

**THE CHROMATIC NUMBER OF RANDOM GRAPHS**

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# THE CHROMATIC NUMBER OF RANDOM GRAPHS

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**Abstract.** Let  $\chi(G(n, p))$  denote the chromatic number of the random graph  $G(n, p)$ . We prove that for every constant  $\epsilon > 0$  there exists a constant  $d_\epsilon$  such that for  $np(n) > d_\epsilon$ ,  $p(n) \rightarrow 0$ , the probability that

$$\frac{np}{2 \log np} \leq \chi(G(n, p)) \leq (1 + \epsilon) \frac{np}{2 \log np}$$

tends to 1 as  $n \rightarrow \infty$ .

Let  $G(n, p)$  be a random graph with vertex set  $[n] = \{1, 2, \dots, n\}$  in which each possible edge is present independently with the probability  $p = p(n)$ . We say that  $G(n, p)$  has some property *a.s.* if the probability that it has this property tends to 1 as  $n \rightarrow \infty$ . In this paper we shall consider the asymptotical behaviour of the chromatic number of  $G(n, p)$ .

The case when  $p$  is a constant,  $0 < p < 1$ , was considered in [1] and [4] (in [1]  $\chi(G(n, p))$  was determined also for  $p(n)$  which tends to 0 slowly enough i.e. for  $p(n) > n^{-\frac{1}{3} + \epsilon}$  where  $\epsilon > 0$ ). We shall solve this problem for every function  $p(n) \rightarrow 0$  such that  $np(n)$  is greater than some large constant. Our main result is the following.

**THEOREM.** For every positive constant  $\epsilon$  there exists a constant  $d_\epsilon$  such that if  $d = d(n) = np(n) > d_\epsilon$  and  $p(n) \rightarrow 0$  then *a.s.*

$$\frac{d}{2 \log d} < \chi(G(n, p)) < (1 + \epsilon) \frac{d}{2 \log d}.$$

The lower bound is an immediate consequence of the well known fact that *a.s.* the independence number of  $G(n, p)$  is less than  $\frac{2 \log d}{d} n$ . Thus it is enough to show the second inequality.

A subset  $S$  of  $G(n, p)$  is *k-independent* if it can be split into some number of disjoint independent sets each of at least  $k$  elements. We shall start with the proof that for  $k_0 = \lfloor (2 - \epsilon) \frac{n \log d}{d} \rfloor$   $G(n, p)$  contains *a.s.* large  $k_0$ -independent set.

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LEMMA 1. Let  $\epsilon > 0$ ,  $k_0 = \lfloor (2-\epsilon) \frac{n \log d}{d} \rfloor$ . Then there exists a constant  $d_\epsilon$  such that whenever  $np(n) = d(n) > d_\epsilon$  then with the probability at least  $1 - n^{-2}$   $G(n, p)$  contains  $k_0$ -independent set with at least  $nd^{-\epsilon}$  vertices.

*Proof.* Let  $z = d^{0.9\epsilon}$  and  $\mathcal{V} = \mathcal{V}(n, p, z)$  denote the partition of  $[n]$  onto sets  $V_1, \dots, V_{\frac{n}{z}}$  such that  $z - 1 \leq |V_i| \leq z + 1$  for  $i = 1, \dots, \frac{n}{z}$ . We shall call a subset  $S$  of  $G(n, p)$   $\mathcal{V}$ -disjoint if no two elements of  $S$  are contained in the same set of partition  $\mathcal{V}$ .

Let  $X$  denote the size of the largest  $k_0$ -independent  $\mathcal{V}$ -disjoint set contained in  $G(n, p)$ . Now define  $X_i$ ,  $i = 1, 2, \dots, z$ , as the expectation of the largest  $k_0$ -independent  $\mathcal{V}$ -disjoint set in a graph  $\mathcal{G}_i$  such that subgraphs induced by  $\bigcup_{r=1}^i V_r$  in  $\mathcal{G}_i$  and  $G(n, p)$  are the same and all edges of  $\mathcal{G}_i$  which are not contained in  $\bigcup_{r=1}^i V_r$  are present in this graph independently with the probability  $p$ . Then  $X = X_{\frac{n}{z}}$  and setting  $X_0 = EX$  we have

$$|X_i - X_{i-1}| \leq 1 \quad \text{for } i = 1, \dots, \frac{n}{z}.$$

Furthermore the sequence  $X_0, \dots, X_{\frac{n}{z}}$  is a martingale known as Doob's Martingale Process (see [5]). Hence, using a martingale inequality of Azuma (see [1], [5]), we obtain

$$\text{Prob} \{|X - EX| \geq t\} \leq 2 \exp\left(-\frac{zt^2}{2n}\right) = 2 \exp\left(-\frac{d^{0.9\epsilon} t^2}{2n}\right).$$

