2008'
0 76 75 44 77 59 13 90 78 79 13 100
0 76 75 44 77 59 13 90 78 79 13 100

Show odi mom for first match of unitycup in year 2006
select mom from ODISTATS where tour_name = 'unitycup' and year = '2006' and matchnumber = ' first '
0 104 60 105 77 59 13 65 78 79 13 94
0 104 60 105 77 59 13 65 78 79 13 94

Show odi mom for first match of superseries in year 2006
select mom from ODISTATS where tour_name = 'superseries' and year = '2006' and matchnumber = ' first '
0 104 60 105 77 59 13 66 78 79 13 94
0 104 60 105 77 59 13 66 78 79 13 94

Find odi mom for first match of dlfcup of year 2006
select mom from ODISTATS where tour_name = 'dlfcup' and year = '2006' and matchnumber = ' first '

Find odi mom for first match of miniworldcup of year 2006
select mom from ODISTATS where tour_name = 'miniworldcup' and year = '2006'
and matchnumber = ' first ' 

Show test mom for first match of hondacup in year 2006
select mom from TESTSTATS where tour_name = 'hondacup' and year = '2006' and 
matchnumber = 'first'

Show test mom for first match of ashes in year 2006
select mom from TESTSTATS where tour_name = 'ashes' and year = '2006' and 
matchnumber = 'first'

Find test mom for first match of ashes of year 2007
select mom from TESTSTATS where tour_name = 'ashes' and year = '2007' and 
matchnumber = 'first'
Find test mom for first match of ashes of year 2008

select mom from TESTSTATS where tour_name = 'ashes' and year = '2008' and matchnumber = 'first'

0 104 75 105 77 59 13 90 78 79 13 100

0 104 75 105 77 59 13 90 78 79 13 100