Thus, to prove Lemma 1, it is enough to show that the probability that  $G(n, p)$  contains a  $k_0$ -independent  $\mathcal{V}$ -disjoint set with more than  $2nd^{-\epsilon}$  elements is greater than  $\exp(-nd^{-1.8\epsilon})$  (this observation was first made by Alan Frieze in [2]).

Let  $Y$  be the number of  $\mathcal{V}$ -disjoint  $k_0$ -independent sets of  $mk_0$  elements,  $m = \lceil \frac{2nd^{-\epsilon}}{k_0} \rceil$ , which can be split into exactly  $m$  independent sets, each of  $k_0$  elements. Then

$$P(Y > 0) \geq \frac{(EY)^2}{EY^2}$$

and

$$\begin{aligned} \frac{EY^2}{(EY)^2} &\leq \prod_{i=1}^m \sum_{\substack{k_1, k_2, \dots, k_{m+1} \\ \sum_{j=1}^{m+1} k_j = k_0}} \frac{\binom{k_0}{k_1} \binom{k_0}{k_2} \dots \binom{k_0}{k_m} \binom{\frac{n}{z} - (i-1)k_0}{k_{m+1}} z^{k_{m+1}}}{\binom{\frac{n}{z} - (i-1)k_0}{k_0} z^{k_0} (1-p)^{\sum_{j=1}^m \binom{k_j}{2}}} \\ &\leq \left[ \sum_{l=0}^{k_0} \frac{a_l \binom{\frac{n}{z} - mk_0}{k_0 - l} z^{k_0 - l}}{\binom{\frac{n}{z} - mk_0}{k_0} z^{k_0}} \right]^m \\ &\leq \left[ \sum_{l=0}^{k_0} \frac{a_l}{(n - k_0 m z)^l} \cdot \frac{k_0!}{(k_0 - l)!} \right]^m \leq \left[ \sum_{l=0}^{k_0} a_l \left( \frac{2k_0}{n} \right)^l \right]^m \end{aligned}$$

where

$$(1) \quad a_l = \sum_{\substack{k_1, k_2, \dots, k_m \\ \sum_{j=1}^m k_j = l}} \binom{k_0}{k_1} \binom{k_0}{k_2} \cdots \binom{k_0}{k_m} (1-p)^{-\sum_{j=1}^m \binom{k_j}{2}}.$$

To estimate  $a_l$  divide all terms of the sum into two groups. We shall call a term *large* if some of indices  $k_1, k_2, \dots, k_m$  are larger than  $0.1\epsilon k_0$  and *small* otherwise.

Denote by  $k'_1, k'_2, \dots, k'_r$  those from  $k_1, k_2, \dots, k_m$  which are greater than  $0.1\epsilon k_0$ . Since  $\sum_{j=1}^m k_j = l$  so  $r \leq 10\epsilon^{-1}$  and we can choose  $k'_1, k'_2, \dots, k'_r$  in at most  $(ml)^{10\epsilon^{-1}}$  ways. Moreover, for every  $k', k''$  such that  $k' > k'' > 0.1\epsilon k_0$  and  $k' + k'' \leq l \leq k_0$ , we have (for  $d$  large enough)

$$\frac{\binom{k_0}{k'} \binom{k_0}{k''} (1-p)^{-\binom{k'}{2} - \binom{k''}{2}}}{\binom{k_0}{k' + k''} (1-p)^{-\binom{k' + k''}{2}}} \leq 2^{k_0} 2^{k_0} \exp(-k'k'') < 1,$$

so

$$\binom{k_0}{k'_1} \binom{k_0}{k'_2} \cdots \binom{k_0}{k'_r} (1-p)^{-\sum_{j=1}^r \binom{k'_j}{2}} \leq \binom{k_0}{l} \exp\left(\frac{l^2 d}{2n}\right).$$

Furthermore, for every choice of  $\bar{k}_1, \bar{k}_2, \dots, \bar{k}_s$ , one can easily get the following inequality

$$(2) \quad \sum_{\substack{\bar{k}_1, \bar{k}_2, \dots, \bar{k}_s \\ \sum_{j=1}^s \bar{k}_j = l \\ \max_{1 \leq j \leq s} \bar{k}_j = b}} \binom{k_0}{\bar{k}_1} \binom{k_0}{\bar{k}_2} \cdots \binom{k_0}{\bar{k}_s} (1-p)^{-\sum_{j=1}^s \binom{\bar{k}_j}{2}} \leq \binom{sk_0}{l} \exp\left(\frac{bld}{2n}\right).$$

Hence, (1) and (2) imply that the sum of all large terms is estimated from above by

$$(ml)^{10\epsilon^{-1}} \binom{k_0}{l} \exp\left(\frac{l^2 d}{2n}\right) \binom{mk_0}{0.1\epsilon l} \exp\left(\frac{0.1\epsilon l^2 d}{2n}\right) \leq \left(\frac{10k_0}{l} \exp\left(\frac{(1+0.2\epsilon)ld}{2n}\right)\right)^l$$

whereas the upper bound for the sum of all small terms, from (2), is given by

$$\binom{mk_0}{l} \exp\left(\frac{0.1\epsilon l^2 d}{2n}\right) \leq \left(\frac{emk_0}{l} \exp\left(\frac{0.1\epsilon ld}{2n}\right)\right)^l.$$

So, finally, we obtain

$$\frac{EY^2}{(EY)^2} \leq \left[ \sum_{l=0}^{k_0} \left( \frac{10 k_0^2 m}{ln} \exp\left(\frac{0.1\epsilon ld}{2n}\right) \right)^l + \left( \frac{20 k_0^2}{ln} \exp\left(\frac{(1+0.2\epsilon)ld}{2n}\right) \right)^l \right]^m.$$

Since the maximal term in the above sum is smaller than  $\exp\left(\frac{100k_0^2 d^{0.15\epsilon} m}{n}\right)$  we have

$$\text{Prob}(Y > 0) > \exp\left(-\frac{200k_0^2 d^{0.15\epsilon} m^2}{n}\right) > \exp(-nd^{-1.8\epsilon}).$$

This completes the proof of Lemma 1.  $\square$

LEMMA 2. Let  $0 < \epsilon < 0.1$ . Then there is a constant  $d_\epsilon$  such that for  $np(n) = d(n) > d_\epsilon$  with the probability greater than  $1 - n^{-1} - d^{-0.1\epsilon}$  more than  $n - nd^{-0.05\epsilon}$  vertices of  $G(n, p)$  can be properly coloured with less than  $(1 + 5\epsilon)\frac{d}{2\log d}$  colours.

*Proof.* In the proof we shall use an “expose-and-merge” technique introduced by David Matula (for details see [3] and [4]).

For  $A \subset [n]$  define  $[A]^2 = \{\{v, w\} : v, w \in A\}$  and let  $\bar{k}_0 = (2 - 0.05\epsilon)(1 - \epsilon)\frac{n \log d}{d}$ ,  $l_0 = \frac{nd^{-1.1\epsilon}}{\bar{k}_0}$ . Consider the following algorithm

**Algorithm**

$E := \emptyset$ ;

$F_0 := \emptyset$ ;

$W_0 := \emptyset$ ;

for  $i = 1$  to  $d^{1.1\epsilon} - d^\epsilon$  do

begin

choose randomly  $A_i \subset [n] \setminus W_{i-1}$  with  $|A_i| = nd^{-\epsilon}$ ;

define  $\mathcal{G}_i$  as the graph with the set of vertices  $A_i$  and the set of edges  $E_i$ , where the probability that  $\{v, w\} \in E_i$  is independent and equal  $p$  for each  $\{v, w\} \in [A_i]^2$ ;

choose a family  $\{\tilde{\mathcal{R}}_1^i, \tilde{\mathcal{R}}_2^i, \dots, \tilde{\mathcal{R}}_{l_0}^i\}$  of disjoint independent sets from  $A_i$ , each of the size  $\bar{k}_0$  - if it is not possible FAIL;

$E'_i := E_i \setminus (E_i \cap F_{i-1})$ ;

$E := E \cup E'_i$ ;

$F_i := F_{i-1} \cup [A_i]^2$ ;

$W_i := W_{i-1} \cup \bigcup_{l=1}^{l_0} \tilde{\mathcal{R}}_l^i$ ;

end

$$\overline{F} := [n]^2 \setminus \bigcup_{i=1}^{d^{1.1\epsilon} - d^\epsilon} F_i;$$

choose  $\overline{E} \subset \overline{F}$  in such a way that each  $e \in \overline{F}$  belong to  $\overline{E}$  with the probability  $p$  independently of each other;

$$E := E \cup \overline{E};$$

$$\text{output } E; \{\tilde{\mathcal{R}}_1^1, \tilde{\mathcal{R}}_2^1, \dots, \tilde{\mathcal{R}}_{l_0}^{d^{1.1\epsilon} - d^\epsilon}\};$$

end

Let us observe first that the probability that  $\{v, w\} \in E$  is equal  $p$  independently for each  $\{v, w\} \in [n]^2$ , so the graph  $\tilde{G}$  with the set of vertices  $[n]$  and the set of edges  $E$  can be treated as  $G(n, p)$ .

Obviously, we can consider each  $\mathcal{G}_i$  as  $G(\bar{n}, p)$  where  $\bar{n} = nd^{-\epsilon}$ . The average degree of such a random graph is given by  $\bar{d} = nd^{-\epsilon}p = d^{1-\epsilon}$  and  $\bar{k}_0 = (2 - 0.05\epsilon) \frac{\bar{n} \log \bar{d}}{\bar{d}}$ . Thus, from Lemma 1, the probability that  $\mathcal{G}_i$  does not contain  $\bar{k}_0$ -independent sets with  $nd^{-1.1\epsilon} < \bar{n} \bar{d}^{-0.05\epsilon}$  elements is less than  $n^{-2}$  so the probability of FAIL in the Algorithm is less than  $n^{-1}$ .

Thus, with the probability  $1 - o(1)$ , the Algorithm finds

$$(d^{1.1\epsilon} - d^\epsilon)l_0 = \frac{n - nd^{-0.1\epsilon}}{\bar{k}_0} < (1 + 5\epsilon) \frac{d}{2 \log d}$$

disjoint sets  $\tilde{\mathcal{R}}_1^1, \tilde{\mathcal{R}}_2^1, \dots, \tilde{\mathcal{R}}_{l_0}^{d^{1.1\epsilon} - d^\epsilon}$  each of the size  $\bar{k}_0$ , so  $\sum_{i,l} |\tilde{\mathcal{R}}_i^l| = n - nd^{-0.1\epsilon}$ . Note however, that although  $\tilde{\mathcal{R}}_i^l$  is an independent subset of  $\mathcal{G}_i$  it is *not* necessarily independent as a subset of  $\tilde{G}$ . Let  $X$  denote the number of edges of  $\tilde{G}$  contained in  $\tilde{\mathcal{R}}_i^l$  for some  $1 \leq i \leq d^{1.1\epsilon} - d^\epsilon$ ,  $1 \leq l \leq l_0$ . We shall estimate from above the size of  $X$ .

Let  $\{v, w\} \in [n]^2$  be such that  $\{v, w\} \in E$  and  $\{v, w\} \in \tilde{\mathcal{R}}_i^l$  for some  $i, l$ . Denote by  $i(v, w)$  the smallest number  $i$  for which  $\{v, w\} \subset A_i$  and let  $j(v, w), l(v, w)$  be such that  $\{v, w\} \subset \tilde{\mathcal{R}}_{l(v,w)}^{j(v,w)}$ . Notice that  $i(v, w) < j(v, w)$  since  $\{v, w\} \in E$  implies that  $\{v, w\}$  is an edge of  $\mathcal{G}_{i(v,w)}$  whereas the set  $\tilde{\mathcal{R}}_{l(v,w)}^{j(v,w)}$  contains no edges of  $\mathcal{G}_{j(v,w)}$ . Since for all  $i$  we have  $|W_i| \leq n - nd^{0.1\epsilon}$ , so the probability that for chosen  $i(v, w), j(v, w)$ , a pair  $\{v, w\}$  is contained in both  $A_{i(v,w)}$  and  $A_{j(v,w)}$  is less than  $(2d^{-0.9\epsilon})^4$ . Now observe that each subset of  $A_{j(v,w)}$  of  $k_0$  elements is equally likely to be chosen as  $\tilde{\mathcal{R}}_{l(v,w)}^{j(v,w)}$  with  $1 \leq l \leq l_0$  (i.e. this event depends only on the structure of  $\mathcal{G}_{j(v,w)}$  which is symmetric with respect to the labelling of vertices). Thus, since  $|A_{j(v,w)}| = nd^{-\epsilon}$ , the probability that both  $v, w$  are in the same set  $\tilde{\mathcal{R}}_{l(v,w)}^{j(v,w)}$  for some  $l(v, w)$  is less than  $d^{1.1\epsilon-1}$ .

So finally we arrive at the following upper bound for the expectation of  $X$

$$\mathbb{E}X \leq \binom{n}{2} \binom{d^{1.1\epsilon} - d^\epsilon}{2} 16d^{-3.6\epsilon} p d^{1.1\epsilon-1} < nd^{-0.2\epsilon}.$$

Thus, from Markov inequality,

$$\text{Prob}(X > nd^{-0.1\epsilon}) \leq d^{-0.1\epsilon}.$$

Now, for all  $i, l$ , delete from  $\tilde{\mathcal{R}}_i^l$  all these vertices which belongs to edges of  $\tilde{G}$  which are contained in  $\tilde{\mathcal{R}}_i^l$  and denote the set obtained in this way by  $\mathcal{R}_i^l$ . Then

$$\sum_{i,l} |\mathcal{R}_i^l| \geq \sum_{i,l} |\tilde{\mathcal{R}}_i^l| - 2X = n - nd^{-0.1\epsilon} - 2X$$

so

$$\text{Prob} \left( \sum_{i,l} |\mathcal{R}_i^l| \geq n - nd^{-0.05\epsilon} \right) > 1 - o(1) - d^{0.1\epsilon}.$$

This completes the proof of Lemma 2.  $\square$

*Proof of Theorem.* Lemma 2 implies that with the probability at least  $1 - o(1) - d^{-0.01\epsilon}$  we can colour  $n - nd^{-0.005\epsilon}$  vertices of  $G(n, p)$  using only  $(1 + 0.5\epsilon) \frac{d}{2 \log d}$  colours. One can easily check that for  $d$  large enough a.s. every subgraph of  $G(n, p)$  on  $s$  vertices,  $s \leq nd^{-0.005\epsilon}$ , has less than  $d^{1-0.001\epsilon} s$  edges. Hence, for large  $d$ , we have

$$(3) \quad \text{Prob} \left( \chi(G(n, p)) \leq (1 + 0.9\epsilon) \frac{d}{2 \log d} \right) > 1 - o(1) - d^{-0.005\epsilon},$$

and for  $d(n) \rightarrow \infty$ , the assertion follows.

Moreover, Shamir and Spencer proved in [5] that for every  $d(n) < \log n$  there exists a function  $u_d(n)$  such that a.s.

$$u_d(n) \leq \chi(G(n, p)) \leq u_d(n) + 5.$$

Thus, since from (3) for  $d(n)$  greater than some constant  $d_\epsilon$  we have

$$\liminf_{n \rightarrow \infty} \text{Prob} \left( \chi(G(n, p)) \leq (1 + 0.9\epsilon) \frac{d}{2 \log d} \right) > 0.5$$

so, for such  $d(n)$ , a.s.

$$\chi(G(n, p)) \leq (1 + 0.9\epsilon) \frac{d}{2 \log d} + 5 \leq (1 + \epsilon) \frac{d}{2 \log d}.$$

This completes the proof of Theorem.  $\square$

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